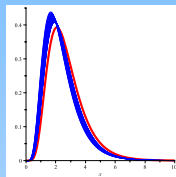
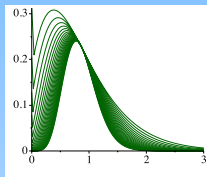


# PROBABILISTIC ANALYSIS OF THE (1 + 1)-EVOLUTIONARY ALGORITHM

Hsien-Kuei Hwang (joint with *Alois Panholzer*,  
*Nicolas Rolin*, *Tsung-Hsi Tsai*, *Wei-Mei Chen*)

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*Evolutionary Computation* Just Accepted MS.

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## Probabilistic Analysis of the $(1 + 1)$ -Evolutionary Algorithm<sup>1</sup>

**Hsien-Kuei Hwang**<sup>2</sup>

hkhwang@stat.sinica.edu.tw

Institute of Statistical Science & Institute of Information Science, Academia Sinica, Taipei 115, Taiwan

**Alois Panholzer**<sup>3</sup>

Alois.Panholzer@tuwien.ac.at

Institut für Diskrete Mathematik und Geometrie, Technische Universität Wien, Wiedner Hauptstraße 8-10/104, 1040 Wien, Austria

**Nicolas Rolin**

nicolasjrolin@gmail.com

LIPN, Institut Galilée, Université Paris 13, 93430, Villetaneuse, France

**Tsung-Hsi Tsai**

chonghi@stat.sinica.edu.tw

Institute of Statistical Science, Academia Sinica, Taipei 115, Taiwan

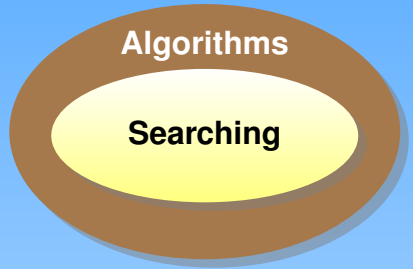
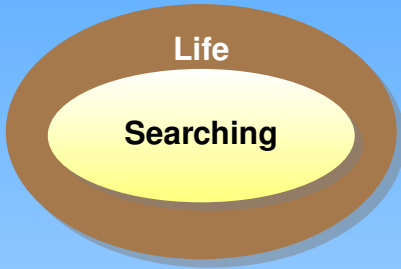
**Wei-Mei Chen**<sup>4</sup>

wmchen@mail.ntust.edu.tw

Department of Electronic and Computer Engineering, National Taiwan University of Science and Technology, Taipei 106, Taiwan



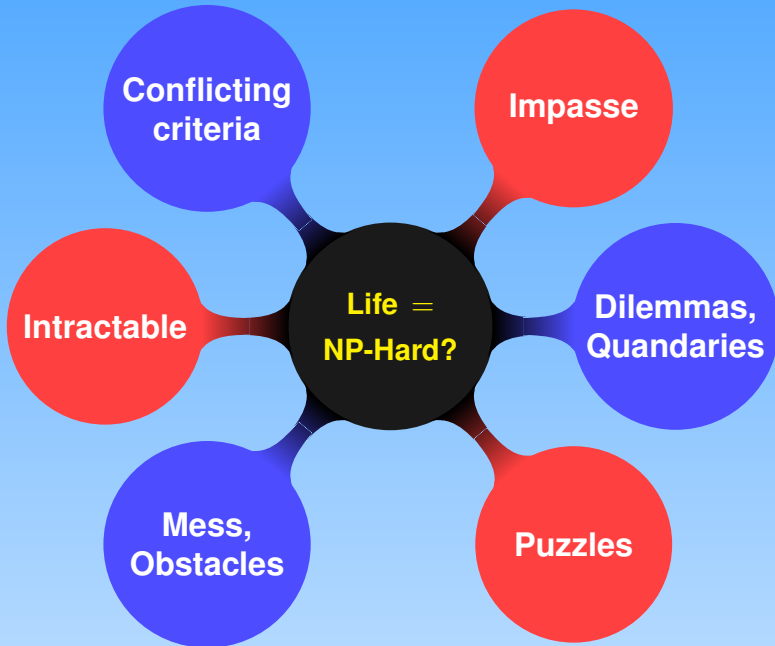
# WE ARE ALWAYS *SEARCHING* & *RESEARCHING*



**Life = Algorithms ?**

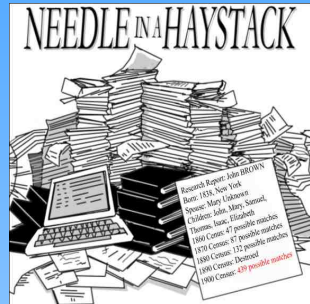


# OPTIMIZATION PROBLEMS EVERYWHERE



# SEARCH ALGORITHMS

- Backtracking
- Branch-and-bound
- Greedy
- Dynamic programming
- Simulated annealing
- Evolutionary algorithms
- Ant colony optimization
- Particle swarm
- Tabu search
- GRASP
- ... Meta-heuristics



## Evolutionary algorithm

From Wikipedia, the free encyclopedia

In **artificial intelligence**, an **evolutionary algorithm** (EA) is a **subset** of **evolutionary computation**, a generic population-based **metaheuristic optimization algorithm**. An EA uses mechanisms inspired by **biological evolution**, such as **reproduction**, **mutation**, **recombination**, and **selection**. **Candidate solutions** to the optimization problem play the role of individuals in a population, and the **fitness function** determines the quality of the solutions (see also **loss function**). **Evolution** of the population then takes place after the repeated application of the above operators. *Artificial evolution* (AE) describes a process involving individual *evolutionary algorithms*; EAs are individual components that participate in an AE.



# EVOLUTIONARY ALGORITHM

The use of Darwinian principles for automated problem solving originated in the 1950s.

*Darwin's theory of evolution: survival of the fittest*

- stochastic evolution on computers
  - ☞ cultivating problem solutions instead of calculating them
- randomized search heuristics
  - ☞ *generate* and *test* (or trial and error)
- useful for global optimization, if the problem is
  - too complex to be handled by an exact method or
  - no exact method is available

**Pioneers: John Holland, Lawrence J. Fogel, Ingo Rechenberg, ...**



# EVOLUTIONARY ALGORITHMS (EAs)

## Anne-Wil Harzing's s/w Publish or Perish

1	Cites	Authors	Title	Year
2	13687	K Deb	Multi-objective optimization using evolutionary algorithms	2001
3	6223	CAC Coello, GB Lamont, DA Van Veldh	Evolutionary algorithms for solving multi-objective problems	2007
4	6113	E Zitzler, L Thiele	Multiobjective evolutionary algorithms: a comparative case study and the str	1999
5	5977	E Zitzler, M Laumanns, L Thiele	SPEA2: Improving the strength Pareto evolutionary algorithm	2001
6	5849	T Back	Evolutionary algorithms in theory and practice: evolution strategies, evolution	1996
7	4562	E Zitzler, K Deb, L Thiele	Comparison of multiobjective evolutionary algorithms: Empirical results	2000
8	3529	K Deb, S Agrawal, A Pratap, T Meyariv	A fast elitist non-dominated sorting genetic algorithm for multi-objective opti	2000
9	3394	HPP Schwefel	Evolution and optimum seeking: the sixth generation	1993
10	3091	T Bäck, DB Fogel, Z Michalewicz	Handbook of evolutionary computation	1997
11	2688	CM Fonseca, PJ Fleming	An overview of evolutionary algorithms in multiobjective optimization	1995
12	2328	P Larrañaga, JA Lozano	Estimation of distribution algorithms: A new tool for evolutionary computatic	2001
13	2195	E Zitzler	Evolutionary algorithms for multiobjective optimization: Methods and applica	1999
14	2168	T Bäck, HP Schwefel	An overview of evolutionary algorithms for parameter optimization	1993
15	2105	DA Van Veldhuizen, GB Lamont	Multiobjective evolutionary algorithms: Analyzing the state-of-the-art	2000
16	1848	RC Eberhart, Y Shi	Comparison between genetic algorithms and particle swarm optimization	1998
17	1838	ÁE Eiben, R Hinterding, ...	Parameter control in evolutionary algorithms	1999
18	1711	E Zitzler, L Thiele	Multiobjective optimization using evolutionary algorithms—a comparative ca	1998
19	1685	Z Michalewicz, M Schoenauer	Evolutionary algorithms for constrained parameter optimization problems	1996
20	1623	CAC Coello	Theoretical and numerical constraint-handling techniques used with evolution	2002

