

Antibandwidth of Hamming graphs

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Workshop on Discrete Mathematics

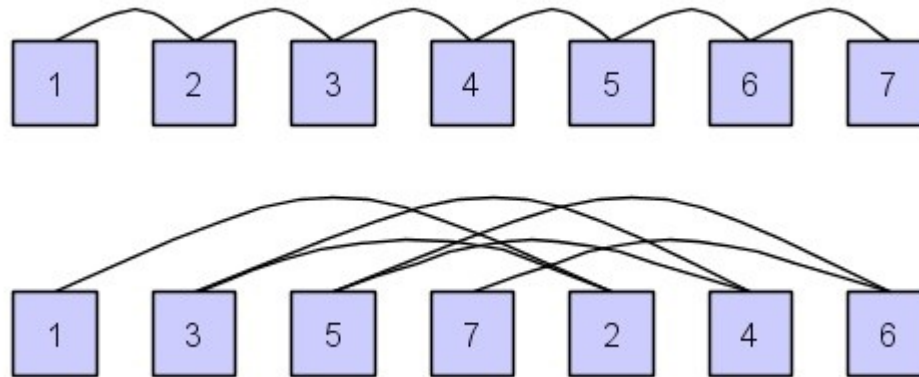
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Antibandwidth problem

- Consists of placing the vertices of a graph on a line in consecutive integer points in such a way that the minimum difference of adjacent vertices is maximized.
- Just another labeling problem (from graph theory point of view)

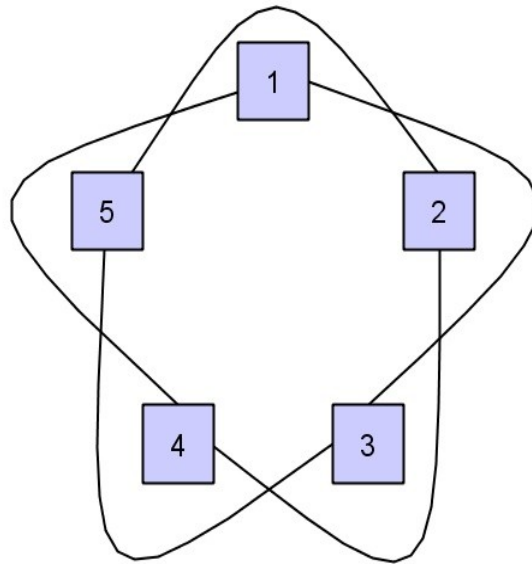
Antibandwidth problem

- Example (linear)



Cyclic antibandwidth

- Example (cyclic)



Motivation

- Introduced in connection with multiprocessor scheduling problems
- Dual problem to bandwidth problem
 - Bandwidth = to find vertex ordering of a graph G s.t. the length of the longest edge is minimized.
 - Antibandwidth = to find vertex ordering of a graph G s.t. the length of the shortest edge is maximized.

Motivation

- Radio frequency assignment problem
N transmitters, N frequencies.

Graph representation: transmitters = vertices,
edges = possible interference

Goal: to assign all frequencies to transmitters in a one-to-one manner s.t. interfering transmitters have as different frequencies as possible.

Motivation

- Enemy facility location problem

Graph representation: vertices = facilities, edges = hostility

Goal: to arrange facilities in a line (cycle, mesh,...) s.t. the critical (smallest) distance between two enemies is maximized

Motivation

- Coding theory

Host graph: n -dimensional hypercube

Guest graph: complete graph K_p

Then $ab(K_p, Q_n) \geq A$ is equivalent to an existence of a code with minimal Hamming distance greater or equal to A .

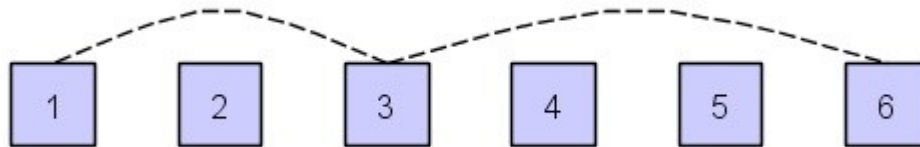
Previous results

- NP-complete (Leung, Vornberger)
- Polynomially solvable for the complements of
 - Interval
 - Arborescent comparability
 - Thresholdgraphs.
(Donnely, Isaak, *Hamiltonian powers in ... graphs*)

Previous results

- Trivial upper bound for connected graphs:

$$ab(G) \leq \lfloor \frac{n}{2} \rfloor$$



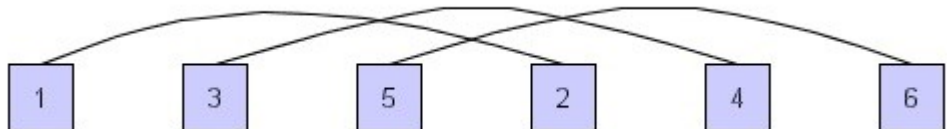
Previous results

- Also interesting for disconnected graphs.
 - Exact values for graphs consisting of copies of simple graphs.

