

# On the second Lyapunov exponent of some multidimensional continued fraction algorithms

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Vienna / Moscow (via zoom), July 2022

# Multidimensional continued fraction algorithms

## Two spaces

$$\Lambda = \{(y_0, y_1, \dots, y_d) \in \mathbb{R}^{d+1} \setminus \{\mathbf{0}\} : y_0 \geq y_1 \geq \dots \geq y_d \geq 0\},$$

$$\Delta = \{(x_1, \dots, x_d) \in \mathbb{R}^d : \mathbf{1} \geq x_1 \geq \dots \geq x_d \geq 0\},$$

We can switch between these spaces.

$$\iota : \Delta \rightarrow \Lambda, \quad (x_1, \dots, x_d) \mapsto (1, x_1, \dots, x_d),$$

$$\kappa : \Lambda \rightarrow \Delta, \quad (y_0, \dots, y_d) \mapsto \left( \frac{y_1}{y_0}, \dots, \frac{y_d}{y_0} \right).$$

The **multidimensional continued fraction algorithm (MCF)** is defined in terms of

$$A : \Delta \rightarrow \text{GL}(d+1, \mathbb{Z}),$$

which selects a matrix for each element of  $\Delta$  (or of  $\Lambda/\mathbb{R}$ ).

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# The mappings

We assume that  $A$  is defined in a way that the **linear version** of the MCF

$$L : \Lambda \rightarrow \Lambda, \quad \mathbf{y} \mapsto \mathbf{y}A(\kappa(\mathbf{y}))^{-1} = \mathbf{y}A\left(\frac{y_1}{y_0}, \dots, \frac{y_d}{y_0}\right)^{-1}$$

is well defined (i.e.,  $L$  maps  $\Lambda$  into itself).

The “**projectivized**” version of the MCF is defined by the commutative diagram

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# Regular continued fraction algorithm

Regular continued fraction algorithm ( $d = 1$ ):

$$\Lambda = \{(y_0, y_1) \in \mathbb{R}^2 \setminus \{\mathbf{0}\} : y_0 \geq y_1 \geq 0\}, \quad \Delta = \{x \in [0, 1]\},$$

Here

$$A(x) = \begin{pmatrix} \lfloor \frac{1}{x} \rfloor & 1 \\ 1 & 0 \end{pmatrix}.$$

Thus we have

$$L(y_0, y_1) = (y_0, y_1) A(y_1/y_0)^{-1}$$

and the Gauss map

$$T(x) = \kappa \left( (1, x) \begin{pmatrix} 0 & 1 \\ 1 & -\lfloor \frac{1}{x} \rfloor \end{pmatrix} \right) = \kappa(x, 1 - \lfloor \frac{1}{x} \rfloor x) = \frac{1}{x} - \left\lfloor \frac{1}{x} \right\rfloor.$$

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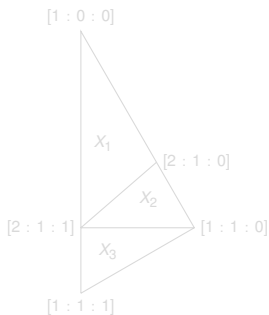
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# Brun's algorithm

Brun's continued fraction algorithm ( $d = 2$ ):

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$$A(y_1, y_2) = M_i \text{ iff } (1, y_1, y_2) \in X_i$$

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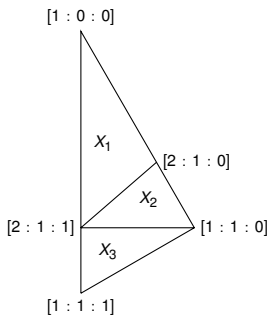
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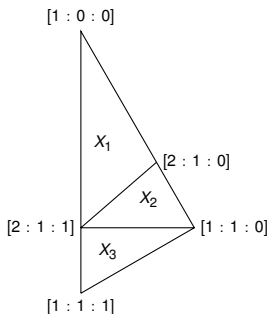
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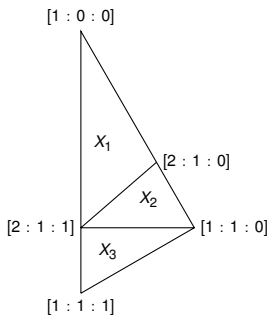
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# Convergents, linear cocycle

The convergents  $\frac{\mathbf{p}_i^{(n)}}{q_i^{(n)}} = \frac{(p_{i,1}^{(n)}, \dots, p_{i,d}^{(n)})}{q_i^{(n)}}$  of  $\mathbf{x} \in \Delta$  are given by the cocycle

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Weak convergence:  $\lim_{n \rightarrow \infty} \frac{\mathbf{p}_i^{(n)}}{q_i^{(n)}} = \mathbf{x}$  for all  $0 \leq i \leq d$

Strong convergence:  $\lim_{n \rightarrow \infty} \|\mathbf{p}_i^{(n)} - q_i^{(n)} \mathbf{x}\| = 0$  for all  $0 \leq i \leq d$

Exponential convergence:  $\|\mathbf{p}_i^{(n)} - q_i^{(n)} \mathbf{x}\| < ce^{-\delta n}$  for all  $0 \leq i \leq d, n \in \mathbb{N}$ .