**Most popular: Multiobjective optimization problems**





# ELITISM IN MULTI-OBJECTIVE EVOLUTIONARY ALGORITHM

nsga

About 43,500 results (0.03 sec)

## A fast and elitist multiobjective genetic algorithm: **NSGA-II**

[K Deb](#), [A Pratap](#), [S Agarwal](#)... - IEEE transactions on ..., 2002 - [ieeexplore.ieee.org](#)

Abstract: Multi-objective evolutionary algorithms (MOEAs) that use non-dominated sorting and sharing have been criticized mainly for:(1) their  $O(MN/\sup 3/)$  computational complexity (where  $M$  is the number of objectives and  $N$  is the population size);(2) their non-elitism

☆ [🔗](#) Cited by 22270 [Related articles](#) [All 45 versions](#)

## A fast elitist non-dominated sorting genetic algorithm for multi-objective optimization: **NSGA-II**

[K Deb](#), [S Agrawal](#), [A Pratap](#), [T Meyarivan](#) - International Conference on ..., 2000 - Springer

Abstract Multi-objective evolutionary algorithms which use non-dominated sorting and sharing have been mainly criticized for their (i)  $O(MN^3)$  computational complexity (where  $M$  is the number of objectives and  $N$  is the population size),(ii) non-elitism approach, and (iii)

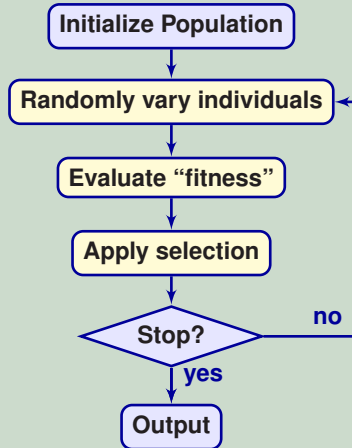
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*Our motivation: from maxima (skylines) to elites to EA*



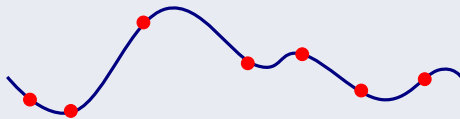
# COMPONENTS OF EAs

- Representation
  - Coding of solutions
- Initialization
- Parent selection
- Evaluation:
  - Fitness function
- Survivor selection
- Offspring Reproduction
  - Genetic operators
- Termination condition

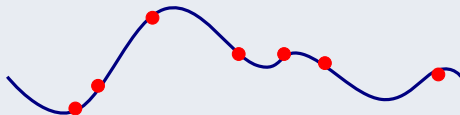


# TYPICAL PROGRESS OF AN EA

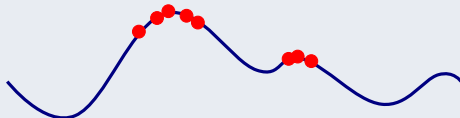
Initialization



Halfway



Termination



# PROS AND CONS OF EAs

## Disadvantages

- **Large convergence time**
- **Difficult adjustment of parameters**
- **Heuristic principle**
- **No guarantee of global max**

## Advantages

- **Reasonably good solutions quickly**
- **Suitable for complex search spaces**
- **Easy to parallelize**
- **Scalable to higher dimensional problems**



# DIFFICULTY OF ANALYSIS OF EA

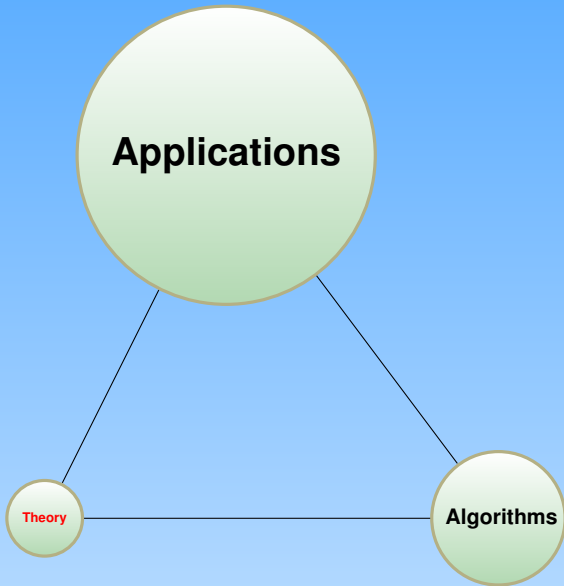
## A typical EA comprises several ingredients

- coding of solution
- population of individuals
- selection for reproduction
- operations for breeding new individuals
- fitness function to evaluate the new individual
- ...

***Mathematical description of the dynamics of the algorithms  
or the asymptotics of the complexity  $\implies$  challenging***



**Droste et al. (2002): *Theory is far behind experimental knowledge ... rigorous research is hard to find.***



# SIMPLEST VERSION: 1 PARENT, 1 CHILD, AND MUTATION ONLY

## **Algorithm (1 + 1)-EA**

- 1 **Choose an initial string  $x \in \{0, 1\}^n$  uniformly at random**
- 2 **Repeat until a terminating condition**
  - **(mutation)** Create  $y$  by flipping each bit of  $x$  independently with probability  $p$
  - **Replace  $x$  by  $y$  iff  $f(y) \geq f(x)$**

**$f$ : fitness (or objective) function**



# ANALYSIS OF (1 + 1)-EA UNDER ONEMAX

Known results for ONEMAX  $f(\mathbf{x}) = x_1 + \dots + x_n$

$X_n :=$  # steps used by the (1 + 1)-EA to reach the optimum state  $f(\mathbf{x}) = n$  when the mutation rate is  $\frac{1}{n}$

- Bäck (1992): transition probabilities
- Mühlenbein (1992):  $\mathbb{E}(X_n) \approx n \log n$
- Droste et al. (1998, 2002):  $\mathbb{E}(X_n) \asymp n \log n$

• • •

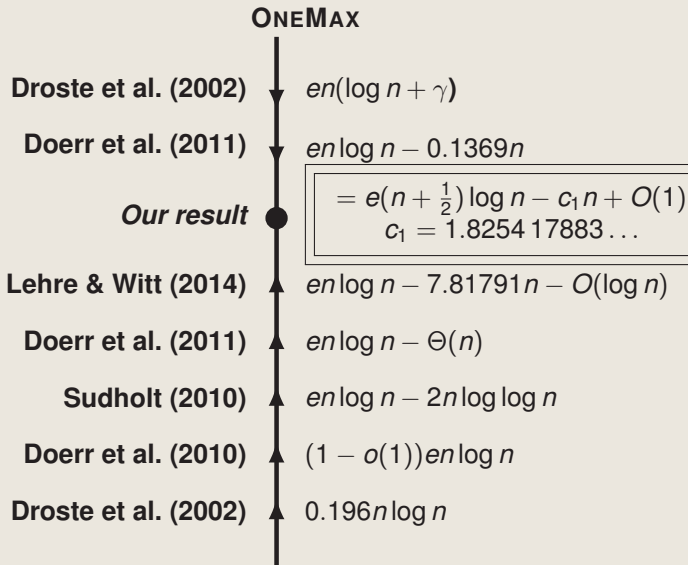
ONEMAX function		Linear functions $\sum w_i x_i$	
Doerr et al. (2010)	lower bound $(1 - o(1))en \log(n)$	Jagerskupper (2011)	upper bound $2.02en \log(n)$
Sudholt 2010	lower bound $en \log(n) - 2n \log \log(n)$	Doerr et al. (2010)	upper bound $1.39en \log(n)$
Doerr et al. (2011)	$en \log(n) - \Theta(n)$	Witt (2013)	upper bound $en \log(n) + O(n)$

