

# Boolean $\pm$ -preclones

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## In memoriam Günther Eigenthaler

This is the first AAA attended by me since the passing of Günther Eigenthaler. Therefore I dedicate this talk to the memory of my esteemed colleague and dear friend Günther.

GÜNTHER EIGENTHALER  
9.2.1950 - 14.2.2025



Günther, Laci Márki, R.P.

September 2023 during an excursion (Lunzer See) in Austria

# Outline

$\pm$ -preclones ( $S$ -preclones)

$\pm$ -relations,  $\pm$ -preservation and a Galois connection

The lattice  ${}^\pm\mathcal{L}_A$  of Boolean  $\pm$ -preclones

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## Motivating example

some history:

Nov. 2021 PALS talk by P. Jipsen on partially ordered algebras (and po-clones): operations which in each argument are *order-preserving* or *order-reversing* (for some given order on the base set).

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How to handle composition? order-reversing composed with order-reversing is order-preserving! Formalization: Collect the properties

in a monoid  $S = (\{+, -\}, \cdot)$ , here a group

.	+	-
+	+	-
-	-	+

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$S := \pm := \{+, -\}$  (in general, finite monoid  $S$ ).

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$f: A^n \rightarrow A$  with  $\text{sgn}(f) = (s_1, \dots, s_n) \in S^n$ ,

i.e., the  $i$ -th argument of  $f$  gets a label (*sign*)  $s_i \in S$  ( $i = 1, \dots, n$ ).

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Boolean  $\pm$ -preclone if  $A = \{0, 1\}$  [Remark: preclone = operad]

## (Boolean) $\pm$ -preclones

$\pm$ -preclone := set  $F \subseteq {}^\pm\text{Op}(A)$  of  $\pm$ -operations closed under:

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## Example (composition)

properties  $S = \{+, -\}$ ,  $A = \{0, 1\}$  with order  $0 < 1$ ,  
+ means order preserving, - means order reversing  
(such functions really form a  $\pm$ -preclone).

Composition:

$$f(x_1, x_2) = \neg x_1 \wedge x_2, \text{sgn}(f) = (s_1, s_2) = (-, +),$$
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# Outline

$\pm$ -preclones ( $S$ -preclones)

$\pm$ -relations,  $\pm$ -preservation and a Galois connection

The lattice  ${}^\pm\mathcal{L}_A$  of Boolean  $\pm$ -preclones

## $\pm$ -relations ( $S$ -relations) and $\pm$ -preservation $\triangleright^\pm$

recall:  $S = \pm := \{+, -\}$

$m$ -ary  $\pm$ -relation:  $\varrho = (\varrho_+, \varrho_-)$  with  $\varrho_s \subseteq A^m$  ( $s \in S$ )

classical notion of preservation:  $f \triangleright \varrho : \iff f(\varrho, \dots, \varrho) \subseteq \varrho$

The “ $S$ -version”:

$f \in {}^\pm \text{Op}^{(n)}(A)$ ,  $\text{sgn}(f) = (s_1, \dots, s_n)$ ,  $\varrho = (\varrho_+, \varrho_-) \in {}^\pm \text{Rel}^{(m)}(A)$

$f \triangleright^\pm (\varrho_+, \varrho_-) : \iff f(\varrho_{s_1}, \dots, \varrho_{s_n}) \subseteq \varrho_+$  and  $f(\varrho_{s_1 \cdots -}, \dots, \varrho_{s_n \cdots -}) \subseteq \varrho_-$

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$$\begin{array}{c}
 f( \begin{array}{|c|c|c|} \hline a_{11} & a_{12} & \dots & a_{1n} \\ \hline a_{21} & a_{22} & \dots & a_{2n} \\ \hline \end{array} ) = \textcolor{red}{\boxed{\bullet}} \\
 f( \begin{array}{|c|c|c|} \hline a_{m1} & a_{m2} & \dots & a_{mn} \\ \hline \end{array} ) = \textcolor{red}{\boxed{\bullet}}
 \end{array}$$

$\in \varrho_{s_1s}$     $\in \varrho_{s_2s}$     $\dots$     $\in \varrho_{s_ns}$   $\Rightarrow$   $\in \varrho_s$     $(s \in S)$

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## ±-relations ( $S$ -relations) and ±-preservation $\triangleright$

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### The “S-version”:

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## The Galois connection $\pm\text{Pol} - \pm\text{Inv}$

$\stackrel{\pm}{\triangleright}$  induces a Galois connection with the operators

$\pm\text{Pol } Q := \{f \in \pm\text{Op}(A) \mid \forall \varrho \in Q: f \stackrel{\pm}{\triangleright} \varrho\}$  ( $\pm$ -polymorphisms),

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for  $F \subseteq \pm\text{Op}(A)$  and  $Q \subseteq \pm\text{Rel}(A)$ .

