

# Boolean $\pm$ -preclones

Peter Jipsen  
Erkko Lehtonen  
Reinhard Pöschel

Chapman University, Orange, CA (USA)  
Khalifa University, Abu Dhabi (United Arab Emirates)  
Technische Universität Dresden (Germany)

AAA108

108. **A**rbeitstagung **A**llgemeine **A**lgebra  
108<sup>th</sup> Workshop on General Algebra  
Wien, February 6-8, 2026

## In memoriam Günther Eigenthaler

This is the first AAA attended by me since the passing of Günther Eigenthaler. Therefore I dedicate this talk to the memory of my esteemed colleague and dear friend Günther.

GÜNTHER EIGENTHALER

9.2.1950 - 14.2.2025



Günther, Laci Márki, R.P.

September 2023 during an excursion (Lunzer See) in Austria

# Outline

$\pm$ -preclones ( $S$ -preclones)

$\pm$ -relations,  $\pm$ -preservation and a Galois connection

The lattice  ${}^{\pm}\mathcal{L}_A$  of Boolean  $\pm$ -preclones

# Outline

$\pm$ -preclones ( $S$ -preclones)

$\pm$ -relations,  $\pm$ -preservation and a Galois connection

The lattice  ${}^{\pm}\mathcal{L}_A$  of Boolean  $\pm$ -preclones

## Motivating example

some history:

**Nov. 2021 PALS** talk by P. Jipsen on partially ordered algebras (and po-clones): operations which in each argument are *order-preserving* **or** *order-reversing* (for some given order on the base set).

Questions: how to characterize such “po-clones”?

R.P.: characterization via invariant relations?

Analogies to many-sorted algebras

(results of E. Lehtonen/ R. Pöschel/ T. Waldhauser),

Let  $P$  be a property for unary functions  $g \in A^A$ .

“motivating example”:  $P = +$  : order-preserving

$P = -$  : order-reversing

An  $n$ -ary operation  $f(x_1, \dots, x_n)$  has property  $P$  in an argument, say  $x_1$ ,  
 $: \iff$  each translation  $x_1 \mapsto f(x_1, c_2, \dots, c_n)$  has this property  $P$  (for all constants  $c_2, \dots, c_n \in A$ ).

How to handle composition? order-reversing composed with order-reversing is order-preserving! Formalization: Collect the properties

in a monoid  $S = (\{+, -\}, \cdot)$ , here a group

$\cdot$	$+$	$-$
$+$	$+$	$-$
$-$	$-$	$+$

## Motivating example

some history:

**Nov. 2021 PALS** talk by P. Jipsen on partially ordered algebras (and po-clones): operations which in each argument are *order-preserving* **or** *order-reversing* (for some given order on the base set).

Questions: how to characterize such “po-clones”?

R.P.: characterization via invariant relations?

Analogies to many-sorted algebras

(results of E. Lehtonen/ R. Pöschel/ T. Waldhauser),

Let  $P$  be a property for unary functions  $g \in A^A$ .

“motivating example”:  $P = +$  : order-preserving

$P = -$  : order-reversing

An  $n$ -ary operation  $f(x_1, \dots, x_n)$  has property  $P$  in an argument, say  $x_1$ ,  
 $: \iff$  each translation  $x_1 \mapsto f(x_1, c_2, \dots, c_n)$  has this property  $P$  (for all  
 constants  $c_2, \dots, c_n \in A$ ).

How to handle composition? order-reversing composed with  
 order-reversing is order-preserving! Formalization: Collect the properties

in a monoid  $S = (\{+, -\}, \cdot)$ , here a group

$\cdot$	$+$	$-$
$+$	$+$	$-$
$-$	$-$	$+$

## Motivating example

some history:

**Nov. 2021 PALS** talk by P. Jipsen on partially ordered algebras (and po-clones): operations which in each argument are *order-preserving* **or** *order-reversing* (for some given order on the base set).

Questions: how to characterize such “po-clones”?

R.P.: characterization via invariant relations?

Analogies to many-sorted algebras

(results of E. Lehtonen/ R. Pöschel/ T. Waldhauser),

Let  $P$  be a property for unary functions  $g \in A^A$ .

“motivating example”:  $P = +$  : order-preserving

$P = -$  : order-reversing

An  $n$ -ary operation  $f(x_1, \dots, x_n)$  has property  $P$  in an argument, say  $x_1$ ,  
 $: \iff$  each translation  $x_1 \mapsto f(x_1, c_2, \dots, c_n)$  has this property  $P$  (for all constants  $c_2, \dots, c_n \in A$ ).

How to handle composition? order-reversing composed with order-reversing is order-preserving! Formalization: Collect the properties

in a monoid  $S = (\{+, -\}, \cdot)$ , here a group

$\cdot$	$+$	$-$
$+$	$+$	$-$
$-$	$-$	$+$

## Motivating example

some history:

**Nov. 2021 PALS** talk by P. Jipsen on partially ordered algebras (and po-clones): operations which in each argument are *order-preserving* **or** *order-reversing* (for some given order on the base set).

Questions: how to characterize such “po-clones”?

R.P.: characterization via invariant relations?

Analogies to many-sorted algebras

(results of E. Lehtonen/ R. Pöschel/ T. Waldhauser),

Let  $P$  be a property for unary functions  $g \in A^A$ .

“motivating example”:  $P = +$  : order-preserving

$P = -$  : order-reversing

An  $n$ -ary operation  $f(x_1, \dots, x_n)$  has property  $P$  in an argument, say  $x_1$ ,  
 $:\iff$  each translation  $x_1 \mapsto f(x_1, c_2, \dots, c_n)$  has this property  $P$  (for all constants  $c_2, \dots, c_n \in A$ ).

How to handle composition? order-reversing composed with order-reversing is order-preserving! Formalization: Collect the properties

in a monoid  $S = (\{+, -\}, \cdot)$ , here a group

$\cdot$	$+$	$-$
$+$	$+$	$-$
$-$	$-$	$+$



## Motivating example

some history:

**Nov. 2021 PALS** talk by P. Jipsen on partially ordered algebras (and po-clones): operations which in each argument are *order-preserving* **or** *order-reversing* (for some given order on the base set).

Questions: how to characterize such “po-clones”?

R.P.: characterization via invariant relations?

Analogies to many-sorted algebras

(results of E. Lehtonen/ R. Pöschel/ T. Waldhauser),

Let  $P$  be a property for unary functions  $g \in A^A$ .

“motivating example”:  $P = +$  : order-preserving

$P = -$  : order-reversing

An  $n$ -ary operation  $f(x_1, \dots, x_n)$  has property  $P$  in an argument, say  $x_1$ ,  
:  $\iff$  each translation  $x_1 \mapsto f(x_1, c_2, \dots, c_n)$  has this property  $P$  (for all constants  $c_2, \dots, c_n \in A$ ).

How to handle composition? order-reversing composed with order-reversing is order-preserving! Formalization: Collect the properties

in a monoid  $S = (\{+, -\}, \cdot)$ , here a group

$\cdot$	$+$	$-$
$+$	$+$	$-$
$-$	$-$	$+$

## Motivating example

some history:

**Nov. 2021 PALS** talk by P. Jipsen on partially ordered algebras (and po-clones): operations which in each argument are *order-preserving* **or** *order-reversing* (for some given order on the base set).

Questions: how to characterize such “po-clones”?

R.P.: characterization via invariant relations?

Analogies to many-sorted algebras

(results of E. Lehtonen/ R. Pöschel/ T. Waldhauser),

Let  $P$  be a property for unary functions  $g \in A^A$ .

“motivating example”:  $P = +$  : order-preserving

$P = -$  : order-reversing

An  $n$ -ary *operation*  $f(x_1, \dots, x_n)$  *has property*  $P$  *in an argument*, say  $x_1$ ,  
 $: \iff$  each translation  $x_1 \mapsto f(x_1, c_2, \dots, c_n)$  has this property  $P$  (for all constants  $c_2, \dots, c_n \in A$ ).

How to handle composition? order-reversing composed with order-reversing is order-preserving! Formalization: Collect the properties

in a monoid  $S = (\{+, -\}, \cdot)$ , here a group

$\cdot$	$+$	$-$
$+$	$+$	$-$
$-$	$-$	$+$

## Motivating example

some history:

**Nov. 2021 PALS** talk by P. Jipsen on partially ordered algebras (and po-clones): operations which in each argument are *order-preserving* **or** *order-reversing* (for some given order on the base set).

Questions: how to characterize such “po-clones”?

R.P.: characterization via invariant relations?

Analogies to many-sorted algebras

(results of E. Lehtonen/ R. Pöschel/ T. Waldhauser),

Let  $P$  be a property for unary functions  $g \in A^A$ .

“motivating example”:  $P = +$  : order-preserving

$P = -$  : order-reversing

An  $n$ -ary *operation*  $f(x_1, \dots, x_n)$  *has property*  $P$  *in an argument*, say  $x_1$ ,  
:  $\iff$  each translation  $x_1 \mapsto f(x_1, c_2, \dots, c_n)$  has this property  $P$  (for all constants  $c_2, \dots, c_n \in A$ ).

How to handle composition? order-reversing composed with order-reversing is order-preserving! Formalization: Collect the properties

in a monoid  $S = (\{+, -\}, \cdot)$ , here a group

$\cdot$	$+$	$-$
$+$	$+$	$-$
$-$	$-$	$+$

## Motivating example

some history:

**Nov. 2021 PALS** talk by P. Jipsen on partially ordered algebras (and po-clones): operations which in each argument are *order-preserving* **or** *order-reversing* (for some given order on the base set).

Questions: how to characterize such “po-clones”?

R.P.: characterization via invariant relations?

Analogies to many-sorted algebras

(results of E. Lehtonen/ R. Pöschel/ T. Waldhauser),

Let  $P$  be a property for unary functions  $g \in A^A$ .

“motivating example”:  $P = +$  : order-preserving

$P = -$  : order-reversing

An  $n$ -ary *operation*  $f(x_1, \dots, x_n)$  *has property*  $P$  *in an argument*, say  $x_1$ ,  
:  $\iff$  each translation  $x_1 \mapsto f(x_1, c_2, \dots, c_n)$  has this property  $P$  (for all constants  $c_2, \dots, c_n \in A$ ).

How to handle composition? order-reversing composed with

order-reversing is order-preserving! Formalization: Collect the properties

in a monoid  $S = (\{+, -\}, \cdot)$ , here a group

$\cdot$	$+$	$-$
$+$	$+$	$-$
$-$	$-$	$+$

## Motivating example

some history:

**Nov. 2021 PALS** talk by P. Jipsen on partially ordered algebras (and po-clones): operations which in each argument are *order-preserving* **or** *order-reversing* (for some given order on the base set).

Questions: how to characterize such “po-clones”?

