

Reconstructing structures from their abstract clones

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Outline

- Reconstructing structures from their automorphism groups and polymorphism clones

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- The topology of algebras

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- Reconstructing structures from their automorphism groups and polymorphism clones
- The topology of algebras
- Reconstruction the topology of function clones



Part I

Reconstructing structures from their
automorphism groups and polymorphism clones

Reconstructing structures up to first-order ...



Reconstructing structures up to first-order ...



countable

Reconstructing structures up to first-order ...



countable, ω -categorical

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Aut()

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Theorem (Ryll-Nardzewski)

Let Δ, Γ be ω -categorical structures on the same domain. Then $\text{Aut}(\Delta) = \text{Aut}(\Gamma)$ iff Δ, Γ are first-order interdefinable.

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Theorem (Ahlbrandt + Ziegler '86)

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Reconstruction from the abstract group

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- Can we reconstruct the **topological structure** of closed oligomorphic permutation groups from their **algebraic structure**?

Better reconstruction plans

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Observe: $\text{Pol}(\Delta) \supseteq \text{End}(\Delta) \supseteq \text{Aut}(\Delta)$.

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Theorem (Bodirsky + Nešetřil '03)

Let Δ, Γ be ω -categorical structures on the same domain. Then:
 $\text{Pol}(\Delta) = \text{Pol}(\Gamma)$ iff Δ, Γ are primitive positive interdefinable.

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Definition (Constraint Satisfaction Problem)

$\text{CSP}(\Delta)$ is the computational problem to decide whether a given primitive positive τ -sentence holds in Δ .

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Part II

The topology of algebras

Clones from algebras

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Structural conclusions about \mathfrak{A} from abstract clone $\text{Clo}(\mathfrak{A})$: **Varieties**.

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Theorem (Birkhoff 1935)

Let $\mathfrak{A}, \mathfrak{B}$ be finite.

\mathfrak{B} is in $\text{HSP}^{\text{fin}}(\mathfrak{A}) \leftrightarrow$

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Theorem ('Topological Birkhoff'; Bodirsky + MP '12)

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Part III

Reconstructing the topology

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Observation. (2) \implies (1).

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Fact. For groups (3) \implies (2).

Groups: the small index property

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Our favorite automorphism groups have automatic continuity:

- $(\mathbb{N}; =)$ (Dixon+Neumann+Thomas'86)
- $(\mathbb{Q}; <)$ and the atomless Boolean algebra (Truss'89)
- the random graph (Hodges+Hodkinson+Lascar+Shelah'93)
- the random K_n -free graphs (Herwig'98)

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various ω -categorical semilinear orders

the random partial order

the random tournament

(Rubin '94)

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the random tournament
(Rubin '94)
- the random k -hypergraphs
the Henson digraphs
(Barbina+MacPherson '07).



Part IV
Negative results

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Important in constraint satisfaction:

“main reason” for NP-hardness of the CSP of a structure.

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Theorem (Bodirsky + MP + Pongrácz '13)

There exists a closed oligomorphic clone with a discontinuous homomorphism to **1**.

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Theorem (Bodirsky + MP + Pongrácz '13)

There exists a closed oligomorphic clone with a discontinuous homomorphism to $\mathbf{1}$.

- Involves non-principal ultrafilter: unfair in the CSP context.
- Also has a continuous homomorphism to $\mathbf{1}$.

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Theorem (Bodirsky + MP + Pongrácz '13)

There exists a closed oligomorphic clone \mathbf{C} and $\xi: \mathbf{C} \rightarrow \mathbf{C}$ such that:

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Theorem (Evans + Hewitt '90)

There exists an closed oligomorphic group which does not have reconstruction.



Part V

Positive results

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Theorem (Bodirsky + MP + Pongrácz '13)

Any closed subclone of \mathbf{O} containing ω^ω has automatic continuity and automatic homeomorphicity.

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Theorem (Bodirsky + MP + Pongrácz '13)

Let G be the random graph.

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- All minimal tractable clones containing $\text{Aut}(G)$.



Part VI

Open problems

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- Which topological clones are closed subclones of \mathbf{O} ?









Thank you!