Bandwidth:



Antibandwidth:



Previous results

- Known results for:
 - Paths, cycles, special trees
 - complete and complete bipartite graphs
 - 2D, 3D and n-dimensional meshes
 - tori, hypercubes
- Cyclic antibandwidth
 - Embedding into cycle

Previous results

Raspaud, Schröder, Sýkora, Török, Vrto:

Antibandwidth and Cyclic Antibandwidth of Meshes and Hypercubes

Calamoneri, Massini, Török, Vrto:

Antibandwidth of complete k -ary trees

Török, Vrto:

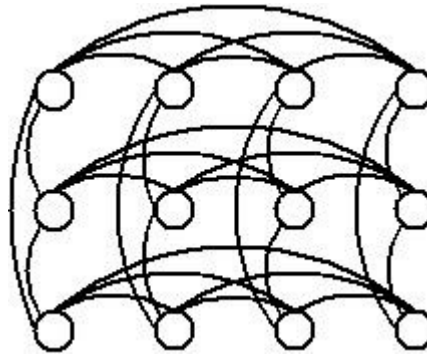
Antibandwidth of three-dimensional meshes

Antibandwidth of d -dimensional meshes

Hamming graph

- Definition
 - Cartesian product of d complete graphs

$K_{n_k}, k=1,2,\dots,d$, i.e. $\Pi_{k=1}^d K_{n_k}$



Hamming graph $K_3 \times K_4$

Hamming graph

- Definition

- Vertices are d-tuples

$$(i_1, i_2, \dots, i_d), i_k \in 0, 1, 2, \dots, n_k - 1$$

- Two vertices (i_1, i_2, \dots, i_d) and (j_1, j_2, \dots, j_d) are adjacent iff the two d-tuples differ in precisely one coordinate

Results

- Work in progress
- Theorem

*For $d \geq 2$,
 $2 \leq n_1 \leq n_2 \leq \dots \leq n_d$*

$ab(\prod_{k=1}^d K_{n_k}) = n_1 n_2 \dots n_{d-1}$, if $n_{d-1} \neq n_d$

$ab(\prod_{k=1}^d K_{n_k}) = n_1 n_2 \dots n_{d-1} - 1$, if $n_{d-1} = n_d$, $n_{d-2} \neq n_{d-1}$

$n_1 n_2 \dots n_{d-1} - \min(n_1 n_2 \dots n_{d-2}, n_{q+1} \dots n_{d-1}) \leq ab(\prod_{k=1}^d K_{n_k}) \leq n_1 n_2 \dots n_{d-1} - 1$

where q is minimal index such that $q \leq d-2$ and $n_q = n_d$

Results

- Labelling

- Case I. $n_{d-1} \neq n_d$

$$\begin{aligned} f(i_1, i_2, \dots, i_d) = & ((N_{d-1} + N_{d-1}/N_1)i_1 \\ & + (N_{d-1} + N_{d-1}/N_2)i_2 + (N_{d-1} + N_{d-1}/N_3)i_3 + \dots \\ & + (N_{d-1} + N_{d-1}/N_{d-1})i_{d-1} + N_{d-1}i_d) \bmod N_d \end{aligned}$$

$N_j = \left| \prod_{k=1}^j K_{n_k} \right|$, number of vertices of the product

- Example $K_2 \times K_3 \times K_4$

$$f(i_1, i_2, i_3) = (9i_1 + 7i_2 + 6i_3) \bmod 24$$

Results

- Example of labelling of $K_2 \times K_3 \times K_4$

0	7	14
9	16	23

6	13	20
15	22	5

12	19	2
21	4	11

18	1	8
3	10	17

Results

- Labelling

- Case II. $n_{d-1} = n_d$, let q be minimal s.t.

- $q \leq d-1$ and $n_q = n_d$

$$\begin{aligned} f(i_1, i_2, \dots, i_d) = & ((N_{d-1} + N_{d-1}/N_1)i_1 \\ & + (N_{d-1} + N_{d-1}/N_2)i_2 + \dots + (N_{d-1} + N_{d-1}/N_{q-1})i_{q-1} \\ & + (N_{d-1} - N_{d-1}/N_q)i_q + \dots \\ & + (N_{d-1} - N_{d-1}/N_{d-1})i_{d-1} + N_{d-1}i_d) \bmod N_d \end{aligned}$$

Results

- Labelling

- Case III. $n_{d-1} = n_d$, define the alternative labelling

$$f(i_1, i_2, \dots, i_d) = ((N_{d-1} - N_0)i_1 + (N_{d-1} - N_1)i_2 + \dots + (N_{d-1} - N_{d-2})i_{d-1} + N_{d-1}i_d) \bmod N_d$$

Results

- Antibandwidth and cyclic antibandwidth have the same value
- Open problem

$$n^2 - n \leq ab(K_n \times K_n \times K_n) \leq n^2 - 1$$

- Conjecture

$$ab(K_n \times K_n \times K_n) = n^2 - n$$

The End

Thank You