**Approaches used: Markov chain, martingale, coupon collection, ...**





# KNOWN BOUNDS FOR $\mathbb{E}(X_n)$ UNDER ONEMAX



## “Rigorous hitting times for binary mutations”

Strongest results obtained so far but proof incomplete  
(probabilistic arguments)

- $\mathbb{E}(X_n) = en \log n + c_1 n + o(n)$ , **where**  $c_1 \approx -1.9$
- $\frac{X_n}{en} - \log n - c_1 \xrightarrow{d} \log \mathbf{Exp}(1)$  (**double-exponential**)

*their results had remained obscure in the EA-literature*



# OUR RESULTS

$$\mathbb{E}(X_n) = en \log n + c_1 n + \frac{1}{2}e \log n + c_2 + O\left(\frac{\log n}{n}\right)$$

$$c_1 = -e \left( \log 2 - \gamma - \phi_1\left(\frac{1}{2}\right) \right) \\ \approx -1.89254\ 17883\ 44686\ 82302\ 25714 \dots,$$

where  $\gamma$  is Euler's constant,

$$\phi_1(z) := \int_0^z \left( \frac{1}{S_1(t)} - \frac{1}{t} \right) dt,$$

$$S_1(z) := \sum_{\ell \geq 1} \frac{z^\ell}{\ell!} \sum_{0 \leq j < \ell} (\ell - j) \frac{(1 - z)^j}{j!}.$$

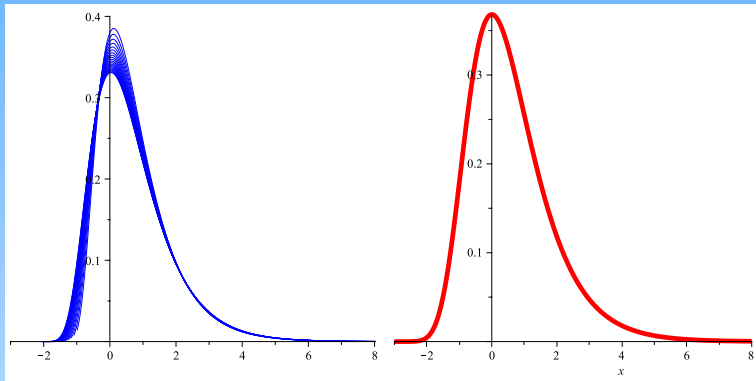
$$\textit{Indeed } \mathbb{E}(X_n) \sim n \sum_{k \geq 0} \frac{c'_k \log n + c_k}{n^k}$$



# LIMIT GUMBEL DISTRIBUTION

$$\mathbb{P} \left( \frac{X_n}{en} - \log n + \log 2 - \phi_1\left(\frac{1}{2}\right) \leq x \right) \rightarrow e^{-e^{-x}}$$

$$\phi_1\left(\frac{1}{2}\right) \approx -0.58029\ 56799\ 84283\ 81332\ 29240 \dots$$



left:  $n = 10..30$

right:  $e^{-x} - e^{-x}$

# OUR APPROACH

***RECURRENCE***  
 **$\Rightarrow$  *ASYMPTOTICS***



# THE RANDOM VARIABLE $X_{n,m}$

$$X_n := \sum_{0 \leq m \leq n} \binom{n}{m} p^m q^{n-m} X_{n,m}$$

$X_{n,m} := \#$  steps used by the  $(1 + 1)$ -EA to reach  $f(\mathbf{x}) = n$  when starting from  $f(\mathbf{x}) = n - m$

Let  $Q_{n,m}(t) := \mathbb{E}(t^{X_{n,m}})$ . Then  $Q_{n,0}(t) = 1$  and

$$Q_{n,m}(t) = \frac{t \sum_{1 \leq \ell \leq m} \lambda_{n,n-m,\ell} Q_{n,m-\ell}(t)}{1 - \left(1 - \sum_{1 \leq \ell \leq m} \lambda_{n,n-m,\ell}\right) t}$$

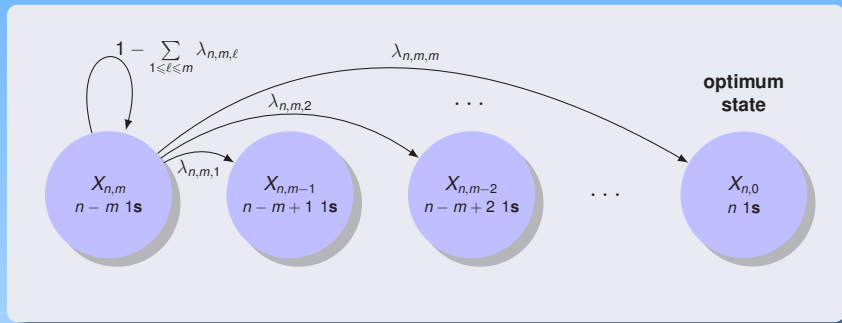
for  $1 \leq m \leq n$ , where  $(\mathbb{P}(m \text{ 1's} \mapsto m + \ell \text{ 1's}))$

$$\lambda_{n,m,\ell} := \left(1 - \frac{1}{n}\right)^n \sum_{0 \leq j \leq \min\{m, n-m-\ell\}} \binom{m}{j} \binom{n-m}{j+\ell} (n-1)^{-\ell-2j}$$



$$\lambda_{n,m,\ell} := \mathbb{P}(m \text{ 1's} \mapsto m + \ell \text{ 1's})$$

$$\sum_{0 \leq j \leq \min\{m, n-m-\ell\}} \underbrace{\binom{m}{j} \left(\frac{1}{n}\right)^j \left(1 - \frac{1}{n}\right)^{m-j}}_{1 \rightarrow 0} \underbrace{\binom{n-m}{j+\ell} \left(\frac{1}{n}\right)^{j+\ell} \left(1 - \frac{1}{n}\right)^{n-m-j-\ell}}_{0 \rightarrow 1}$$



Q: How to solve this recurrence?

$$Q_{n,m}(t) = \frac{t \sum_{1 \leq l \leq m} \lambda_{n,n-m,l} Q_{n,m-l}(t)}{1 - \left(1 - \sum_{1 \leq l \leq m} \lambda_{n,n-m,l}\right) t}$$

$$m = O(1)$$





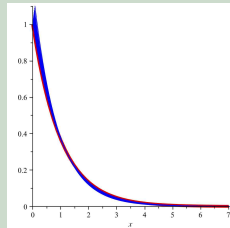
# SIMPLEST CASE: $m = 1$

$$m = 1 \quad (n - 1 \text{ 1s} \implies n \text{ 1s})$$

$$\lambda_{n,n-1,1} = \frac{1}{n} \left(1 - \frac{1}{n}\right)^{n-1} \implies \text{Geometric} \left( \frac{1}{n} \left(1 - \frac{1}{n}\right)^{n-1} \right)$$

$$Q_{n,1}(t) = \frac{\frac{1}{n} \left(1 - \frac{1}{n}\right)^{n-1} t}{1 - \left(1 - \frac{1}{n} \left(1 - \frac{1}{n}\right)^{n-1}\right) t}$$

$$\implies \frac{X_{n,1}}{en} \xrightarrow{d} \text{Exp}(1)$$



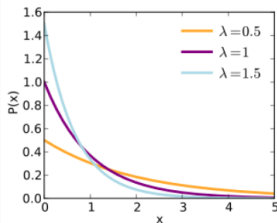
$$\mathbb{P} \left( \frac{X_{n,1}}{en} \leq x \right) \rightarrow 1 - e^{-x}$$