Theorem (The Galois closures)

$\pm\langle F \rangle = \pm\text{Pol} \pm\text{Inv } F$  ( $\pm$ -preclone generated by  $F$ ),

$\pm[Q] = \pm\text{Inv} \pm\text{Pol } Q$  ( $\pm$ -relational clone generated by  $Q$ ).

[JipLP2023]: *S*-preclones and the Galois connection  ${}^S\text{Pol} - {}^S\text{Inv}$ , Part I,  
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## once more: our “motivating” Example

$(A, \leq)$  poset ( $0 < 1$ ),  $S = \{+, -\}$  (group).

For the  $\pm$ -preclone  $F$  (of  $\pm$ -operations where  $+$  means order preserving und  $-$  means order reversing) we have the following relational characterization:

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$g \stackrel{+}{\triangleright} \varrho$  implies

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Example: For  $g(x_1, x_2) = x_1 \vee \neg x_2$ ,  $\text{sgn}(g) = (+, -)$ ,

$g \stackrel{\pm}{\triangleright} \varrho$  implies

$g\left(\begin{smallmatrix} 0 \\ 1 \end{smallmatrix}, \begin{smallmatrix} c \\ c \end{smallmatrix}\right) \in f(\varrho_+, \varrho_-) \subseteq \varrho_+ = \left(\begin{smallmatrix} 0 & 0 & 1 \\ 0 & 1 & 1 \end{smallmatrix}\right)$ , i.e., order-preserving in  $x_1$ ,

$g\left(\begin{smallmatrix} c \\ c \end{smallmatrix}, \begin{smallmatrix} 0 \\ 1 \end{smallmatrix}\right) \in f(\varrho_-, \varrho_+) = f(\varrho_{+-}, \varrho_{--}) \subseteq \varrho_- = \left(\begin{smallmatrix} 0 & 1 & 1 \\ 0 & 0 & 1 \end{smallmatrix}\right)$ , i.e.,  
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## once more: our “motivating” Example

$(A, \leq)$  poset ( $0 < 1$ ),  $S = \{+, -\}$  (group).

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# Outline

$\pm$ -preclones ( $S$ -preclones)

$\pm$ -relations,  $\pm$ -preservation and a Galois connection

The lattice  ${}^\pm\mathcal{L}_A$  of Boolean  $\pm$ -preclones

## Some properties of the lattice ${}^\pm\mathcal{L}_A$

${}^\pm\mathcal{L}_A :=$  lattice of all  $\pm$ -preclones on  $A$  w.r.t.  $\subseteq$   
( ${}^\pm\mathcal{L}_2$  for Boolean  $\pm$ -preclones,  $A = \{0,1\}$ )

Some properties (hold also for arbitrary monoids  $S$  instead of  $\pm$ )

- least  $\pm$ -preclone:  ${}^\pm\mathcal{J}_A = \pm$ -projections =  ${}^\pm\langle \text{id}_A \rangle$
- largest  $\pm$ -preclone:  ${}^\pm\text{Op}(A)$
- ${}^\pm\mathcal{L}_A$  is atomic and coatomic (each  $\pm$ -preclone contains an atom and is contained in a coatom).
- ${}^\pm\mathcal{L}_A$  has finitely many atoms and coatoms.

${}^\pm\text{Op}(A)$  is finitely generated (by at most binary  $\pm$ -operations)

${}^\pm\text{Rel}(A)$  is finitely generated (by at most ternary  $\pm$ -relations)

Problem: Describe all maximal or minimal  $\pm$ -preclones (coatoms or atoms)

Recall:  $\mathcal{L}_2$ , the Post lattice of Boolean clones, is countable and has 5 maximal and 7 minimal clones.

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# The maximal Boolean $\pm$ -preclones

## Theorem

*There are nine maximal Boolean  $\pm$ -preclones listed below. Each such preclone is of the form  $F = {}^{\pm}\text{Pol } \varrho$  for some  $\pm$ -relation  $\varrho = (\varrho_+, \varrho_-)$ :*

- (a)  ${}^{\pm}\text{Pol}(\sigma, \sigma)$  with  $\sigma \in \{\sigma_0, \sigma_1, \sigma_2, \sigma_3, \sigma_4\}$  where  $\text{Pol } \sigma_i$  is maximal in  $\mathcal{L}_2$  (0-preserving, 1-preserving, monotone, self-dual, linear operations)  
 $\sigma_0 = \{0\}$ ,  $\sigma_1 = \{1\}$ ,  $\sigma_2 = \leq = \{(0, 0), (0, 1), (1, 1)\}$ ,  
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- (b)  ${}^{\pm}\text{Pol}(\leq, \geq)$  our motivating example! all  $\pm$ -operations where each +argument is order-preserving and each -argument is order-reversing.
- (c)  ${}^{\pm}\text{Pol}(A, \emptyset)$  = all functions with positive or mixed signum.
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(A)  $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle x + y + z \rangle$   
*where the generators have signum  $\lambda = (+, +, +, -)$ ,  
 (majority and minority operation, the last argument is fictitious)* (#2)