R.P.: characterization via invariant relations?

Analogies to many-sorted algebras

(results of E. Lehtonen/ R. Pöschel/ T. Waldhauser),

Let  $P$  be a property for unary functions  $g \in A^A$ .

“motivating example”:  $P = +$  : order-preserving

$P = -$  : order-reversing

An  $n$ -ary *operation*  $f(x_1, \dots, x_n)$  *has property*  $P$  *in an argument*, say  $x_1$ ,  
:  $\iff$  each translation  $x_1 \mapsto f(x_1, c_2, \dots, c_n)$  has this property  $P$  (for all constants  $c_2, \dots, c_n \in A$ ).

How to handle composition? order-reversing composed with order-reversing is order-preserving! Formalization: Collect the properties

in a monoid  $S = (\{+, -\}, \cdot)$ , here a group

$\cdot$	$+$	$-$
$+$	$+$	$-$
$-$	$-$	$+$

## $\pm$ -operations ( $S$ -operations)

$S := \pm := \{+, -\}$  (in general, finite monoid  $S$ ).

*$n$ -ary  $S$ -operation ( $\pm$ -operation): operation  $f$  together with its signum*

$$f: A^n \rightarrow A \text{ with } \text{sgn}(f) = (s_1, \dots, s_n) \in S^n,$$

i.e., the  $i$ -th argument of  $f$  gets a label (*sign*)  $s_i \in S$  ( $i = 1, \dots, n$ ).

${}^{\pm}\text{Op}(A) :=$  all finitary  $\pm$ -operations

## $\pm$ -operations ( $S$ -operations)

$S := \pm := \{+, -\}$  (in general, finite monoid  $S$ ).

$n$ -ary  $S$ -operation ( $\pm$ -operation): operation  $f$  together with its *signum*

$$f: A^n \rightarrow A \text{ with } \text{sgn}(f) = (s_1, \dots, s_n) \in S^n,$$

i.e., the  $i$ -th argument of  $f$  gets a label (*sign*)  $s_i \in S$  ( $i = 1, \dots, n$ ).

${}^{\pm}\text{Op}(A) :=$  all finitary  $\pm$ -operations

## $\pm$ -operations ( $S$ -operations)

$S := \pm := \{+, -\}$  (in general, finite monoid  $S$ ).

$n$ -ary  $S$ -operation ( $\pm$ -operation): operation  $f$  together with its *signum*

$$f: A^n \rightarrow A \text{ with } \text{sgn}(f) = (s_1, \dots, s_n) \in S^n,$$

i.e., the  $i$ -th argument of  $f$  gets a label (*sign*)  $s_i \in S$  ( $i = 1, \dots, n$ ).

$${}^{\pm}\text{Op}(A) := \text{all finitary } \pm\text{-operations}$$



## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

- (1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),
- (2) permutation of arguments (operations  $\zeta, \tau$ ),
- (3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),
- (4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,  
e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where  
 $\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,
- (5) “linearized” composition  
 $\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

*Boolean*  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

(1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),

(2) permutation of arguments (operations  $\zeta, \tau$ ),

(3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),

(4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,

e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where

$\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,

(5) “linearized” composition

$\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

*Boolean*  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

(1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),

(2) permutation of arguments (operations  $\zeta, \tau$ ),

(3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),

(4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,

e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where

$\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,

(5) “linearized” composition

$\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

*Boolean*  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

(1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),

(2) permutation of arguments (operations  $\zeta, \tau$ ),

(3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),

(4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,

e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where

$\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,

(5) “linearized” composition

$\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

*Boolean*  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

- (1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),
- (2) permutation of arguments (operations  $\zeta, \tau$ ),
- (3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),

- (4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,  
e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where  
 $\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,

- (5) “linearized” composition  
 $\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

*Boolean*  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

- (1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),
- (2) permutation of arguments (operations  $\zeta, \tau$ ),
- (3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),
- (4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,  
e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where  
 $\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,

- (5) “linearized” composition

$\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

*Boolean*  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

- (1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),
- (2) permutation of arguments (operations  $\zeta, \tau$ ),
- (3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),
- (4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,  
e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where  
 $\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,

(5) “linearized” composition

$\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

Boolean  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

- (1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  ( $+$  unit element of  $S$ ),
- (2) permutation of arguments (operations  $\zeta, \tau$ ),
- (3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),
- (4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,  
e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where  
 $\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,
- (5) “linearized” composition

$\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

Boolean  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]



## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

- (1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),
- (2) permutation of arguments (operations  $\zeta, \tau$ ),
- (3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),
- (4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,  
e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where  
 $\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,
- (5) “linearized” composition  
 $\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

Boolean  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

- (1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),
- (2) permutation of arguments (operations  $\zeta, \tau$ ),
- (3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),
- (4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,  
e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where  
 $\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,
- (5) “linearized” composition  
 $\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

Boolean  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

- (1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),
- (2) permutation of arguments (operations  $\zeta, \tau$ ),
- (3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),
- (4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,  
e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where  
 $\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,
- (5) “linearized” composition  
 $\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

Boolean  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

- (1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),
- (2) permutation of arguments (operations  $\zeta, \tau$ ),
- (3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),
- (4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,  
e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where  
 $\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,
- (5) “linearized” composition  
 $\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

Boolean  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

- (1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  ( $+$  unit element of  $S$ ),
- (2) permutation of arguments (operations  $\zeta, \tau$ ),
- (3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),
- (4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,  
e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where  
 $\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,
- (5) “linearized” composition  
 $\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

**Boolean**  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone  $:=$  set  $F \subseteq {}^{\pm}\text{Op}(A)$  of  $\pm$ -operations closed under:

- (1)  $\text{id}_A \in F$ ,  $\text{id}_A(x) = x$ ,  $\text{sgn}(\text{id}_A) := (+)$  (+ unit element of  $S$ ),
- (2) permutation of arguments (operations  $\zeta, \tau$ ),
- (3) identification of arguments *with the same sign*  $s$  ( $\Delta^s$ ),
- (4) adding fictitious arguments of (arbitrary) sign  $s \in S$ ,  
e.g.,  $(\nabla^s f)(x_1, x_2, \dots, x_{n+1}) := f(x_2, \dots, x_{n+1})$ , where  
 $\text{sgn}(\nabla^s f) = (s, s_1, \dots, s_n)$  for  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,
- (5) “linearized” composition  
 $\text{sgn}(f) = (s_1, \dots, s_n)$  and  $\text{sgn}(g) = (s'_1, \dots, s'_m)$ . Then

$$\begin{aligned} (f \circ g)(x_1, \dots, x_m, x_{m+1}, \dots, x_{m+n-1}) \\ := f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}) \end{aligned}$$

with  $\text{sgn}(f \circ g) = (s'_1 s_1, \dots, s'_m s_1, s_2, \dots, s_n)$ .

${}^{\pm}\langle F \rangle :=$   $\pm$ -preclone generated by  $F \subseteq {}^{\pm}\text{Op}(A)$ .

**Boolean**  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = **operad**]

## Example (composition)

properties  $S = \{+, -\}$ ,  $A = \{0, 1\}$  with order  $0 < 1$ ,  
 $+$  means order preserving,  $-$  means order reversing  
 (such functions really form a  $\pm$ -preclone).

Composition:

$$f(x_1, x_2) = \neg x_1 \wedge x_2, \operatorname{sgn}(f) = (s_1, s_2) = (-, +),$$

$$g(x_1, x_2) = x_1 \vee \neg x_2, \operatorname{sgn}(g) = (s'_1, s'_2) = (+, -).$$

$$(f \circ g)(x_1, x_2, x_3) = f(g(x_1, x_2), x_3) = \neg(x_1 \vee \neg x_2) \wedge x_3 = \neg x_1 \wedge x_2 \wedge x_3$$

$$\text{has signum } (s'_1 s_1, s'_2 s_1, s_2) = (+ \cdot -, - \cdot -, +) = (-, +, +).$$

One is allowed to identify  $x_2$  and  $x_3$ , but not  $x_2$  and  $x_1$ .

## Example (composition)

properties  $S = \{+, -\}$ ,  $A = \{0, 1\}$  with order  $0 < 1$ ,  
 $+$  means order preserving,  $-$  means order reversing  
 (such functions really form a  $\pm$ -preclone).

Composition:

$$f(x_1, x_2) = \neg x_1 \wedge x_2, \text{ sgn}(f) = (s_1, s_2) = (-, +),$$

$$g(x_1, x_2) = x_1 \vee \neg x_2, \text{ sgn}(g) = (s'_1, s'_2) = (+, -).$$

$$(f \circ g)(x_1, x_2, x_3) = f(g(x_1, x_2), x_3) = \neg(x_1 \vee \neg x_2) \wedge x_3 = \neg x_1 \wedge x_2 \wedge x_3$$

$$\text{has signum } (s'_1 s_1, s'_2 s_1, s_2) = (+ \cdot -, - \cdot -, +) = (-, +, +).$$

One is allowed to identify  $x_2$  and  $x_3$ , but not  $x_2$  and  $x_1$ .



## Example (composition)

properties  $S = \{+, -\}$ ,  $A = \{0, 1\}$  with order  $0 < 1$ ,  
 $+$  means order preserving,  $-$  means order reversing  
 (such functions really form a  $\pm$ -preclone).

Composition:

$$f(x_1, x_2) = \neg x_1 \wedge x_2, \operatorname{sgn}(f) = (s_1, s_2) = (-, +),$$

$$g(x_1, x_2) = x_1 \vee \neg x_2, \operatorname{sgn}(g) = (s'_1, s'_2) = (+, -).$$

$$(f \circ g)(x_1, x_2, x_3) = f(g(x_1, x_2), x_3) = \neg(x_1 \vee \neg x_2) \wedge x_3 = \neg x_1 \wedge x_2 \wedge x_3$$

$$\text{has signum } (s'_1 s_1, s'_2 s_1, s_2) = (+ \cdot -, - \cdot -, +) = (-, +, +).$$

One is allowed to identify  $x_2$  and  $x_3$ , but not  $x_2$  and  $x_1$ .

## Example (composition)

properties  $S = \{+, -\}$ ,  $A = \{0, 1\}$  with order  $0 < 1$ ,  
 $+$  means order preserving,  $-$  means order reversing  
 (such functions really form a  $\pm$ -preclone).

Composition:

$$f(x_1, x_2) = \neg x_1 \wedge x_2, \text{sgn}(f) = (s_1, s_2) = (-, +),$$

$$g(x_1, x_2) = x_1 \vee \neg x_2, \text{sgn}(g) = (s'_1, s'_2) = (+, -).$$

$$(f \circ g)(x_1, x_2, x_3) = f(g(x_1, x_2), x_3) = \neg(x_1 \vee \neg x_2) \wedge x_3 = \neg x_1 \wedge x_2 \wedge x_3$$

has signum  $(s'_1 s_1, s'_2 s_1, s_2) = (+ \cdot -, - \cdot -, +) = (-, +, +)$ .