# EXPONENTIAL DISTRIBUTION

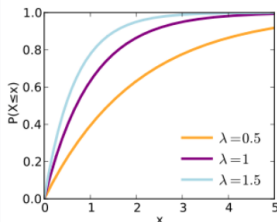
<b>Parameters</b>	$\lambda > 0$ rate, or inverse <b>scale</b>
<b>Support</b>	$x \in [0, \infty)$
<b>pdf</b>	$\lambda e^{-\lambda x}$
<b>CDF</b>	$1 - e^{-\lambda x}$
<b>Mean</b>	$\lambda^{-1}$
<b>Median</b>	$\lambda^{-1} \ln 2$
<b>Mode</b>	0
<b>Variance</b>	$\lambda^{-2}$
<b>Skewness</b>	2
<b>Ex. kurtosis</b>	6
<b>Entropy</b>	$1 - \ln(\lambda)$
<b>MGF</b>	$\left(1 - \frac{t}{\lambda}\right)^{-1}$ for $t < \lambda$
<b>CF</b>	$\left(1 - \frac{it}{\lambda}\right)^{-1}$

## Exponential

Probability density function



Cumulative distribution function



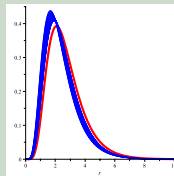
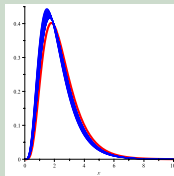
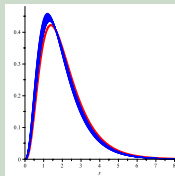
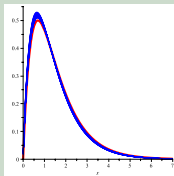
# BY INDUCTION $m = O(1)$

$$\frac{X_{n,m}}{en} \xrightarrow{d} \text{Exp}(1) + \cdots + \text{Exp}(m)$$

$$\mathbb{P}\left(\frac{X_{n,m}}{en} \leq x\right) \rightarrow (1 - e^{-x})^m$$

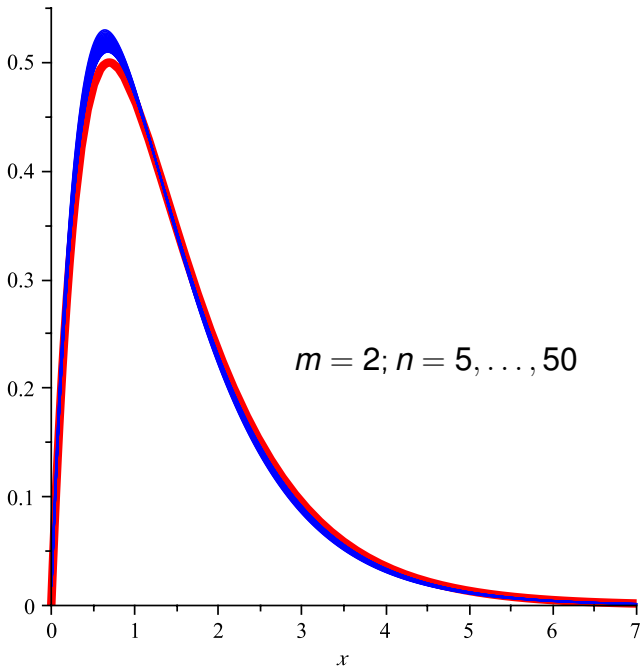
$$\mathbb{E}(X_{n,m}) \sim eH_m n \quad \text{and} \quad \mathbb{V}(X_{n,m}) \sim e^2 H_m^{(2)} n^2$$

$m = 2, 4, 6, 8$  &  $n = 5, \dots, 50$



**An LLT also holds**





# SKETCH OF PROOF

$$\begin{aligned}\lambda_{n,n-m,\ell} &= \left(1 - \frac{1}{n}\right)^n \sum_{0 \leq j \leq \min\{n-m, m-\ell\}} \binom{m}{j+\ell} \binom{n-m}{j} (n-1)^{-\ell-2j} \\ &= \binom{m}{\ell} e^{-1} n^{-\ell} \left(1 + O\left(\frac{m-\ell}{n(\ell+1)} + \frac{\ell}{n}\right)\right)\end{aligned}$$

$m = O(1) \implies j = 1$  *is dominant*



# SKETCH OF PROOF

$$\begin{aligned}\lambda_{n,n-m,\ell} &= \left(1 - \frac{1}{n}\right)^n \sum_{0 \leq j \leq \min\{n-m, m-\ell\}} \binom{m}{j+\ell} \binom{n-m}{j} (n-1)^{-\ell-2j} \\ &= \binom{m}{\ell} e^{-1} n^{-\ell} \left(1 + O\left(\frac{m-\ell}{n(\ell+1)} + \frac{\ell}{n}\right)\right)\end{aligned}$$

$m = O(1) \implies j = 1$  is dominant

$$\begin{aligned}Q_{n,m}(t) &= \frac{t \sum_{1 \leq \ell \leq m} \lambda_{n,n-m,\ell} Q_{n,m-\ell}(t)}{1 - \left(1 - \sum_{1 \leq \ell \leq m} \lambda_{n,n-m,\ell}\right) t} \\ \implies Q_{n,m}(t) &\sim \frac{\frac{m}{en} t}{1 - \left(1 - \frac{m}{en}\right) t} Q_{n,m-1}(t)\end{aligned}$$



# SKETCH OF PROOF: BY INDUCTION

$$Q_{n,m}(t) \sim \prod_{1 \leq r \leq m} \frac{\frac{r}{en} t}{1 - \left(1 - \frac{r}{en}\right) t}$$



# SKETCH OF PROOF: BY INDUCTION

$$Q_{n,m}(t) \sim \prod_{1 \leq r \leq m} \frac{\frac{r}{en} t}{1 - \left(1 - \frac{r}{en}\right) t}$$

$$\implies Q_{n,m}(e^{s/(en)}) \sim \prod_{1 \leq r \leq m} \frac{1}{1 - \frac{s}{r}} \implies \sum_{1 \leq r \leq m} \mathbf{Exp}(r)$$

**Fails when  $m \rightarrow \infty$**





# SKETCH OF PROOF: BY INDUCTION

$$Q_{n,m}(t) \sim \prod_{1 \leq r \leq m} \frac{\frac{r}{en} t}{1 - \left(1 - \frac{r}{en}\right) t}$$
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**Fails when  $m \rightarrow \infty$**

Let  $Y_m := \sum_{1 \leq r \leq m} \mathbf{Exp}(r)$ . Then as  $m \rightarrow \infty$

$$\mathbb{E}(e^{(Y_m - H_m)i\theta}) = \prod_{1 \leq r \leq m} \frac{e^{-\frac{i\theta}{r}}}{1 - \frac{i\theta}{r}} \rightarrow \prod_{r \geq 1} \frac{e^{-\frac{i\theta}{r}}}{1 - \frac{i\theta}{r}} = e^{-\gamma i\theta} \Gamma(1 - i\theta)$$