(B)  $\pm\langle h_0 \rangle, \pm\langle h_1 \rangle, \pm\langle h_y \rangle$  where  $h_\dagger(x, y, z, u) = \begin{cases} x & \text{if } x = y \text{ or } z = u, \\ \dagger & \text{otherwise,} \end{cases}$   
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(C)  $\pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle, \pm\langle x \vee (y \wedge z) \rangle, \pm\langle x \wedge (y \vee z) \rangle,$   
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There are twenty minimal Boolean  $\pm$ -preclones. Each such  $\pm$ -preclone is of the form  $\pm\langle f \rangle$  with one  $\pm$ -operation  $f$  as generator:

(A)  $\pm\langle (x \wedge y) \vee (y \wedge z) \vee (z \wedge x) \rangle, \pm\langle x + y + z \rangle$   
*where the generators have signum  $\lambda = (+, +, +, -)$ ,  
 (majority and minority operation, the last argument is fictitious)* (#2)

(B)  $\pm\langle h_0 \rangle, \pm\langle h_1 \rangle, \pm\langle h_y \rangle$  where  $h_{\dagger}(x, y, z, u) = \begin{cases} x & \text{if } x = y \text{ or } z = u, \\ \dagger & \text{otherwise,} \end{cases}$   
*where the generators have signum  $\lambda = (+, +, -, -)$*  (#3)

(C)  $\pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle, \pm\langle x \vee (y \wedge z) \rangle, \pm\langle x \wedge (y \vee z) \rangle,$   
 $\pm\langle x \vee (y \wedge \neg z) \rangle, \pm\langle x \wedge (y \vee \neg z) \rangle,$   
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*where the generators have signum  $\lambda = (+, +, -)$*  (#8)

(D)  $\pm\langle 0 \rangle, \pm\langle 1 \rangle, \pm\langle y \rangle, \pm\langle \neg y \rangle, \pm\langle \neg x \rangle, \pm\langle x \wedge y \rangle, \pm\langle x \vee y \rangle$   
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## Preimage classes

(a tool for investigations of the structure of  ${}^\pm \mathcal{L}_2$ )

Let  $F \in \mathcal{L}_2$  be a Boolean clone.

$F^\square := \{P \in {}^\pm \mathcal{L}_2 \mid \langle \mathring{P} \rangle = F\}$  preimage class of  $F$

$\langle \mathring{P} \rangle$  = “underlying clone” forgetting all signs

Remark: relational characterization  $\langle \mathring{P} \rangle = \text{Pol}\{\sigma \in \text{Rel}(A) \mid (\sigma, \sigma) \in {}^\pm \text{Inv } P\}$

Structure: semi-interval with greatest element  $P_F := \{f \in {}^\pm \text{Op}(A) \mid \mathring{f} \in F\}$

The lattice  ${}^\pm \mathcal{L}_2$  of Boolean  $\pm$ -preclones is the (disjoint) union of all preimage classes of Boolean clones:

$${}^\pm \mathcal{L}_2 = \bigcup \{F^\square \mid F \in \mathcal{L}_2\}.$$

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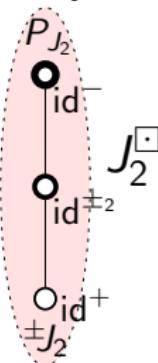
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## How preimage classes look like? Some examples

unary clones  $F$  (i.e., generated by unary Boolean functions):

$$F = J_2 := \langle \text{id} \rangle: |J_2^\square| = 3$$



$$F_0 := \langle c_0 \rangle: |F_0^\square| = 6,$$

$$F_1 := \langle c_1 \rangle: |F_1^\square| = 6,$$

$$F_{01} := \langle c_0, c_1 \rangle: |F_{01}^\square| = 10,$$

$$F_{\neg} := \langle \neg \rangle: |F_{\neg}^\square| = 7,$$

$$F_2 := \langle \text{Op}^{(1)}(A) \rangle: |F_2^\square| = 19.$$

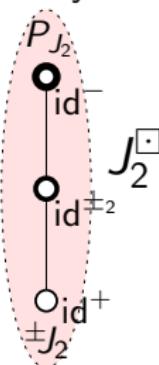
$$\begin{aligned} \text{id}^-(x) &= x \\ \text{sgn}(\text{id}^-) &= (-) \\ \text{id}^{\pm_2}(x, y) &= y \\ \text{sgn}(\text{id}^{\pm_2}) &= (+, -) \\ \text{id}^+(x) &= x \\ \text{sgn}(\text{id}^+) &= (+) \end{aligned}$$

There are 12 join-irreducible elements generated by  
 $\text{id}_A^{\pm_2}, \text{id}_A^-, c_0^+, c_0^-, c_1^+, c_1^-, \neg^{\pm_1}, \neg^{\pm_2}, \neg^+, \neg^-$ .