One is allowed to identify  $x_2$  and  $x_3$ , but not  $x_2$  and  $x_1$ .

# Outline

$\pm$ -preclones ( $S$ -preclones)

$\pm$ -relations,  $\pm$ -preservation and a Galois connection

The lattice  ${}^{\pm}\mathcal{L}_A$  of Boolean  $\pm$ -preclones

# $\pm$ -relations ( $S$ -relations) and $\pm$ -preservation $\triangleright^{\pm}$

recall:  $S = \pm := \{+, -\}$

$m$ -ary  $\pm$ -relation:  $\varrho = (\varrho_+, \varrho_-)$  with  $\varrho_s \subseteq A^m$  ( $s \in S$ )

classical notion of preservation:  $f \triangleright \varrho : \iff f(\varrho, \dots, \varrho) \subseteq \varrho$

The “ $S$ -version”:

$f \in {}^{\pm}\text{Op}^{(n)}(A)$ ,  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,  $\varrho = (\varrho_+, \varrho_-) \in {}^{\pm}\text{Rel}^{(m)}(A)$

$f \triangleright^{\pm} (\varrho_+, \varrho_-) : \iff f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_+$  and  $f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_-$

$f \triangleright^{\pm} \varrho$ :  $f$   $\pm$ -preserves  $\varrho$ ,  $f$  is an  $\pm$ -polymorphism of  $\varrho$ ,  $\varrho$  is  $(\pm)$ -invariant for  $f$

# $\pm$ -relations ( $S$ -relations) and $\pm$ -preservation $\triangleright^{\pm}$

recall:  $S = \pm := \{+, -\}$

$m$ -ary  **$\pm$ -relation**:  $\varrho = (\varrho_+, \varrho_-)$  with  $\varrho_s \subseteq A^m$  ( $s \in S$ )

classical notion of preservation:  $f \triangleright \varrho : \iff f(\varrho, \dots, \varrho) \subseteq \varrho$

The “ $S$ -version”:

$f \in {}^{\pm}\text{Op}^{(n)}(A)$ ,  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,  $\varrho = (\varrho_+, \varrho_-) \in {}^{\pm}\text{Rel}^{(m)}(A)$

$f \triangleright^{\pm} (\varrho_+, \varrho_-) : \iff f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_+$  and  $f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_-$

$f \triangleright^{\pm} \varrho$ :  $f$   $\pm$ -preserves  $\varrho$ ,  $f$  is an  $\pm$ -polymorphism of  $\varrho$ ,  $\varrho$  is  $(\pm)$ -invariant for  $f$

# $\pm$ -relations ( $S$ -relations) and $\pm$ -preservation $\triangleright^{\pm}$

recall:  $S = \pm := \{+, -\}$

$m$ -ary  $\pm$ -relation:  $\varrho = (\varrho_+, \varrho_-)$  with  $\varrho_s \subseteq A^m$  ( $s \in S$ )

classical notion of preservation:  $f \triangleright \varrho : \iff f(\varrho, \dots, \varrho) \subseteq \varrho$

The “ $S$ -version”:

$f \in {}^{\pm}\text{Op}^{(n)}(A)$ ,  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,  $\varrho = (\varrho_+, \varrho_-) \in {}^{\pm}\text{Rel}^{(m)}(A)$

$f \triangleright^{\pm} (\varrho_+, \varrho_-) : \iff f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_+$  and  $f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_-$

$f \triangleright^{\pm} \varrho$ :  $f$   $\pm$ -preserves  $\varrho$ ,  $f$  is an  $\pm$ -polymorphism of  $\varrho$ ,  $\varrho$  is  $(\pm)$ -invariant for  $f$

# $\pm$ -relations ( $S$ -relations) and $\pm$ -preservation $\triangleright^{\pm}$

recall:  $S = \pm := \{+, -\}$

$m$ -ary  $\pm$ -relation:  $\varrho = (\varrho_+, \varrho_-)$  with  $\varrho_s \subseteq A^m$  ( $s \in S$ )

classical notion of preservation:  $f \triangleright \varrho : \iff f(\varrho, \dots, \varrho) \subseteq \varrho$

The “ $S$ -version”:

$f \in {}^{\pm}\text{Op}^{(n)}(A)$ ,  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,  $\varrho = (\varrho_+, \varrho_-) \in {}^{\pm}\text{Rel}^{(m)}(A)$

$f \triangleright^{\pm} (\varrho_+, \varrho_-) : \iff f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_+$  and  $f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_-$

$f \triangleright^{\pm} \varrho$ :  $f$   $\pm$ -preserves  $\varrho$ ,  $f$  is an  $\pm$ -polymorphism of  $\varrho$ ,  $\varrho$  is  $(\pm)$ -invariant for  $f$

## $\pm$ -relations ( $S$ -relations) and $\pm$ -preservation $\triangleright^{\pm}$

recall:  $S = \pm := \{+, -\}$

$m$ -ary  $\pm$ -relation:  $\varrho = (\varrho_+, \varrho_-)$  with  $\varrho_s \subseteq A^m$  ( $s \in S$ )

classical notion of preservation:  $f \triangleright \varrho : \iff f(\varrho, \dots, \varrho) \subseteq \varrho$

The “ $S$ -version”:

$f \in {}^{\pm}\text{Op}^{(n)}(A)$ ,  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,  $\varrho = (\varrho_+, \varrho_-) \in {}^{\pm}\text{Rel}^{(m)}(A)$

$f \triangleright^{\pm} (\varrho_+, \varrho_-) : \iff f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_+$  and  $f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_-$

$f \triangleright^{\pm} \varrho$ :  $f$   $\pm$ -preserves  $\varrho$ ,  $f$  is an  $\pm$ -polymorphism of  $\varrho$ ,  $\varrho$  is  $(\pm)$ -invariant for  $f$



# $\pm$ -relations ( $S$ -relations) and $\pm$ -preservation $\triangleright^\pm$

recall:  $S = \pm := \{+, -\}$

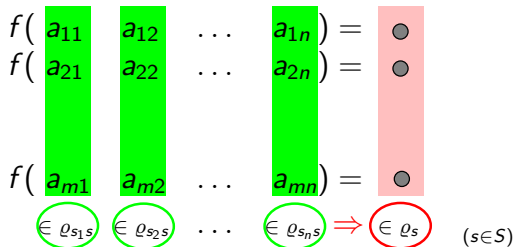
$m$ -ary  $\pm$ -relation:  $\varrho = (\varrho_+, \varrho_-)$  with  $\varrho_s \subseteq A^m$  ( $s \in S$ )

classical notion of preservation:  $f \triangleright \varrho : \iff f(\varrho, \dots, \varrho) \subseteq \varrho$

The “ $S$ -version”:

$f \in \pm\text{Op}^{(n)}(A)$ ,  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,  $\varrho = (\varrho_+, \varrho_-) \in \pm\text{Rel}^{(m)}(A)$

$f \triangleright^\pm (\varrho_+, \varrho_-) : \iff f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_+$  and  $f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_-$



$f \triangleright^\pm \varrho$ :  $f$   $\pm$ -preserves  $\varrho$ ,  $f$  is an  $\pm$ -polymorphism of  $\varrho$ ,  $\varrho$  is  $(\pm)$ -invariant for  $f$

# $\pm$ -relations ( $S$ -relations) and $\pm$ -preservation $\triangleright^\pm$

recall:  $S = \pm := \{+, -\}$

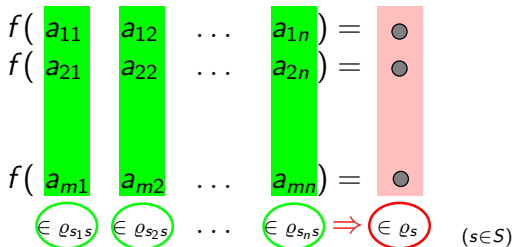
$m$ -ary  $\pm$ -relation:  $\varrho = (\varrho_+, \varrho_-)$  with  $\varrho_s \subseteq A^m$  ( $s \in S$ )

classical notion of preservation:  $f \triangleright \varrho : \iff f(\varrho, \dots, \varrho) \subseteq \varrho$

The “ $S$ -version”:

$f \in \pm\text{Op}^{(n)}(A)$ ,  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,  $\varrho = (\varrho_+, \varrho_-) \in \pm\text{Rel}^{(m)}(A)$

$f \triangleright^\pm (\varrho_+, \varrho_-) : \iff f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_+$  and  $f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_-$



$f \triangleright^\pm \varrho$ :  $f$   $\pm$ -preserves  $\varrho$ ,  $f$  is an  $\pm$ -polymorphism of  $\varrho$ ,  $\varrho$  is  $(\pm)$ -invariant for  $f$

## The Galois connection ${}^{\pm}\text{Pol} - {}^{\pm}\text{Inv}$

${}^{\pm}\triangleright$  induces a Galois connection with the operators

${}^{\pm}\text{Pol } Q := \{f \in {}^{\pm}\text{Op}(A) \mid \forall \varrho \in Q: f \triangleright^{\pm} \varrho\}$  ( $\pm$ -polymorphisms),

${}^{\pm}\text{Inv } F := \{\varrho \in {}^{\pm}\text{Rel}(A) \mid \forall f \in F: f \triangleright^{\pm} \varrho\}$  (invariant  $\pm$ -relations).

for  $F \subseteq {}^{\pm}\text{Op}(A)$  and  $Q \subseteq {}^{\pm}\text{Rel}(A)$ .

Theorem (The Galois closures)

${}^{\pm}\langle F \rangle = {}^{\pm}\text{Pol } {}^{\pm}\text{Inv } F$  ( $\pm$ -preclone generated by  $F$ ),

${}^{\pm}[Q] = {}^{\pm}\text{Inv } {}^{\pm}\text{Pol } Q$  ( $\pm$ -relational clone generated by  $Q$ ).

[JipLP2023]: *S-preclones and the Galois connection  ${}^S\text{Pol} - {}^S\text{Inv}$ , Part I*,  
*Algebra Universalis* 85, 2024

# The Galois connection ${}^{\pm}\text{Pol} - {}^{\pm}\text{Inv}$

${}^{\pm}\triangleright$  induces a Galois connection with the operators

${}^{\pm}\text{Pol } Q := \{f \in {}^{\pm}\text{Op}(A) \mid \forall \varrho \in Q: f \stackrel{\pm}{\triangleright} \varrho\}$  ( $\pm$ -polymorphisms),

${}^{\pm}\text{Inv } F := \{\varrho \in {}^{\pm}\text{Rel}(A) \mid \forall f \in F: f \stackrel{\pm}{\triangleright} \varrho\}$  (invariant  $\pm$ -relations).

for  $F \subseteq {}^{\pm}\text{Op}(A)$  and  $Q \subseteq {}^{\pm}\text{Rel}(A)$ .