# SKETCH OF PROOF: BY INDUCTION

$$Q_{n,m}(t) \sim \prod_{1 \leq r \leq m} \frac{\frac{r}{en} t}{1 - \left(1 - \frac{r}{en}\right) t}$$
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$$\implies \mathbb{P}(Y_m - \log m \leq x) \rightarrow e^{-e^{-x}} \quad (x \in \mathbb{R})$$



$$m \rightarrow \infty$$



# EXPECTED VALUES $\mathbb{E}(X_{n,m})$

$$\mu_{n,m} := \mathbb{E}(X_{n,m}) = Q'_{n,m}(1) \quad (\mu_{n,0} = 0)$$

$$\mu_{n,m} = \frac{1 + \sum_{1 \leq \ell \leq m} \lambda_{n,n-m,\ell} \mu_{n,m-\ell}}{\sum_{1 \leq \ell \leq m} \lambda_{n,n-m,\ell} \mu_{n,m-\ell}} \quad (1 \leq m \leq n)$$

$$\implies \mathbb{E}(X_n) = 2^{-n} \sum_{0 \leq m \leq n} \binom{n}{m} \mu_{n,m}$$

$$\text{Let } e_n := \left(1 - \frac{1}{n+1}\right)^{n+1} \text{ and } \mu_{n,m}^* := \frac{e_n}{n} \mu_{n+1,m}$$

$$\sum_{1 \leq \ell \leq m} \lambda_{n,m,\ell}^* (\mu_{n,m}^* - \mu_{n,m-\ell}^*) = \frac{1}{n}$$

$$\lambda_{n,m,\ell}^* := \frac{\lambda_{n+1,n+1-m,\ell}}{e_n} = \sum_{0 \leq j \leq \min\{n+1-m, m-\ell\}} \binom{n+1-m}{j} \binom{m}{j+\ell} n^{-\ell-2j}$$



$$\mu_{n,1}^* = 1 \text{ \& } \sum_{1 \leq l \leq m} \lambda_{n,m,l}^* (\mu_{n,m}^* - \mu_{n,m-l}^*) = \frac{1}{n}$$

$$\mu_{n,2}^* = \frac{3n^2 + n - 1}{2n^2 + 2n - 1}$$

$$\mu_{n,3}^* = \frac{22n^6 + 40n^5 - 19n^4 - 42n^3 + 14n^2 + 15n - 6}{(2n^2 + 2n - 1)(6n^4 + 12n^3 - 7n^2 - 9n + 6)}$$

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$$\mu_{n,4}^* = \frac{\begin{pmatrix} 600n^{12} + 2616n^{11} + 1128n^{10} - 7460n^9 \\ -4958n^8 + 11506n^7 + 6167n^6 - 10887n^5 \\ -2862n^4 + 5917n^3 - 153n^2 - 1398n + 360 \end{pmatrix}}{(2n^2 + 2n - 1)(6n^4 + 12n^3 - 7n^2 - 9n + 6) \times (24n^6 + 72n^5 - 48n^4 - 140n^3 + 93n^2 + 83n - 60)}$$

$$\mu_{n,5}^* = \frac{\begin{pmatrix} 78912n^{20} + 626112n^{19} + 1150848n^{18} - 2455104n^{17} \\ -8313432n^{16} + 4491096n^{15} + 27182504n^{14} - 5263508n^{13} \\ -55021022n^{12} + 7628986n^{11} + 74466297n^{10} - 15193087n^9 \\ -67391443n^8 + 21902962n^7 + 38443857n^6 - 18491957n^5 \\ -11698973n^4 + 8358804n^3 + 827844n^2 - 1576800n + 302400 \end{pmatrix}}{(2n^2 + 2n - 1)(6n^4 + 12n^3 - 7n^2 - 9n + 6) \times (24n^6 + 72n^5 - 48n^4 - 140n^3 + 93n^2 + 83n - 60) \times (120n^8 + 480n^7 - 360n^6 - 1720n^5 + 1145n^4 + 2394n^3 - 1685n^2 - 1118n + 840)}$$

# ASYMPTOTICS OF $\mu_{n,m}^*$

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$$\mu_{n,3}^* = \frac{11}{6} - \frac{13}{6} n^{-1} + \frac{155}{36} n^{-2} - \frac{323}{36} n^{-3} + \frac{4007}{216} n^{-4} + \dots$$

$$\mu_{n,4}^* = \frac{25}{12} - \frac{41}{12} n^{-1} + \frac{329}{36} n^{-2} - \frac{917}{36} n^{-3} + \frac{61841}{864} n^{-4} + \dots$$

$$\mu_{n,5}^* = \frac{137}{60} - \frac{283}{60} n^{-1} + \frac{2839}{180} n^{-2} - \frac{19859}{360} n^{-3} + \frac{848761}{4320} n^{-4} + \dots$$

$$\mu_{n,6}^* = \frac{49}{20} - \frac{121}{20} n^{-1} + \frac{1453}{60} n^{-2} - \frac{36709}{360} n^{-3} + \frac{70451}{160} n^{-4} + \dots$$



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$$\left\{ H_m = \sum_{1 \leq j \leq m} \frac{1}{j} \right\} = \left\{ 1, \frac{3}{2}, \frac{11}{6}, \frac{25}{12}, \frac{137}{60}, \frac{49}{20}, \dots \right\}$$





# HEURISTICS

An Ansatz approximation:  $\mu_{n,m}^* \approx \sum_{k \geq 0} \frac{d_k(m)}{n^k}$

$$d_0(m) = H_m \quad (m \geq 0)$$

$$d_1(m) = H_m + \frac{1}{2} - \frac{3}{2} m \quad (m \geq 1)$$

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$$d_3(m) = \frac{1}{2} H_m + \frac{7}{24} - \frac{575}{432} m + \frac{23}{18} m^2 - \frac{283}{432} m^3 \quad (m \geq 2)$$

$$d_4(m) = \frac{5}{18} H_m - \frac{59}{720} - \frac{3439}{3456} m + \frac{15101}{11520} m^2 - \frac{19951}{17280} m^3 + \frac{5759}{11520} m^4 \quad (m \geq 4)$$

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**Complication:**  $d_k(m)$  holds for  $m \geq 2 \lfloor \frac{k}{2} \rfloor$



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**General pattern:**  $\mu_{n,m}^* \approx \sum_{k \geq 0} n^{-k} \left( b_k H_m + \sum_{0 \leq j \leq k} \varpi_{k,j} m^j \right)$



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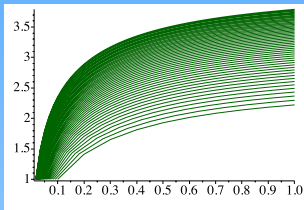
**General pattern:**  $\mu_{n,m}^* \approx \sum_{k \geq 0} n^{-k} \left( b_k H_m + \sum_{0 \leq j \leq k} \varpi_{k,j} m^j \right)$

$$\alpha := \frac{m}{n} \implies \boxed{\mu_{n,m}^* \approx H_m + \phi(\alpha) \quad \text{for } 1 \leq m \leq n}$$

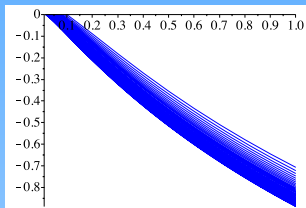


# A MORE GENERAL ANSATZ

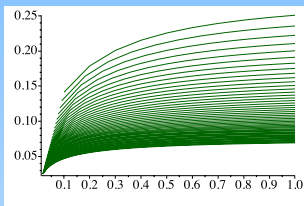
$$\mu_{n,m}^* \approx H_m + \phi_1(\alpha) + \frac{b_1 H_m + \phi_2(\alpha)}{n} + \frac{b_2 H_m + \phi_3(\alpha)}{n^2} + \dots$$



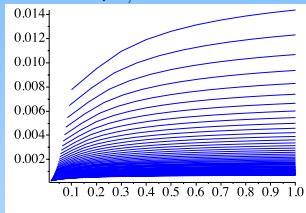
$\mu_{n,m}^*$



$\mu_{n,m}^* - H_m$



$\mu_{n,m}^* - (H_m + \phi_1(\alpha))$



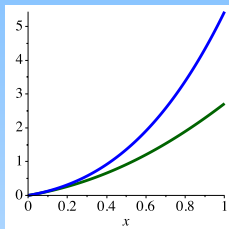
$\mu_{n,m}^* - (H_m + \phi_1(\alpha) + \frac{H_m + \phi_2(\alpha)}{n})$

$$S_r(z) := \sum_{\ell \geq 1} \frac{z^\ell}{\ell!} \sum_{0 \leq j < \ell} (\ell - j)^r \frac{(1 - z)^j}{j!}$$