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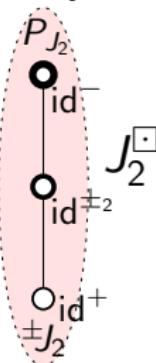
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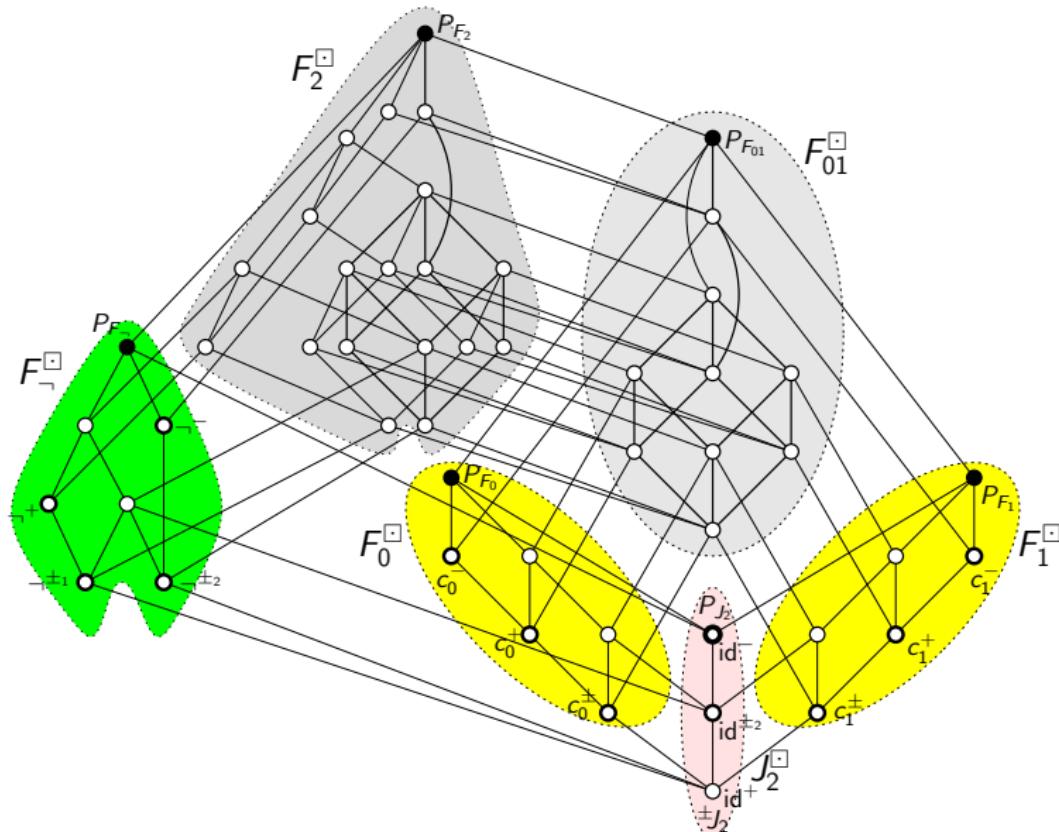
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## The lattice of unary $\pm$ -preclones



## A challenging open problem

up to now, we investigated few preimage classes  $F^\square$  (with  $F$  Boolean clone), all of them are finite

Does there exist a preimage class of infinite cardinality?

Does there exist a preimage class of uncountable cardinality?

Is the lattice  ${}^\pm\mathcal{L}_2$  of Boolean  $\pm$ -preclones countable?

(compare: the lattice  $\mathcal{L}_2$  of Boolean clones is countable  
[E.L. Post, 1921])

all what we know about Boolean  $\pm$ -preclones is contained in

P. JIPSEN, E. LEHTONEN, AND R. PÖSCHEL,  *$S$ -preclones and the Galois connection  ${}^S\text{Pol} - {}^S\text{Inv}$ , Part II: Boolean  $\pm$ -preclones*  
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## References

### ===== The classical Galois connection $\text{Pol} - \text{Inv}$ =====

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### ===== preclones (operads) =====

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### ===== Analogy to multi-sorted algebras =====

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### ===== $S$ -preclones (New) =====

- P. JIPSEN, E. LEHTONEN, AND R. PÖSCHEL,  *$S$ -preclones and the Galois connection  ${}^S\text{Pol} - {}^S\text{Inv}$ , Part I*, Algebra Universalis 85, 2024  
(arXiv 2023: <http://arxiv.org/abs/2306.00493>).

*Thank you for your  
ATTENTION!*