Theorem (The Galois closures)

${}^{\pm}\langle F \rangle = {}^{\pm}\text{Pol } {}^{\pm}\text{Inv } F$  ( $\pm$ -preclone generated by  $F$ ),

${}^{\pm}[Q] = {}^{\pm}\text{Inv } {}^{\pm}\text{Pol } Q$  ( $\pm$ -relational clone generated by  $Q$ ).

[JipLP2023]: *S-preclones and the Galois connection  ${}^S\text{Pol} - {}^S\text{Inv}$ , Part I*,  
*Algebra Universalis* 85, 2024

## The Galois connection ${}^{\pm}\text{Pol} - {}^{\pm}\text{Inv}$

${}^{\pm}\triangleright$  induces a Galois connection with the operators

${}^{\pm}\text{Pol } Q := \{f \in {}^{\pm}\text{Op}(A) \mid \forall \varrho \in Q: f \triangleright^{\pm} \varrho\}$  ( $\pm$ -polymorphisms),

${}^{\pm}\text{Inv } F := \{\varrho \in {}^{\pm}\text{Rel}(A) \mid \forall f \in F: f \triangleright^{\pm} \varrho\}$  (invariant  $\pm$ -relations).

for  $F \subseteq {}^{\pm}\text{Op}(A)$  and  $Q \subseteq {}^{\pm}\text{Rel}(A)$ .

### Theorem (The Galois closures)

$${}^{\pm}\langle F \rangle = {}^{\pm}\text{Pol } {}^{\pm}\text{Inv } F \quad (\pm\text{-preclone generated by } F),$$

$${}^{\pm}[Q] = {}^{\pm}\text{Inv } {}^{\pm}\text{Pol } Q \quad (\pm\text{-relational clone generated by } Q).$$

[JipLP2023]: *S-preclones and the Galois connection  ${}^S\text{Pol} - {}^S\text{Inv}$ , Part I, Algebra Universalis 85, 2024*

## once more: our “motivating” Example

$(A, \leq)$  poset ( $0 < 1$ ),  $S = \{+, -\}$  (group).

For the  $\pm$ -preclone  $F$  (of  $\pm$ -operations where  $+$  means order preserving und  $-$  means order reversing) we have the following relational characterization:

$$F = {}^{\pm}\text{Pol } \varrho \text{ for the } \pm\text{-relation } \varrho = (\varrho_+, \varrho_-) := (\leq, \geq).$$

Example: For  $g(x_1, x_2) = x_1 \vee \neg x_2$ ,  $\text{sgn}(g) = (+, -)$ ,

$g \stackrel{\pm}{\triangleright} \varrho$  implies

$g\left(\begin{smallmatrix} 0 \\ 1 \end{smallmatrix}, \begin{smallmatrix} c \\ c \end{smallmatrix}\right) \in f(\varrho_+, \varrho_-) \subseteq \varrho_+ = \left(\begin{smallmatrix} 0 & 0 & 1 \\ 0 & 1 & 1 \end{smallmatrix}\right)$ , i.e., order-preserving in  $x_1$ ,

$g\left(\begin{smallmatrix} c \\ c \end{smallmatrix}, \begin{smallmatrix} 0 \\ 1 \end{smallmatrix}\right) \in f(\varrho_-, \varrho_+) = f(\varrho_+, -, \varrho_+, -) \subseteq \varrho_- = \left(\begin{smallmatrix} 0 & 1 & 1 \\ 0 & 0 & 1 \end{smallmatrix}\right)$ , i.e., order-reversing in  $x_2$ .

## once more: our “motivating” Example

$(A, \leq)$  poset ( $0 < 1$ ),  $S = \{+, -\}$  (group).

For the  $\pm$ -preclone  $F$  (of  $\pm$ -operations where  $+$  means order preserving und  $-$  means order reversing) we have the following relational characterization:

$$F = {}^{\pm}\text{Pol } \varrho \text{ for the } \pm\text{-relation } \varrho = (\varrho_+, \varrho_-) := (\leq, \geq).$$

Example: For  $g(x_1, x_2) = x_1 \vee \neg x_2$ ,  $\text{sgn}(g) = (+, -)$ ,

$g \stackrel{\pm}{\triangleright} \varrho$  implies

$g(\begin{smallmatrix} 0 \\ 1 \end{smallmatrix}, \begin{smallmatrix} c \\ c \end{smallmatrix}) \in f(\varrho_+, \varrho_-) \subseteq \varrho_+ = (\begin{smallmatrix} 0 & 0 \\ 0 & 1 \end{smallmatrix} \begin{smallmatrix} 1 \\ 1 \end{smallmatrix})$ , i.e., order-preserving in  $x_1$ ,

$g(\begin{smallmatrix} c \\ c \end{smallmatrix}, \begin{smallmatrix} 0 \\ 1 \end{smallmatrix}) \in f(\varrho_-, \varrho_+) = f(\varrho_+, -, \varrho_+ \dots) \subseteq \varrho_- = (\begin{smallmatrix} 0 & 1 \\ 0 & 0 \end{smallmatrix} \begin{smallmatrix} 1 \\ 1 \end{smallmatrix})$ , i.e., order-reversing in  $x_2$ .

## once more: our “motivating” Example

$(A, \leq)$  poset ( $0 < 1$ ),  $S = \{+, -\}$  (group).

For the  $\pm$ -preclone  $F$  (of  $\pm$ -operations where  $+$  means order preserving und  $-$  means order reversing) we have the following relational characterization:

$$F = {}^{\pm}\text{Pol } \varrho \text{ for the } \pm\text{-relation } \varrho = (\varrho_+, \varrho_-) := (\leq, \geq).$$

Example: For  $g(x_1, x_2) = x_1 \vee \neg x_2$ ,  $\text{sgn}(g) = (+, -)$ ,

$g \stackrel{\pm}{\triangleright} \varrho$  implies

$g\left(\begin{smallmatrix} 0 \\ 1 \end{smallmatrix}, \begin{smallmatrix} c \\ c \end{smallmatrix}\right) \in f(\varrho_+, \varrho_-) \subseteq \varrho_+ = \left(\begin{smallmatrix} 0 & 0 \\ 0 & 1 \end{smallmatrix}\right)$ , i.e., order-preserving in  $x_1$ ,

$g\left(\begin{smallmatrix} c \\ c \end{smallmatrix}, \begin{smallmatrix} 0 \\ 1 \end{smallmatrix}\right) \in f(\varrho_-, \varrho_+) = f(\varrho_+, -, \varrho_+ -) \subseteq \varrho_- = \left(\begin{smallmatrix} 0 & 1 \\ 0 & 0 \end{smallmatrix}\right)$ , i.e.,  
order-reversing in  $x_2$ .



## once more: our “motivating” Example

$(A, \leq)$  poset ( $0 < 1$ ),  $S = \{+, -\}$  (group).

For the  $\pm$ -preclone  $F$  (of  $\pm$ -operations where  $+$  means order preserving und  $-$  means order reversing) we have the following relational characterization:

$$F = {}^{\pm}\text{Pol } \varrho \text{ for the } \pm\text{-relation } \varrho = (\varrho_+, \varrho_-) := (\leq, \geq).$$

Example: For  $g(x_1, x_2) = x_1 \vee \neg x_2$ ,  $\text{sgn}(g) = (+, -)$ ,

$g \stackrel{\pm}{\triangleright} \varrho$  implies

$g(\begin{smallmatrix} 0 \\ 1 \end{smallmatrix}, \begin{smallmatrix} c \\ c \end{smallmatrix}) \in f(\varrho_+, \varrho_-) \subseteq \varrho_+ = (\begin{smallmatrix} 0 & 0 \\ 0 & 1 \end{smallmatrix} \begin{smallmatrix} 1 \\ 1 \end{smallmatrix})$ , i.e., **order-preserving in  $x_1$** ,

$g(\begin{smallmatrix} c \\ c \end{smallmatrix}, \begin{smallmatrix} 0 \\ 1 \end{smallmatrix}) \in f(\varrho_-, \varrho_+) = f(\varrho_{+,-}, \varrho_{-,-}) \subseteq \varrho_- = (\begin{smallmatrix} 0 & 1 \\ 0 & 0 \end{smallmatrix} \begin{smallmatrix} 1 \\ 1 \end{smallmatrix})$ , i.e.,  
**order-reversing in  $x_2$** .

# Outline

$\pm$ -preclones ( $S$ -preclones)

$\pm$ -relations,  $\pm$ -preservation and a Galois connection

The lattice  ${}^{\pm}\mathcal{L}_A$  of Boolean  $\pm$ -preclones

## Some properties of the lattice ${}^{\pm}\mathcal{L}_A$

${}^{\pm}\mathcal{L}_A :=$  lattice of all  $\pm$ -preclones on  $A$  w.r.t.  $\subseteq$   
( ${}^{\pm}\mathcal{L}_2$  for Boolean  $\pm$ -preclones,  $A = \{0, 1\}$ )

Some properties (hold also for arbitrary monoids  $S$  instead of  $\pm$ )

- least  $\pm$ -preclone:  ${}^{\pm}J_A = \pm$ -projections  $= {}^{\pm}\langle \text{id}_A \rangle$
- largest  $\pm$ -preclone:  ${}^{\pm}\text{Op}(A)$
- ${}^{\pm}\mathcal{L}_A$  is atomic and coatomic (each  $\pm$ -preclone contains an atom and is contained in a coatom).
- ${}^{\pm}\mathcal{L}_A$  has finitely many atoms and coatoms.

${}^{\pm}\text{Op}(A)$  is finitely generated (by at most binary  $\pm$ -operations)

${}^{\pm}\text{Rel}(A)$  is finitely generated (by at most ternary  $\pm$ -relations)

Problem: Describe all maximal or minimal  $\pm$ -preclones (coatoms or atoms)

Recall:  $\mathcal{L}_2$ , the Post lattice of Boolean clones, is countable and has 5 maximal and 7 minimal clones.

## Some properties of the lattice ${}^{\pm}\mathcal{L}_A$

${}^{\pm}\mathcal{L}_A :=$  lattice of all  $\pm$ -preclones on  $A$  w.r.t.  $\subseteq$   
( ${}^{\pm}\mathcal{L}_2$  for Boolean  $\pm$ -preclones,  $A = \{0, 1\}$ )

Some properties (hold also for arbitrary monoids  $S$  instead of  $\pm$ )

- least  $\pm$ -preclone:  ${}^{\pm}J_A = \pm$ -projections =  ${}^{\pm}\langle \text{id}_A \rangle$
- largest  $\pm$ -preclone:  ${}^{\pm}\text{Op}(A)$
- ${}^{\pm}\mathcal{L}_A$  is atomic and coatomic (each  $\pm$ -preclone contains an atom and is contained in a coatom).
- ${}^{\pm}\mathcal{L}_A$  has finitely many atoms and coatoms.