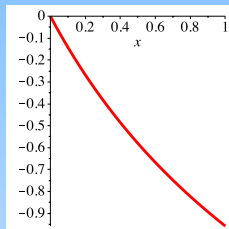
$$\phi_1(z) := \int_0^z \left( \frac{1}{S_1(t)} - \frac{1}{t} \right) dt$$

$$\phi_2(z) = \frac{1}{2} - \int_0^z \left( \frac{S_2(t)S_1'(t)}{2S_1(t)^3} - \frac{S_0(t)}{S_1(t)^2} - \frac{1}{2S_1(t)} - \frac{1}{2t^2} + \frac{1}{t} \right) dt$$

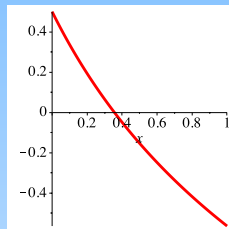
(analytic in  $|z| \leq 1$ )



$S_1(x)$  &  $S_2(x)$



$\phi_1(x)$



$\phi_2(x)$



# A GENERATING FUNCTION APPROACH?

$$\lambda_{n,m,\ell}^* = \sum_{0 \leq j \leq \min\{n+1-m, m-\ell\}} \binom{n+1-m}{j} \binom{m}{j+\ell} n^{-\ell-2j}$$

$$f_n(z) := \sum_{m \geq 1} \mu_{n,m}^* z^m$$

$$\begin{aligned} \sum_{0 \leq \ell < m} \lambda_{n,m,m-\ell}^* (\mu_{n,m}^* - \mu_{n,\ell}^*) &= \frac{1}{n} \\ \Rightarrow \frac{1}{2\pi i} \oint \left( \frac{1}{1-t} - \frac{z(t + \frac{1}{n})}{t(1 + \frac{t}{n} - z(t + \frac{1}{n}))} \right) \\ &\quad \times \left(1 + \frac{t}{n}\right)^{n+1} f_n \left( \frac{z(t + \frac{1}{n})}{t(1 + \frac{t}{n})} \right) dt = \frac{z}{n(1-z)} \end{aligned}$$



# A HEURISTIC

$$\frac{1}{2\pi i} \oint f_n(w) \Phi_n(z, w) dw = \frac{z}{n(1-z)}$$

$$\begin{aligned} \Phi_n(z, w) &:= \left(1 + \frac{\tau}{n}\right)^{n+1} \left( \frac{1}{1-\tau} - \frac{z(\tau + \frac{1}{n})}{\tau(1 + \frac{\tau}{n} - z(\tau + \frac{1}{n}))} \right) \frac{d\tau}{dw} \\ &\sim \frac{z(w-1)}{n(w-z)^2} + \dots \end{aligned}$$

Assume  $f_n(w) \sim \phi(w)$ .

$$\frac{z}{2\pi i n} \oint \frac{\phi(w)(w-1)}{(w-z)^2} dw = \frac{z}{n} (\phi(z) - (1-z)\phi'(z)) = \frac{z}{n} \frac{1}{1-z} = \mathbf{RHS}$$

Then ( $\phi(0) = 0$ )

$$\phi(z) - (1-z)\phi'(z) = \frac{1}{1-z} \implies \phi(z) = \frac{1}{1-z} \log \frac{1}{1-z}.$$

$$\implies \mu_{n,m}^* \sim H_m.$$





# HOW TO GUESS $\phi_1(\alpha)$ ?

Assume  $\mu_{n,m}^* \sim H_m + \phi(\alpha)$  ( $\alpha := \frac{m}{n}$ )

$$H_m - H_{m-\ell} = \frac{\ell}{m} + \frac{\ell(\ell-1)}{2m^2} + \dots$$

$$\phi\left(\frac{m}{n}\right) - \phi\left(\frac{m-\ell}{n}\right) = \phi'(\alpha) \frac{\ell}{m} + O\left(\frac{\ell^2}{m^2}\right)$$

## Matched asymptotics

$$\begin{aligned} \frac{1}{n} &= \sum_{1 \leq \ell \leq m} \lambda_{n,m,\ell}^* (\mu_{n,m}^* - \mu_{n,m-\ell}^*) \\ &\sim \sum_{1 \leq \ell \leq m} \lambda_{n,m,\ell}^* \left( \frac{\ell}{m} + \phi'(\alpha) \frac{\ell}{n} \right) \\ &\sim \frac{1}{n} \left( \frac{1}{\alpha} + \phi'(\alpha) \right) \sum_{1 \leq \ell \leq m} \ell \lambda_{n,m,\ell}^* \end{aligned}$$



$$\frac{1}{n} \sim \frac{1}{n} \left( \frac{1}{\alpha} + \phi'(\alpha) \right) \sum_{1 \leq l \leq m} l \lambda_{n,m,l}^*$$

$$\begin{aligned} \sum_{1 \leq l \leq m} l \lambda_{n,m,l}^* &= \sum_{j \geq 1} \binom{n+1-m}{j} n^{-j} \sum_{j < l \leq m} l \binom{m}{l} n^{-l} \\ &\sim \sum_{j \geq 1} \frac{(1-\alpha)^j}{j!} \sum_{l > j} l \frac{\alpha^l}{l!} \\ &= S_1(\alpha) \end{aligned}$$

Then we see that  $\phi$  must satisfy

$$\phi'(x) = \frac{1}{S_1(x)} - \frac{1}{x} = -\frac{3}{2} + \frac{11}{6}x - \dots \implies \boxed{\phi = \phi_1}$$

*The justification relies on a careful error analysis*



# TOOLS NEEDED

**Lemma 1. Asymptotics of  $A_{n,m}^* := \sum_{1 \leq \ell \leq m} a_\ell \lambda_{n,m,\ell}^*$**

**Assume that  $A(z) = \sum_{\ell \geq 1} a_\ell z^{\ell-1}$  has a nonzero radius of convergence in the  $z$ -plane. Then**

$$A_{n,m}^* = \tilde{A}_0(\alpha) - \frac{\tilde{A}_1(\alpha)}{2n} + O(n^{-2}),$$

**where**

$$\tilde{A}_0(\alpha) := \sum_{\ell \geq 1} \frac{\alpha^\ell}{\ell!} \sum_{0 \leq j < \ell} a_{\ell-j} \frac{(1-\alpha)^j}{j!}$$

$$\tilde{A}_1(\alpha) := \sum_{\ell \geq 1} \frac{\alpha^\ell}{\ell!} \sum_{0 \leq j < \ell} a_{\ell-j} \left( \alpha \frac{(1-\alpha)^{j+2}}{(j+2)!} - 2 \frac{(1-\alpha)^{j-1}}{(j-1)!} + (1-\alpha) \frac{(1-\alpha)^{j-2}}{(j-2)!} \right)$$

$$A_{n,m}^* = \frac{1}{2\pi i} \oint_{|z|=c} A(z) \left(1 + \frac{1}{nz}\right)^m \left(1 + \frac{z}{n}\right)^{n+1-m} dz$$



# TOOLS NEEDED

**Lemma 2. (Asymptotic transfer)**

$$\sum_{1 \leq l \leq m} \lambda_{n,m,l}^* (a_{n,m} - a_{n,m-l}) = b_{n,m}$$

**If  $|b_{n,m}| \leq c/n$ , uniformly for  $1 \leq m \leq n$  and  $n \geq 1$ , where  $c > 0$ , then**

$$|a_{n,m}| \leq cH_m \quad (1 \leq m \leq n).$$

**In particular,  $\mu_{n,m}^* \leq H_m$**

$$\Lambda_{n,m}^* := \sum_{1 \leq l \leq m} \lambda_{n,m,l}^* \geq \frac{m}{n} \quad (1 \leq m \leq n)$$

$$|a_{n,m}| \leq \frac{|b_{n,m}|}{\Lambda_{n,m}^*} + |a_{n,m-1}| \leq \frac{c}{n} \cdot \frac{n}{m} + cH_{m-1} = cH_m,$$

**Useful for error analysis**



# TOOLS NEEDED

## Lemma 3

If  $\phi \in C^2[0, 1]$  and  $\phi'(x) \neq 0$  for  $x \in [0, 1]$ , then

$$\begin{aligned} & \sum_{1 \leq \ell \leq m} \lambda_{n,m,\ell}^* \left( \phi \left( \frac{m}{n} \right) - \phi \left( \frac{m-\ell}{n} \right) \right) \\ &= \frac{\phi'(\alpha)}{n} \sum_{1 \leq \ell \leq m} \ell \lambda_{n,m,\ell}^* + O(n^{-2}) \\ &= \frac{\phi'(\alpha) \mathcal{S}_1(\alpha)}{n} + O(n^{-2}) \end{aligned}$$

uniformly for  $1 \leq m \leq n$ .