${}^{\pm}\text{Op}(A)$  is finitely generated (by at most binary  $\pm$ -operations)

${}^{\pm}\text{Rel}(A)$  is finitely generated (by at most ternary  $\pm$ -relations)

Problem: Describe all maximal or minimal  $\pm$ -preclones (coatoms or atoms)

Recall:  $\mathcal{L}_2$ , the Post lattice of Boolean clones, is countable and has 5 maximal and 7 minimal clones.

## Some properties of the lattice ${}^{\pm}\mathcal{L}_A$

${}^{\pm}\mathcal{L}_A :=$  lattice of all  $\pm$ -preclones on  $A$  w.r.t.  $\subseteq$   
( ${}^{\pm}\mathcal{L}_2$  for Boolean  $\pm$ -preclones,  $A = \{0, 1\}$ )

Some properties (hold also for arbitrary monoids  $S$  instead of  $\pm$ )

- least  $\pm$ -preclone:  ${}^{\pm}J_A = \pm$ -projections =  ${}^{\pm}\langle \text{id}_A \rangle$
- largest  $\pm$ -preclone:  ${}^{\pm}\text{Op}(A)$
- ${}^{\pm}\mathcal{L}_A$  is atomic and coatomic (each  $\pm$ -preclone contains an atom and is contained in a coatom).
- ${}^{\pm}\mathcal{L}_A$  has finitely many atoms and coatoms.

${}^{\pm}\text{Op}(A)$  is finitely generated (by at most binary  $\pm$ -operations)

${}^{\pm}\text{Rel}(A)$  is finitely generated (by at most ternary  $\pm$ -relations)

Problem: Describe all maximal or minimal  $\pm$ -preclones (coatoms or atoms)

Recall:  $\mathcal{L}_2$ , the Post lattice of Boolean clones, is countable and has 5 maximal and 7 minimal clones.

## Some properties of the lattice ${}^{\pm}\mathcal{L}_A$

${}^{\pm}\mathcal{L}_A :=$  lattice of all  $\pm$ -preclones on  $A$  w.r.t.  $\subseteq$   
( ${}^{\pm}\mathcal{L}_2$  for Boolean  $\pm$ -preclones,  $A = \{0, 1\}$ )

Some properties (hold also for arbitrary monoids  $S$  instead of  $\pm$ )

- least  $\pm$ -preclone:  ${}^{\pm}J_A = \pm$ -projections =  ${}^{\pm}\langle \text{id}_A \rangle$
- largest  $\pm$ -preclone:  ${}^{\pm}\text{Op}(A)$
- ${}^{\pm}\mathcal{L}_A$  is atomic and coatomic (each  $\pm$ -preclone contains an atom and is contained in a coatom).
- ${}^{\pm}\mathcal{L}_A$  has finitely many atoms and coatoms.

${}^{\pm}\text{Op}(A)$  is finitely generated (by at most binary  $\pm$ -operations)

${}^{\pm}\text{Rel}(A)$  is finitely generated (by at most ternary  $\pm$ -relations)

**Problem:** Describe all maximal or minimal  $\pm$ -preclones (coatoms or atoms)

Recall:  $\mathcal{L}_2$ , the Post lattice of Boolean clones, is countable and has 5 maximal and 7 minimal clones.

## Some properties of the lattice ${}^{\pm}\mathcal{L}_A$

${}^{\pm}\mathcal{L}_A :=$  lattice of all  $\pm$ -preclones on  $A$  w.r.t.  $\subseteq$   
( ${}^{\pm}\mathcal{L}_2$  for Boolean  $\pm$ -preclones,  $A = \{0, 1\}$ )

Some properties (hold also for arbitrary monoids  $S$  instead of  $\pm$ )

- least  $\pm$ -preclone:  ${}^{\pm}J_A = \pm$ -projections =  ${}^{\pm}\langle \text{id}_A \rangle$
- largest  $\pm$ -preclone:  ${}^{\pm}\text{Op}(A)$
- ${}^{\pm}\mathcal{L}_A$  is atomic and coatomic (each  $\pm$ -preclone contains an atom and is contained in a coatom).
- ${}^{\pm}\mathcal{L}_A$  has finitely many atoms and coatoms.

${}^{\pm}\text{Op}(A)$  is finitely generated (by at most binary  $\pm$ -operations)

${}^{\pm}\text{Rel}(A)$  is finitely generated (by at most ternary  $\pm$ -relations)

Problem: Describe all maximal or minimal  $\pm$ -preclones (coatoms or atoms)

Recall:  $\mathcal{L}_2$ , the Post lattice of Boolean clones, is countable and has 5 maximal and 7 minimal clones.

## Some properties of the lattice ${}^{\pm}\mathcal{L}_A$

${}^{\pm}\mathcal{L}_A :=$  lattice of all  $\pm$ -preclones on  $A$  w.r.t.  $\subseteq$   
( ${}^{\pm}\mathcal{L}_2$  for Boolean  $\pm$ -preclones,  $A = \{0, 1\}$ )

Some properties (hold also for arbitrary monoids  $S$  instead of  $\pm$ )

- least  $\pm$ -preclone:  ${}^{\pm}J_A = \pm$ -projections =  ${}^{\pm}\langle \text{id}_A \rangle$
- largest  $\pm$ -preclone:  ${}^{\pm}\text{Op}(A)$
- ${}^{\pm}\mathcal{L}_A$  is atomic and coatomic (each  $\pm$ -preclone contains an atom and is contained in a coatom).
- ${}^{\pm}\mathcal{L}_A$  has finitely many atoms and coatoms.

${}^{\pm}\text{Op}(A)$  is finitely generated (by at most binary  $\pm$ -operations)

${}^{\pm}\text{Rel}(A)$  is finitely generated (by at most ternary  $\pm$ -relations)

**Problem:** Describe all maximal or minimal  $\pm$ -preclones (coatoms or atoms)

Recall:  $\mathcal{L}_2$ , the Post lattice of Boolean clones, is countable and has 5 maximal and 7 minimal clones.



## Some properties of the lattice ${}^{\pm}\mathcal{L}_A$

${}^{\pm}\mathcal{L}_A :=$  lattice of all  $\pm$ -preclones on  $A$  w.r.t.  $\subseteq$   
( ${}^{\pm}\mathcal{L}_2$  for Boolean  $\pm$ -preclones,  $A = \{0, 1\}$ )

Some properties (hold also for arbitrary monoids  $S$  instead of  $\pm$ )

- least  $\pm$ -preclone:  ${}^{\pm}J_A = \pm$ -projections =  ${}^{\pm}\langle \text{id}_A \rangle$
- largest  $\pm$ -preclone:  ${}^{\pm}\text{Op}(A)$
- ${}^{\pm}\mathcal{L}_A$  is atomic and coatomic (each  $\pm$ -preclone contains an atom and is contained in a coatom).
- ${}^{\pm}\mathcal{L}_A$  has finitely many atoms and coatoms.

${}^{\pm}\text{Op}(A)$  is finitely generated (by at most binary  $\pm$ -operations)

${}^{\pm}\text{Rel}(A)$  is finitely generated (by at most ternary  $\pm$ -relations)

**Problem:** Describe all maximal or minimal  $\pm$ -preclones (coatoms or atoms)

Recall:  $\mathcal{L}_2$ , the Post lattice of Boolean clones, is countable and has 5 maximal and 7 minimal clones.

## The maximal Boolean $\pm$ -preclones

### Theorem

*There are nine maximal Boolean  $\pm$ -preclones listed below. Each such preclone is of the form  $F = {}^{\pm}\text{Pol } \varrho$  for some  $\pm$ -relation  $\varrho = (\varrho_+, \varrho_-)$ :*

- (a)  ${}^{\pm}\text{Pol}(\sigma, \sigma)$  with  $\sigma \in \{\sigma_0, \sigma_1, \sigma_2, \sigma_3, \sigma_4\}$  where  $\text{Pol } \sigma_i$  is maximal in  $\mathcal{L}_2$  (0-preserving, 1-preserving, monotone, self-dual, linear operations)  
 $\sigma_0 = \{0\}, \sigma_1 = \{1\}, \sigma_2 = \leq = \{(0, 0), (0, 1), (1, 1)\},$   
 $\sigma_3 = \{(0, 1), (1, 0)\}, \sigma_4 = \{(x, y, z, u) \in A^4 \mid x + y + z + u = 0\}.$
- (b)  ${}^{\pm}\text{Pol}(\leq, \geq)$  our motivating example! all  $\pm$ -operations where each  $+$ -argument is order-preserving and each  $-$ -argument is order-reversing.
- (c)  ${}^{\pm}\text{Pol}(A, \emptyset) =$  all functions with positive or mixed signum.
- (d)  ${}^{\pm}\text{Pol}(A^2, \Delta_A) =$  all Boolean  $\pm$ -operations, where each negative argument is fictitious (including all negative constants).
- (e)  ${}^{\pm}\text{Pol}(\{0\}, \{1\})$  e.g., for  $f$  with  $\text{sgn}(f) = (+, +, -, -, -)$ :  
 $f \in {}^{\pm}\text{Pol}(\{0\}, \{1\}) \iff f(0, 0, 1, 1, 1) = 0 \text{ and } f(1, 1, 0, 0, 0) = 1.$

## The maximal Boolean $\pm$ -preclones

### Theorem

*There are nine maximal Boolean  $\pm$ -preclones listed below. Each such preclone is of the form  $F = {}^{\pm}\text{Pol } \varrho$  for some  $\pm$ -relation  $\varrho = (\varrho_+, \varrho_-)$ :*

- (a)  ${}^{\pm}\text{Pol}(\sigma, \sigma)$  with  $\sigma \in \{\sigma_0, \sigma_1, \sigma_2, \sigma_3, \sigma_4\}$  where  $\text{Pol } \sigma_i$  is maximal in  $\mathcal{L}_2$  (0-preserving, 1-preserving, monotone, self-dual, linear operations)  
 $\sigma_0 = \{0\}, \sigma_1 = \{1\}, \sigma_2 = \leq = \{(0, 0), (0, 1), (1, 1)\},$   
 $\sigma_3 = \{(0, 1), (1, 0)\}, \sigma_4 = \{(x, y, z, u) \in A^4 \mid x + y + z + u = 0\}.$
- (b)  ${}^{\pm}\text{Pol}(\leq, \geq)$  our motivating example! all  $\pm$ -operations where each  $+$ argument is order-preserving and each  $-$ argument is order-reversing.
- (c)  ${}^{\pm}\text{Pol}(A, \emptyset) =$  all functions with positive or mixed signum.
- (d)  ${}^{\pm}\text{Pol}(A^2, \Delta_A) =$  all Boolean  $\pm$ -operations, where each negative argument is fictitious (including all negative constants).
- (e)  ${}^{\pm}\text{Pol}(\{0\}, \{1\})$  e.g., for  $f$  with  $\text{sgn}(f) = (+, +, -, -, -)$ :  
 $f \in {}^{\pm}\text{Pol}(\{0\}, \{1\}) \iff f(0, 0, 1, 1, 1) = 0 \text{ and } f(1, 1, 0, 0, 0) = 1.$

## The maximal Boolean $\pm$ -preclones

### Theorem

*There are nine maximal Boolean  $\pm$ -preclones listed below. Each such preclone is of the form  $F = {}^{\pm}\text{Pol } \varrho$  for some  $\pm$ -relation  $\varrho = (\varrho_+, \varrho_-)$ :*

- (a)  ${}^{\pm}\text{Pol}(\sigma, \sigma)$  with  $\sigma \in \{\sigma_0, \sigma_1, \sigma_2, \sigma_3, \sigma_4\}$  where  $\text{Pol } \sigma_i$  is maximal in  $\mathcal{L}_2$  (0-preserving, 1-preserving, monotone, self-dual, linear operations)  
 $\sigma_0 = \{0\}, \sigma_1 = \{1\}, \sigma_2 = \leq = \{(0, 0), (0, 1), (1, 1)\},$   
 $\sigma_3 = \{(0, 1), (1, 0)\}, \sigma_4 = \{(x, y, z, u) \in A^4 \mid x + y + z + u = 0\}.$
- (b)  ${}^{\pm}\text{Pol}(\leq, \geq)$  our motivating example! all  $\pm$ -operations where each  $+$ -argument is order-preserving and each  $-$ -argument is order-reversing.
- (c)  ${}^{\pm}\text{Pol}(A, \emptyset) =$  all functions with positive or mixed signum.
- (d)  ${}^{\pm}\text{Pol}(A^2, \Delta_A) =$  all Boolean  $\pm$ -operations, where each negative argument is fictitious (including all negative constants).
- (e)  ${}^{\pm}\text{Pol}(\{0\}, \{1\})$  e.g., for  $f$  with  $\text{sgn}(f) = (+, +, -, -, -)$ :  
 $f \in {}^{\pm}\text{Pol}(\{0\}, \{1\}) \iff f(0, 0, 1, 1, 1) = 0 \text{ and } f(1, 1, 0, 0, 0) = 1.$

## The maximal Boolean $\pm$ -preclones

### Theorem

*There are nine maximal Boolean  $\pm$ -preclones listed below. Each such preclone is of the form  $F = {}^{\pm}\text{Pol } \varrho$  for some  $\pm$ -relation  $\varrho = (\varrho_+, \varrho_-)$ :*

- (a)  ${}^{\pm}\text{Pol}(\sigma, \sigma)$  with  $\sigma \in \{\sigma_0, \sigma_1, \sigma_2, \sigma_3, \sigma_4\}$  where  $\text{Pol } \sigma_i$  is maximal in  $\mathcal{L}_2$  (0-preserving, 1-preserving, monotone, self-dual, linear operations)  
 $\sigma_0 = \{0\}, \sigma_1 = \{1\}, \sigma_2 = \leq = \{(0, 0), (0, 1), (1, 1)\},$   
 $\sigma_3 = \{(0, 1), (1, 0)\}, \sigma_4 = \{(x, y, z, u) \in A^4 \mid x + y + z + u = 0\}.$
- (b)  ${}^{\pm}\text{Pol}(\leq, \geq)$  our motivating example! all  $\pm$ -operations where each  $+$ -argument is order-preserving and each  $-$ -argument is order-reversing.
- (c)  ${}^{\pm}\text{Pol}(A, \emptyset) =$  all functions with positive or mixed signum.
- (d)  ${}^{\pm}\text{Pol}(A^2, \Delta_A) =$  all Boolean  $\pm$ -operations, where each negative argument is fictitious (including all negative constants).
- (e)  ${}^{\pm}\text{Pol}(\{0\}, \{1\})$  e.g., for  $f$  with  $\text{sgn}(f) = (+, +, -, -, -)$ :  
 $f \in {}^{\pm}\text{Pol}(\{0\}, \{1\}) \iff f(0, 0, 1, 1, 1) = 0 \text{ and } f(1, 1, 0, 0, 0) = 1.$

## The maximal Boolean $\pm$ -preclones

### Theorem

*There are nine maximal Boolean  $\pm$ -preclones listed below. Each such preclone is of the form  $F = {}^{\pm}\text{Pol } \varrho$  for some  $\pm$ -relation  $\varrho = (\varrho_+, \varrho_-)$ :*

- (a)  ${}^{\pm}\text{Pol}(\sigma, \sigma)$  with  $\sigma \in \{\sigma_0, \sigma_1, \sigma_2, \sigma_3, \sigma_4\}$  where  $\text{Pol } \sigma_i$  is maximal in  $\mathcal{L}_2$  (0-preserving, 1-preserving, monotone, self-dual, linear operations)  
 $\sigma_0 = \{0\}, \sigma_1 = \{1\}, \sigma_2 = \leq = \{(0, 0), (0, 1), (1, 1)\},$   
 $\sigma_3 = \{(0, 1), (1, 0)\}, \sigma_4 = \{(x, y, z, u) \in A^4 \mid x + y + z + u = 0\}.$
- (b)  ${}^{\pm}\text{Pol}(\leq, \geq)$  our motivating example! all  $\pm$ -operations where each  $+$ -argument is order-preserving and each  $-$ -argument is order-reversing.
- (c)  ${}^{\pm}\text{Pol}(A, \emptyset) =$  all functions with positive or mixed signum.
- (d)  ${}^{\pm}\text{Pol}(A^2, \Delta_A) =$  all Boolean  $\pm$ -operations, where each negative argument is fictitious (including all negative constants).
- (e)  ${}^{\pm}\text{Pol}(\{0\}, \{1\})$  e.g., for  $f$  with  $\text{sgn}(f) = (+, +, -, -, -)$ :  
 $f \in {}^{\pm}\text{Pol}(\{0\}, \{1\}) \iff f(0, 0, 1, 1, 1) = 0 \text{ and } f(1, 1, 0, 0, 0) = 1.$

## The maximal Boolean $\pm$ -preclones

### Theorem

*There are nine maximal Boolean  $\pm$ -preclones listed below. Each such preclone is of the form  $F = {}^{\pm}\text{Pol } \varrho$  for some  $\pm$ -relation  $\varrho = (\varrho_+, \varrho_-)$ :*

- (a)  ${}^{\pm}\text{Pol}(\sigma, \sigma)$  with  $\sigma \in \{\sigma_0, \sigma_1, \sigma_2, \sigma_3, \sigma_4\}$  where  $\text{Pol } \sigma_i$  is maximal in  $\mathcal{L}_2$  (0-preserving, 1-preserving, monotone, self-dual, linear operations)  
 $\sigma_0 = \{0\}, \sigma_1 = \{1\}, \sigma_2 = \leq = \{(0, 0), (0, 1), (1, 1)\},$   
 $\sigma_3 = \{(0, 1), (1, 0)\}, \sigma_4 = \{(x, y, z, u) \in A^4 \mid x + y + z + u = 0\}.$
- (b)  ${}^{\pm}\text{Pol}(\leq, \geq)$  our motivating example! all  $\pm$ -operations where each  $+$ -argument is order-preserving and each  $-$ -argument is order-reversing.
- (c)  ${}^{\pm}\text{Pol}(A, \emptyset) =$  all functions with positive or mixed signum.
- (d)  ${}^{\pm}\text{Pol}(A^2, \Delta_A) =$  all Boolean  $\pm$ -operations, where each negative argument is fictitious (including all negative constants).
- (e)  ${}^{\pm}\text{Pol}(\{0\}, \{1\})$  e.g., for  $f$  with  $\text{sgn}(f) = (+, +, -, -, -)$ :  
 $f \in {}^{\pm}\text{Pol}(\{0\}, \{1\}) \iff f(0, 0, 1, 1, 1) = 0 \text{ and } f(1, 1, 0, 0, 0) = 1.$

## The minimal Boolean $\pm$ -preclones

### Theorem

*There are twenty minimal Boolean  $\pm$ -preclones. Each such  $\pm$ -preclone is of the form  $\pm\langle f \rangle$  with one  $\pm$ -operation  $f$  as generator:*

(A)  $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle x + y + z \rangle$   
*where the generators have signum  $\lambda = (+, +, +, -)$ ,  
 (majority and minority operation, the last argument is fictitious) (#2)*

(B)  $\pm\langle h_0 \rangle, \pm\langle h_1 \rangle, \pm\langle h_y \rangle$  where  $h_{\dagger}(x, y, z, u) = \begin{cases} x & \text{if } x = y \text{ or } z = u, \\ \dagger & \text{otherwise,} \end{cases}$   
*where the generators have signum  $\lambda = (+, +, -, -)$ , (#3)*

(C)  $\pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle, \pm\langle x \vee (y \wedge z) \rangle, \pm\langle x \wedge (y \vee z) \rangle,$   
 $\pm\langle x \vee (y \wedge \neg z) \rangle, \pm\langle x \wedge (y \vee \neg z) \rangle,$   
 $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle (x \wedge y) \vee (y \wedge \neg z) \vee (\neg z \wedge x) \rangle,$   
*where the generators have signum  $\lambda = (+, +, -)$ , (#8)*

(D)  $\pm\langle 0 \rangle, \pm\langle 1 \rangle, \pm\langle y \rangle, \pm\langle \neg y \rangle, \pm\langle \neg x \rangle, \pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle$   
*where the generators have signum  $\lambda = (+, -)$ . (#7)*



## The minimal Boolean $\pm$ -preclones

### Theorem

*There are twenty minimal Boolean  $\pm$ -preclones. Each such  $\pm$ -preclone is of the form  $\pm\langle f \rangle$  with one  $\pm$ -operation  $f$  as generator:*

(A)  $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle x + y + z \rangle$   
*where the generators have signum  $\lambda = (+, +, +, -)$ ,  
 (majority and minority operation, the last argument is fictitious) (#2)*

(B)  $\pm\langle h_0 \rangle, \pm\langle h_1 \rangle, \pm\langle h_y \rangle$  where  $h_{\dagger}(x, y, z, u) = \begin{cases} x & \text{if } x = y \text{ or } z = u, \\ \dagger & \text{otherwise,} \end{cases}$   
*where the generators have signum  $\lambda = (+, +, -, -)$ , (#3)*