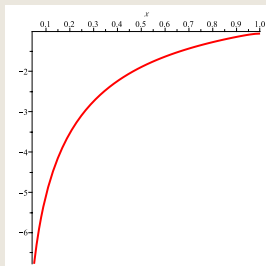
**Bootstrapping & induction**  $\implies \mu_{n,m}^* \sim \sum_{k \geq 0} \frac{b_k H_{m+\phi_{k+1}(\alpha)}}{n^k}$



# INITIAL BITS ARE BERNOULLI( $p$ ):

$$\sum_m \binom{n}{m} p^m q^{n-m} \mu_{n,m}^*$$

$$\begin{aligned} \frac{\mathbb{E}(X_n)}{en} &= \log qn + \gamma + \phi_1(q) \\ &+ \frac{1}{2n} \left( \log qn + \gamma + 3 - \phi_1(q) \right. \\ &\left. + 2q\phi'(q) + pq\phi''(q) + 2\phi_2(q) \right) \\ &+ O(n^{-2} \log n) \end{aligned}$$



$$p = \frac{1}{2}$$

$$\begin{aligned} \frac{\mathbb{E}(X_n)}{en} &= \log n - \log 2 + \gamma + \phi_1\left(\frac{1}{2}\right) + \frac{1}{2n} \left( \log n - \log 2 + \gamma \right. \\ &\left. + 3 - \phi_1\left(\frac{1}{2}\right) + \phi'\left(\frac{1}{2}\right) + \frac{1}{4}\phi''\left(\frac{1}{2}\right) + 2\phi_2\left(\frac{1}{2}\right) \right) + \dots \end{aligned}$$



# VARIANCE OF $X_{n,m}$

Uniformly for  $1 \leq m \leq n$

$$\mathbb{V}(X_{n,m}) = e^2 H_m^{(2)} n^2 - e(2e + 1) \left(n + \frac{1}{2}\right) H_m \\ + \psi_1(\alpha)n + \psi_2(\alpha) + O(n^{-1}H_m)$$

The two dominating terms independent of  $p$

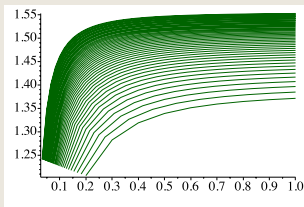
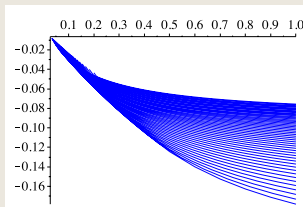
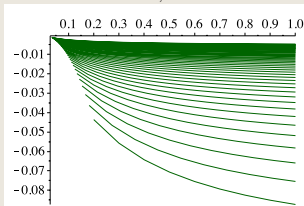
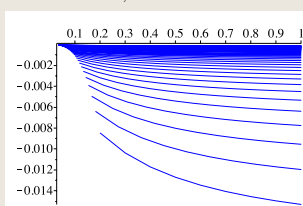
$$\mathbb{V}(X_n) = \frac{e^2 \pi^2}{6} n^2 - e(2e + 1) \left(n + \frac{1}{2}\right) \log n \\ + c'_1 n + c'_2 + O(n^{-1} \log n)$$

$$\psi_1(\alpha) = \int_0^\alpha \left( \frac{S_2(x)}{S_1(x)^3} - \frac{1}{x^2} + \frac{2}{x} \right) dx$$

$$\psi_2(\alpha) = \frac{7}{12} - \int_0^\alpha \left( \frac{5S_1'(x)S_2(x)^2}{2S_1(x)^5} - \frac{2S_1'(x)S_3(x) + S_2(x)S_2'(x) + 6S_0(x)S_2(x)}{2S_1(x)^4} \right. \\ \left. - \frac{S_0(x)}{S_1(x)^3} + \frac{2}{S_1(x)^2} - \frac{1}{x^3} + \frac{3}{x^2} - \frac{11}{2x} \right) dx$$



$$\begin{aligned}
 V_{n,m}^* &:= \frac{\left(1 - \frac{1}{n+1}\right)^{n+1}}{n^2} (\mathbb{V}(X_{n+1,m}) + \mathbb{E}(X_{n+1,m})) \\
 &= H_m^{(2)} + \sum_{1 \leq k < K} \frac{r_k H_m + s_k H_m^{(2)} + t_k}{n^k} + O(H_m n^{-K})
 \end{aligned}$$


 $V_{n,m}^*$ 

 $V_{n,m}^* - H_m^{(2)}$ 

 $K = 2$ 

 $K = 3$



# LIMIT GUMBEL DISTRIBUTION OF $X_{n,m}$

$$\mathbb{P} \left( \sum_{1 \leq r \leq m} \mathbf{Exp}(r) - \log m \leq x \right) \rightarrow e^{-e^{-x}}$$

If  $m \rightarrow \infty$  with  $n$  and  $m \leq n$ , then

$$\mathbb{P} \left( \frac{X_{n,m}}{en} - \log m - \phi_1\left(\frac{m}{n}\right) \leq x \right) \rightarrow e^{-e^{-x}}$$

By induction

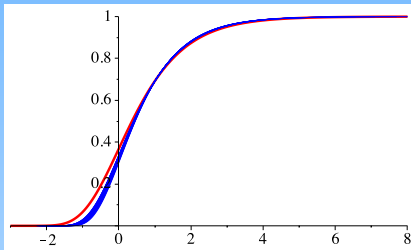
$$\mathbb{E} \left( e^{X_{n,m}s/(en) - (H_m + \phi_1(\frac{m}{n}))s} \right) = \left( 1 + O\left(\frac{H_m}{n}\right) \right) \prod_{1 \leq r \leq m} \frac{e^{-s/r}}{1 - \frac{s}{r}},$$

uniformly for  $1 \leq m \leq n$  (**proof long and messy**).

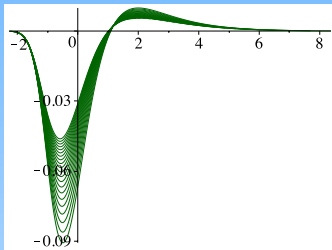


$$X_n = \sum_{0 \leq m \leq n} \binom{n}{m} p^m q^{n-m} X_{n,m}$$

$$\mathbb{P} \left( \frac{X_n}{en} - \log pn - \phi_1(\rho) \leq x \right) \rightarrow e^{-e^{-x}} \quad (x \in \mathbb{R})$$



$$\mathbb{P} \left( \frac{X_n}{en} - \log \frac{n}{2} - \phi_1\left(\frac{1}{2}\right) \leq x \right)$$



**convergence rates**



# MAIN STEPS

$$\begin{aligned} F_{n,m}(s) &:= \frac{\mathbb{E} \left( e^{X_{n,m} s / (en)} \right) e^{-\phi\left(\frac{m}{n}\right)s}}{\prod_{1 \leq r \leq m} \frac{1}{1 - \frac{s}{r}}} \\ &= \frac{Q_{n,m} \left( e^{s/(en)} \right) e^{-H_m s - \phi\left(\frac{m}{n}\right)s}}{\prod_{1 \leq r \leq m} \frac{e^{-s/r}}{1 - \frac{s}{r}}}. \end{aligned}$$

$$F_{n,m}(s) = \frac{\sum_{1 \leq \ell \leq m} \lambda_{n,n-m,\ell} F_{n,m-\ell}(s) e^{-\left(\phi\left(\frac{m}{n}\right) - \phi\left(\frac{m-\ell}{n}\right)\right)s} \prod_{m-\ell+1 \leq r \leq m} \left(1 - \frac{s}{r}\right)}{e^{-s/(en)} - \left(1 - \sum_{1 \leq \ell \leq m} \lambda_{n,n-m,\ell}\right)}$$