(C)  $\pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle, \pm\langle x \vee (y \wedge z) \rangle, \pm\langle x \wedge (y \vee z) \rangle,$   
 $\pm\langle x \vee (y \wedge \neg z) \rangle, \pm\langle x \wedge (y \vee \neg z) \rangle,$   
 $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle (x \wedge y) \vee (y \wedge \neg z) \vee (\neg z \wedge x) \rangle,$   
*where the generators have signum  $\lambda = (+, +, -)$ , (#8)*

(D)  $\pm\langle 0 \rangle, \pm\langle 1 \rangle, \pm\langle y \rangle, \pm\langle \neg y \rangle, \pm\langle \neg x \rangle, \pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle$   
*where the generators have signum  $\lambda = (+, -)$ . (#7)*

## The minimal Boolean $\pm$ -preclones

### Theorem

*There are twenty minimal Boolean  $\pm$ -preclones. Each such  $\pm$ -preclone is of the form  $\pm\langle f \rangle$  with one  $\pm$ -operation  $f$  as generator:*

(A)  $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle x + y + z \rangle$   
*where the generators have signum  $\lambda = (+, +, +, -)$ ,  
 (majority and minority operation, the last argument is fictitious)* (#2)

(B)  $\pm\langle h_0 \rangle, \pm\langle h_1 \rangle, \pm\langle h_y \rangle$  where  $h_{\dagger}(x, y, z, u) = \begin{cases} x & \text{if } x = y \text{ or } z = u, \\ \dagger & \text{otherwise,} \end{cases}$   
*where the generators have signum  $\lambda = (+, +, -, -)$ ,* (#3)

(C)  $\pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle, \pm\langle x \vee (y \wedge z) \rangle, \pm\langle x \wedge (y \vee z) \rangle,$   
 $\pm\langle x \vee (y \wedge \neg z) \rangle, \pm\langle x \wedge (y \vee \neg z) \rangle,$   
 $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle (x \wedge y) \vee (y \wedge \neg z) \vee (\neg z \wedge x) \rangle,$   
*where the generators have signum  $\lambda = (+, +, -)$ ,* (#8)

(D)  $\pm\langle 0 \rangle, \pm\langle 1 \rangle, \pm\langle y \rangle, \pm\langle \neg y \rangle, \pm\langle \neg x \rangle, \pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle$   
*where the generators have signum  $\lambda = (+, -)$ .* (#7)

## The minimal Boolean $\pm$ -preclones

### Theorem

*There are twenty minimal Boolean  $\pm$ -preclones. Each such  $\pm$ -preclone is of the form  $\pm\langle f \rangle$  with one  $\pm$ -operation  $f$  as generator:*

(A)  $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle x + y + z \rangle$   
*where the generators have signum  $\lambda = (+, +, +, -)$ ,  
 (majority and minority operation, the last argument is fictitious)* (#2)

(B)  $\pm\langle h_0 \rangle, \pm\langle h_1 \rangle, \pm\langle h_y \rangle$  where  $h_{\dagger}(x, y, z, u) = \begin{cases} x & \text{if } x = y \text{ or } z = u, \\ \dagger & \text{otherwise,} \end{cases}$   
*where the generators have signum  $\lambda = (+, +, -, -)$ ,* (#3)

(C)  $\pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle, \pm\langle x \vee (y \wedge z) \rangle, \pm\langle x \wedge (y \vee z) \rangle,$   
 $\pm\langle x \vee (y \wedge \neg z) \rangle, \pm\langle x \wedge (y \vee \neg z) \rangle,$   
 $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle (x \wedge y) \vee (y \wedge \neg z) \vee (\neg z \wedge x) \rangle,$   
*where the generators have signum  $\lambda = (+, +, -)$ ,* (#8)

(D)  $\pm\langle 0 \rangle, \pm\langle 1 \rangle, \pm\langle y \rangle, \pm\langle \neg y \rangle, \pm\langle \neg x \rangle, \pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle$   
*where the generators have signum  $\lambda = (+, -)$ .* (#7)

## The minimal Boolean $\pm$ -preclones

### Theorem

*There are twenty minimal Boolean  $\pm$ -preclones. Each such  $\pm$ -preclone is of the form  $\pm\langle f \rangle$  with one  $\pm$ -operation  $f$  as generator:*

(A)  $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle x + y + z \rangle$   
*where the generators have signum  $\lambda = (+, +, +, -)$ ,  
 (majority and minority operation, the last argument is fictitious)* (#2)

(B)  $\pm\langle h_0 \rangle, \pm\langle h_1 \rangle, \pm\langle h_y \rangle$  where  $h_{\dagger}(x, y, z, u) = \begin{cases} x & \text{if } x = y \text{ or } z = u, \\ \dagger & \text{otherwise,} \end{cases}$   
*where the generators have signum  $\lambda = (+, +, -, -)$ ,* (#3)

(C)  $\pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle, \pm\langle x \vee (y \wedge z) \rangle, \pm\langle x \wedge (y \vee z) \rangle,$   
 $\pm\langle x \vee (y \wedge \neg z) \rangle, \pm\langle x \wedge (y \vee \neg z) \rangle,$   
 $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle (x \wedge y) \vee (y \wedge \neg z) \vee (\neg z \wedge x) \rangle,$   
*where the generators have signum  $\lambda = (+, +, -)$ ,* (#8)

(D)  $\pm\langle 0 \rangle, \pm\langle 1 \rangle, \pm\langle y \rangle, \pm\langle \neg y \rangle, \pm\langle \neg x \rangle, \pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle$   
*where the generators have signum  $\lambda = (+, -)$ .* (#7)

# Preimage classes

(a tool for investigations of the structure of  ${}^{\pm}\mathcal{L}_2$ )

Let  $F \in \mathcal{L}_2$  be a Boolean clone.

$$F^{\square} := \{P \in {}^{\pm}\mathcal{L}_2 \mid \langle \mathring{P} \rangle = F\} \text{ *preimage class of } F*$$

$\langle \mathring{P} \rangle$  = “underlying clone” forgetting all signa

Remark: relational characterization  $\langle \mathring{P} \rangle = \text{Pol}\{\sigma \in \text{Rel}(A) \mid (\sigma, \sigma) \in {}^{\pm}\text{Inv } P\}$

Structure: semi-interval with greatest element  $P_F := \{f \in {}^{\pm}\text{Op}(A) \mid \mathring{f} \in F\}$

The lattice  ${}^{\pm}\mathcal{L}_2$  of Boolean  $\pm$ -preclones is the (disjoint) union of all preimage classes of Boolean clones:

$${}^{\pm}\mathcal{L}_2 = \bigcup \{F^{\square} \mid F \in \mathcal{L}_2\}.$$

## Preimage classes

(a tool for investigations of the structure of  ${}^{\pm}\mathcal{L}_2$ )

Let  $F \in \mathcal{L}_2$  be a Boolean clone.

$$F^{\square} := \{P \in {}^{\pm}\mathcal{L}_2 \mid \langle \mathring{P} \rangle = F\} \text{ *preimage class of } F*$$

$\langle \mathring{P} \rangle$  = “underlying clone” forgetting all signa

Remark: relational characterization  $\langle \mathring{P} \rangle = \text{Pol}\{\sigma \in \text{Rel}(A) \mid (\sigma, \sigma) \in {}^{\pm}\text{Inv } P\}$

Structure: semi-interval with greatest element  $P_F := \{f \in {}^{\pm}\text{Op}(A) \mid \mathring{f} \in F\}$

The lattice  ${}^{\pm}\mathcal{L}_2$  of Boolean  $\pm$ -preclones is the (disjoint) union of all preimage classes of Boolean clones:

$${}^{\pm}\mathcal{L}_2 = \bigcup \{F^{\square} \mid F \in \mathcal{L}_2\}.$$

## Preimage classes

(a tool for investigations of the structure of  ${}^{\pm}\mathcal{L}_2$ )

Let  $F \in \mathcal{L}_2$  be a Boolean clone.

$$F^{\square} := \{P \in {}^{\pm}\mathcal{L}_2 \mid \langle \dot{P} \rangle = F\} \text{ *preimage class of } F*$$

$\langle \dot{P} \rangle$  = “underlying clone” forgetting all signa

Remark: relational characterization  $\langle \dot{P} \rangle = \text{Pol}\{\sigma \in \text{Rel}(A) \mid (\sigma, \sigma) \in {}^{\pm}\text{Inv } P\}$

Structure: semi-interval with greatest element  $P_F := \{f \in {}^{\pm}\text{Op}(A) \mid \dot{f} \in F\}$

The lattice  ${}^{\pm}\mathcal{L}_2$  of Boolean  $\pm$ -preclones is the (disjoint) union of all preimage classes of Boolean clones:

$${}^{\pm}\mathcal{L}_2 = \bigcup \{F^{\square} \mid F \in \mathcal{L}_2\}.$$

## Preimage classes

(a tool for investigations of the structure of  ${}^{\pm}\mathcal{L}_2$ )

Let  $F \in \mathcal{L}_2$  be a Boolean clone.

$$F^{\square} := \{P \in {}^{\pm}\mathcal{L}_2 \mid \langle \dot{P} \rangle = F\} \text{ *preimage class of } F*$$

$\langle \dot{P} \rangle$  = “underlying clone” forgetting all signa

Remark: relational characterization  $\langle \dot{P} \rangle = \text{Pol}\{\sigma \in \text{Rel}(A) \mid (\sigma, \sigma) \in {}^{\pm}\text{Inv } P\}$

Structure: semi-interval with greatest element  $P_F := \{f \in {}^{\pm}\text{Op}(A) \mid \dot{f} \in F\}$

The lattice  ${}^{\pm}\mathcal{L}_2$  of Boolean  $\pm$ -preclones is the (disjoint) union of all preimage classes of Boolean clones:

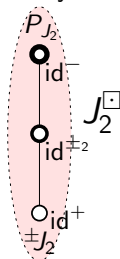
$${}^{\pm}\mathcal{L}_2 = \bigcup \{F^{\square} \mid F \in \mathcal{L}_2\}.$$



## How preimage classes look like? Some examples

unary clones  $F$  (i.e., generated by unary Boolean functions):

$$F = J_2 := \langle \text{id} \rangle: |J_2^{\square}| = 3$$



$$\text{id}^{-}(x) = x$$

$$\text{sgn}(\text{id}^{-}) = (-)$$

$$\text{id}^{\pm 2}(x, y) = y$$

$$\text{sgn}(\text{id}^{\pm 2}) = (+, -)$$

$$\text{id}^{+}(x) = x$$

$$\text{sgn}(\text{id}^{+}) = (+)$$

$$F_0 := \langle c_0 \rangle: |F_0^{\square}| = 6,$$

$$F_1 := \langle c_1 \rangle: |F_1^{\square}| = 6,$$

$$F_{01} := \langle c_0, c_1 \rangle: |F_{01}^{\square}| = 10,$$

$$F_{\neg} := \langle \neg \rangle: |F_{\neg}^{\square}| = 7,$$

$$F_2 := \langle \text{Op}^{(1)}(A) \rangle: |F_2^{\square}| = 19.$$

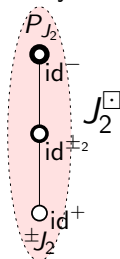
There are 12 join-irreducible elements generated by

$$\text{id}_A^{\pm 2}, \text{id}_A^{-}, c_0^{\pm}, c_0^{+}, c_0^{-}, c_1^{\pm}, c_1^{+}, c_1^{-}, \neg^{\pm 1}, \neg^{\pm 2}, \neg^{+}, \neg^{-}.$$

## How preimage classes look like? Some examples

unary clones  $F$  (i.e., generated by unary Boolean functions):

$$F = J_2 := \langle \text{id} \rangle: |J_2^\square| = 3$$



$$\text{id}^-(x) = x$$

$$\text{sgn}(\text{id}^-) = (-)$$

$$\text{id}^{\pm 2}(x, y) = y$$

$$\text{sgn}(\text{id}^{\pm 2}) = (+, -)$$

$$\text{id}^+(x) = x$$

$$\text{sgn}(\text{id}^+) = (+)$$

$$F_0 := \langle c_0 \rangle: |F_0^\square| = 6,$$

$$F_1 := \langle c_1 \rangle: |F_1^\square| = 6,$$

$$F_{01} := \langle c_0, c_1 \rangle: |F_{01}^\square| = 10,$$

$$F_{\neg} := \langle \neg \rangle: |F_{\neg}^\square| = 7,$$

$$F_2 := \langle \text{Op}^{(1)}(A) \rangle: |F_2^\square| = 19.$$

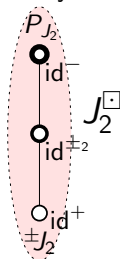
There are 12 join-irreducible elements generated by

$$\text{id}_A^{\pm 2}, \text{id}_A^-, c_0^\pm, c_0^+, c_0^-, c_1^\pm, c_1^+, c_1^-, \neg^{\pm 1}, \neg^{\pm 2}, \neg^+, \neg^-.$$

## How preimage classes look like? Some examples

unary clones  $F$  (i.e., generated by unary Boolean functions):

$$F = J_2 := \langle \text{id} \rangle: |J_2^\square| = 3$$



$$\text{id}^-(x) = x$$

$$\text{sgn}(\text{id}^-) = (-)$$

$$\text{id}^{\pm 2}(x, y) = y$$

$$\text{sgn}(\text{id}^{\pm 2}) = (+, -)$$

$$\text{id}^+(x) = x$$

$$\text{sgn}(\text{id}^+) = (+)$$

$$F_0 := \langle c_0 \rangle: |F_0^\square| = 6,$$

$$F_1 := \langle c_1 \rangle: |F_1^\square| = 6,$$

$$F_{01} := \langle c_0, c_1 \rangle: |F_{01}^\square| = 10,$$

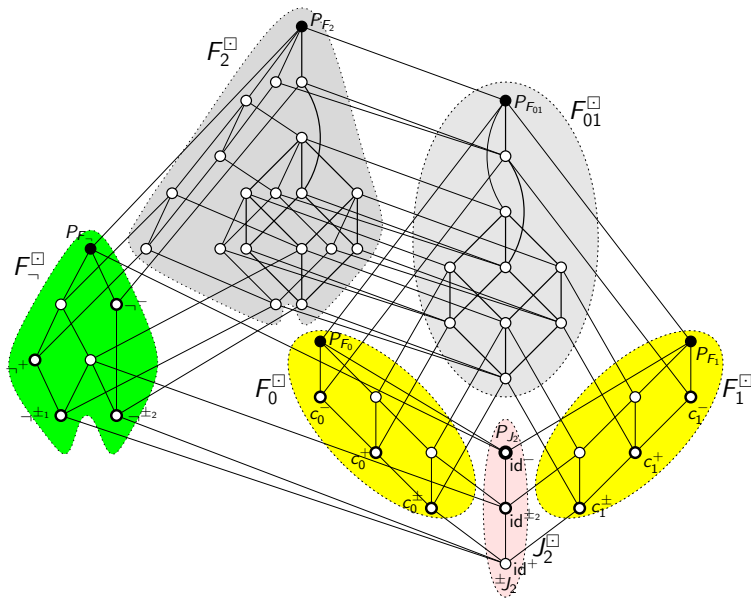
$$F_{\neg} := \langle \neg \rangle: |F_{\neg}^\square| = 7,$$

$$F_2 := \langle \text{Op}^{(1)}(A) \rangle: |F_2^\square| = 19.$$

There are 12 join-irreducible elements generated by

$$\text{id}_A^{\pm 2}, \text{id}_A^-, c_0^\pm, c_0^+, c_0^-, c_1^\pm, c_1^+, c_1^-, \neg^{\pm 1}, \neg^{\pm 2}, \neg^+, \neg^-.$$

## The lattice of unary $\pm$ -preclones



## A challenging open problem

up to now, we investigated few preimage classes  $F^{\square}$  (with  $F$  Boolean clone), all of them are finite

Does there exist a preimage class of infinite cardinality?

Does there exist a preimage class of uncountable cardinality?

Is the lattice  ${}^{\pm}\mathcal{L}_2$  of Boolean  $\pm$ -preclones countable?

(compare: the lattice  $\mathcal{L}_2$  of Boolean clones is countable  
[E.L. POST, 1921])

all what we know about Boolean  $\pm$ -preclones is contained in

P. JIPSEN, E. LEHTONEN, AND R. PÖSCHEL, *S-preclones and the Galois connection  ${}^S\text{Pol} - {}^S\text{Inv}$ , Part II: Boolean  $\pm$ -preclones*  
(in preparation, to be submitted)

## A challenging open problem

up to now, we investigated few preimage classes  $F^{\square}$  (with  $F$  Boolean clone), all of them are finite

Does there exist a preimage class of infinite cardinality?

Does there exist a preimage class of uncountable cardinality?

Is the lattice  ${}^{\pm}\mathcal{L}_2$  of Boolean  $\pm$ -preclones countable?

(compare: the lattice  $\mathcal{L}_2$  of Boolean clones is countable  
[E.L. POST, 1921])

all what we know about Boolean  $\pm$ -preclones is contained in

P. JIPSEN, E. LEHTONEN, AND R. PÖSCHEL, *S-preclones and the Galois connection  ${}^S\text{Pol} - {}^S\text{Inv}$ , Part II: Boolean  $\pm$ -preclones* (in preparation, to be submitted)

## A challenging open problem

up to now, we investigated few preimage classes  $F^{\square}$  (with  $F$  Boolean clone), all of them are finite

Does there exist a preimage class of infinite cardinality?

Does there exist a preimage class of uncountable cardinality?

Is the lattice  ${}^{\pm}\mathcal{L}_2$  of Boolean  $\pm$ -preclones countable?

(compare: the lattice  $\mathcal{L}_2$  of Boolean clones is countable  
[E.L. POST, 1921])

all what we know about Boolean  $\pm$ -preclones is contained in

P. JIPSEN, E. LEHTONEN, AND R. PÖSCHEL, *S-preclones and the Galois connection  ${}^S\text{Pol} - {}^S\text{Inv}$ , Part II: Boolean  $\pm$ -preclones* (in preparation, to be submitted)

## A challenging open problem

up to now, we investigated few preimage classes  $F^{\square}$  (with  $F$  Boolean clone), all of them are finite

Does there exist a preimage class of infinite cardinality?

Does there exist a preimage class of uncountable cardinality?

**Is the lattice  ${}^{\pm}\mathcal{L}_2$  of Boolean  $\pm$ -preclones countable?**

(compare: the lattice  $\mathcal{L}_2$  of Boolean clones is countable  
[E.L. POST, 1921])

all what we know about Boolean  $\pm$ -preclones is contained in

P. JIPSEN, E. LEHTONEN, AND R. PÖSCHEL,  *$S$ -preclones and the Galois connection  ${}^S\text{Pol} - {}^S\text{Inv}$ , Part II: Boolean  $\pm$ -preclones*  
(in preparation, to be submitted)



## A challenging open problem

up to now, we investigated few preimage classes  $F^{\square}$  (with  $F$  Boolean clone), all of them are finite

Does there exist a preimage class of infinite cardinality?

Does there exist a preimage class of uncountable cardinality?

**Is the lattice  ${}^{\pm}\mathcal{L}_2$  of Boolean  $\pm$ -preclones countable?**


(compare: the lattice  $\mathcal{L}_2$  of Boolean clones is countable  
[E.L. POST, 1921])

all what we know about Boolean  $\pm$ -preclones is contained in

P. JIPSEN, E. LEHTONEN, AND R. PÖSCHEL, *S-preclones and the Galois connection  ${}^S\text{Pol} - {}^S\text{Inv}$ , Part II: Boolean  $\pm$ -preclones* (in preparation, to be submitted)


# References

===== The classical Galois connection  $\text{Pol} - \text{Inv}$  =====

 V.G. BODNARČUK, L.A. KALUŽNIN, N.N. KOTOV, AND B.A. ROMOV, *Galois theory for Post algebras I*. Kibernetika (Kiev) (3), (1969), 1–10, (Russian).


 R. PÖSCHEL AND L.A. KALUŽNIN, *Funktionen- und Relationenalgebren*. Deutscher Verlag der Wissenschaften, Berlin, 1979, Birkhäuser Verlag Basel, Math. Reihe Bd. 67, 1979.

===== preclones (operads) =====


 Z. ÉSIK AND P. WEIL, *Algebraic recognizability of regular tree languages*. Theoret. Comput. Sci. 340(2), (2005), 291–321. (notion of preclone)

 E. LEHTONEN, *Characterization of preclones by matrix collections*. Asian-Eur. J. Math. 3(3), (2010), 457–473.

===== Analogy to multi-sorted algebras =====

 E. LEHTONEN, R. PÖSCHEL, AND T. WALDHAUSER, *Reflection-closed varieties of multisorted algebras and minor identities*. Algebra Universalis 79(3), (2018), Art. 70, 22 pages.

=====  $S$ -preclones (New) =====

 P. JIPSEN, E. LEHTONEN, AND R. PÖSCHEL,  *$S$ -preclones and the Galois connection  $S\text{Pol} - S\text{Inv}$ , Part I*, Algebra Universalis 85, 2024  
(arXiv 2023: <http://arxiv.org/abs/2306.00493>).

*Thank you for ATTENTION!*