**Prove**  $|F_{n,m}(s) - 1| \leq Cn^{-1}H_m$  **when**  $\phi = \phi_1$



# AN AUXILIARY FUNCTION

$$G_{n,m}(s) := \frac{\sum_{1 \leq \ell \leq m} \lambda_{n,n-m,\ell} e^{-\left(\phi\left(\frac{m}{n}\right) - \phi\left(\frac{m-\ell}{n}\right)\right)s} \prod_{m-\ell+1 \leq r \leq m} \left(1 - \frac{s}{r}\right)}{e^{-s/(en)} - \left(1 - \sum_{1 \leq \ell \leq m} \lambda_{n,n-m,\ell}\right) e^{s/(en)}}$$

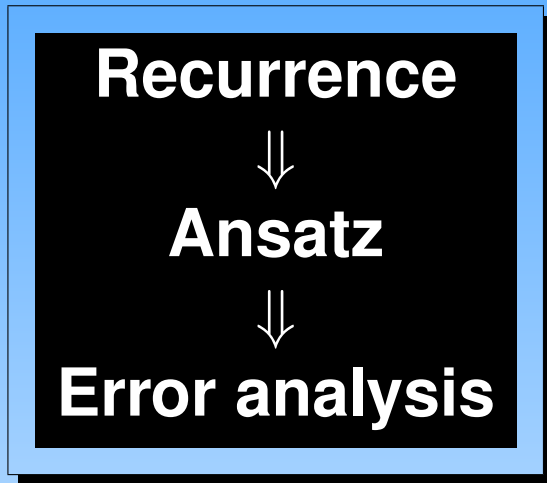
If  $\phi \in C^2[0, 1]$ , then

$$G_{n,m}(s) = \frac{1 - \frac{s}{m} (1 + \alpha \phi'(\alpha)) \frac{S_1(\alpha)}{S(\alpha)} + O\left(\frac{1}{mn}\right)}{1 - \frac{s}{m} \cdot \frac{\alpha}{S(\alpha)} + O\left(\frac{1}{mn}\right)},$$

**If  $\phi = \phi_1$  then  $G_{n,m}(s) = 1 + O((mn)^{-1})$**



# A SUMMARY OF THE APPROACHES



# (1 + 1)-EA FOR LEADINGONES

$$f(\mathbf{x}) = \sum_{1 \leq k \leq n} \prod_{1 \leq j \leq k} x_j; Y_n := \text{optimization time}$$

- Rudolph (1997): introduced LEADINGONES
- Droste et al. (2002):  $\mathbb{E}(Y_n) \asymp n^2$
- Ladret (2005): CLT( $c_1 n^2, c_2 n^3$ )
- Böttcher et al. (2010): re-derived mean
- many other papers

***Prove Ladret's results by a direct analytic approach***



# TIME TO OPTIMUM STATE UNDER LEADING ONES

$Y_n$  (starting with  $n$  random bits (each being 1 with probability  $\frac{1}{2}$ ))

$$\mathbb{E}(e^{Y_n s}) := 2^{-n} + \sum_{1 \leq m \leq n} 2^{m-n-1} Q_{n,m}(s)$$

where the conditional moment generating function  $Q_{n,m}(s)$  satisfies the recurrence relation

$$(1 - (1 - pq^{n-m})e^s) Q_{n,m}(s) = pq^{n-m}e^s \left( 2^{1-m} + \sum_{1 \leq \ell < m} \frac{Q_{n,\ell}(s)}{2^{m-\ell}} \right),$$

for  $1 \leq m \leq n$ , where  $q = 1 - p$ .

$$p \asymp n^{-1}$$



# CLOSED-FORM SOLUTION FOR $Q_{n,m}$

$$Q_{n,m}(s) = \frac{1}{1 - \frac{1-e^{-s}}{pq^{n-m}}} \prod_{1 \leq j < m} \frac{1 - \frac{1-e^{-s}}{2pq^{n-j}}}{1 - \frac{1-e^{-s}}{pq^{n-j}}}$$

$$Y_{n,m} \stackrel{d}{=} \underbrace{Z_{n,m}^{[0]}}_{\text{geom}} + \underbrace{Z_{n,m}^{[m-1]} + \dots + Z_{n,m}^{[m-1]}}_{\frac{1}{2} + \frac{1}{2} \text{geom}}$$

$$R_m(t) := \mathbb{E} \left( t^{Z_{n,m}^{[0]}} \right) = \frac{pq^{n-m}t}{1 - (1 - pq^{n-m})t}$$

$$\mathbb{E} \left( t^{Z_{n,m}^{[j]}} \right) = \frac{1}{2} \cdot \frac{1 - (1 - 2pq^{n-j})t}{1 - (1 - pq^{n-j})t} = \frac{1}{2} + \frac{R_j(t)}{2}$$

$(j = 1, \dots, m-1)$





$$\frac{Y_n - \mathbb{E}(Y_n)}{\sqrt{\mathbb{V}(Y_n)}} \longrightarrow \mathcal{N}(0, 1)$$

$$\begin{aligned} \mathbb{E}(Y_{n,m}) &= \frac{1}{pq^{n-1}} \left( \frac{1 - q^{m-1}}{2p} + q^{m-1} \right) \\ &\stackrel{p = \frac{c}{n}}{=} \frac{n^2}{2c^2} \left( e^c - e^{c(1-\alpha)} + O\left(\frac{c(c+1)}{n}\right) \right). \end{aligned}$$

$$\begin{aligned} \mathbb{E}(Y_n) &= \sum_{1 \leq m \leq n} 2^{-n+m-1} \mathbb{E}(Y_{n,m}) = \frac{q}{2p^2} (q^{-n} - 1) \\ &= \frac{e^c - 1}{2c^2} n^2 + \frac{(c-2)e^c + 2}{4c} n + \frac{ce^c(3c-4)}{48} + \dots \end{aligned}$$

$$\begin{aligned} \mathbb{V}(Y_n) &= \frac{3q^2}{4p^3(1+q)} (q^{-2n} - 1) - \mu_n \\ &\stackrel{p = \frac{c}{n}}{=} \frac{e^{2c} - 1}{8c^3} n^3 + \frac{3e^{2c}(2c-3) - 8e^c + 17}{16c^2} n^2 \\ &\quad + \frac{(6c^2 - 10c + 3)e^{2c} - 8(c-2)e^c - 19}{32c} n + O(1) \end{aligned}$$



Properties	ONEMAX ( $X_n$ )	LEADINGONES ( $Y_n$ )
Mean $\sim$	$e n \log n + c_1 n$	$\frac{e-1}{2} n^2$
Variance $\sim$	$\frac{\pi^2}{6} (en)^2 - (2e+1)en \log n$	$\frac{e^2-1}{8} n^3$
Limit law	<p><b>Gumbel distribution</b></p> $\mathbb{P} \left( \frac{X_n}{en} - \log \frac{n}{2} - \phi_1 \left( \frac{1}{2} \right) \leq x \right)$ $\rightarrow e^{-e^{-x}}$	<p><b>Gaussian distribution</b></p> $\mathbb{P} \left( \frac{Y_n - \frac{e-1}{2} n^2}{\sqrt{\frac{e^2-1}{8} n^3}} \leq x \right)$ $\rightarrow \frac{1}{\sqrt{2\pi}} \int_{-\infty}^x e^{-\frac{t^2}{2}} dt$
Approach	Ansatz & error analysis	Analytic combinatorics



THANK YOU

**1+1=2**

**(1 + 1)-EA = 232**