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# Lyapunov Exponents and Exponential convergence

We assume that  $(\Delta, T, \mu)$  is ergodic for an invariant probability measure  $\mu$ , the cocycle  $A$  is log-integrable, i.e.,

$$\int_X \log \max(1, \|A(x)\|) d\nu(x) < \infty,$$

with Lyapunov exponents

$$\lambda_1(A) \geq \lambda_2(A) \geq \lambda_3(A) \geq \cdots \geq \lambda_{d+1}(A)$$

$\lambda_2(A) < 0 \iff$  almost everywhere exponential convergence

**Problem:** For a given algorithm prove that  $\lambda_2(A) < 0$ .

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# Approximation cocycle

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$$\Pi = \begin{pmatrix} 0 & 1 & 0 & 0 \\ \vdots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & 1 \end{pmatrix}, \quad H(x_1, \dots, x_d) = \begin{pmatrix} -x_1 & \cdots & -x_d \\ 1 & \cdots & 0 \\ 0 & \ddots & 0 \\ 0 & \cdots & 1 \end{pmatrix}.$$

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# Determining $\lambda_2(A) = \lambda_1(D)$

$$|\det A(\mathbf{x})| = 1 \Rightarrow \sum_{k=1}^{d+1} \lambda_k(A) = 0$$

$$d = 2: \lambda_2(A) = -\lambda_1(A) - \lambda_3(A) = \lambda_1({}^tA^{-1}) - \lambda_1(A)$$

can be used to show  $\lambda_2(A) < 0$  for **Jacobi–Perron** (**Paley–Ursell’30**).

$d \geq 2$ : ergodicity, subadditivity  $\Rightarrow$

$$\lambda_1(D) = \inf_{n \in \mathbb{N}} \frac{1}{n} \int_{\Delta} \log \|D^{(n)}(\mathbf{x})\| d\mu(\mathbf{x})$$

gives upper bounds for  $\lambda_2(A)$

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# Selmer algorithm (ordered)

$$\Delta = \{(x_1, \dots, x_d) \in \mathbb{R}^d : 1 \geq x_1 \geq \dots \geq x_d \geq 0, x_{d-1} + x_d \geq 1\}$$

$$T_S(x_1, \dots, x_d) = \kappa(\text{ord}(1-x_d, x_1, \dots, x_d)) = \text{ord}\left(\frac{x_2}{x_1}, \dots, \frac{x_d}{x_1}, \frac{1-x_d}{x_1}\right)$$

$$A_S(\mathbf{x}) = \begin{cases} S_a & \text{if } x_d > 1/2, \\ S_b & \text{if } x_d < 1/2, \end{cases}$$

$$S_a = \begin{pmatrix} 0 & 1 & 0 & \dots & 0 \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \dots & 0 & 1 & 0 \\ 1 & 0 & \dots & 0 & 1 \\ 1 & 0 & \dots & 0 & 0 \end{pmatrix}, \quad S_b = \begin{pmatrix} 0 & 1 & 0 & \dots & 0 \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \dots & 0 & 1 & 0 \\ 1 & 0 & \dots & 0 & 0 \\ 1 & 0 & \dots & 0 & 1 \end{pmatrix}$$

invariant measure  $\mu_S$  has density  $\frac{c}{x_1 \dots x_d}$  (Bruin, Fokkink, Kraaikamp'15).

# Selmer algorithm (ordered)

$$\Delta = \{(x_1, \dots, x_d) \in \mathbb{R}^d : 1 \geq x_1 \geq \dots \geq x_d \geq 0, x_{d-1} + x_d \geq 1\}$$

$$T_S(x_1, \dots, x_d) = \kappa(\text{ord}(1-x_d, x_1, \dots, x_d)) = \text{ord}\left(\frac{x_2}{x_1}, \dots, \frac{x_d}{x_1}, \frac{1-x_d}{x_1}\right)$$

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# Selmer approximation cocycle

$$D_S(\mathbf{x}) = \begin{pmatrix} 0 & 1 & 0 & \cdots & 0 \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & 1 & 0 \\ -x_1 & -x_2 & \cdots & -x_{d-1} & 1 - x_d \\ -x_1 & -x_2 & \cdots & -x_{d-1} & -x_d \end{pmatrix} \quad \text{if } x_d > 1/2$$

$$D_S(\mathbf{x}) = \begin{pmatrix} 0 & 1 & 0 & \cdots & 0 \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & 1 & 0 \\ -x_1 & -x_2 & \cdots & -x_{d-1} & -x_d \\ -x_1 & -x_2 & \cdots & -x_{d-1} & 1 - x_d \end{pmatrix} \quad \text{if } x_d < 1/2$$

Selmer algorithm,  $d = 2$ 

$$D_S^{(2)}(\mathbf{x}) = \begin{cases} \begin{pmatrix} 1 - x_1 & -x_2 \\ 1 & 0 \end{pmatrix} & \text{if } A_S^{(2)}(\mathbf{x}) = S_a^2 = \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 0 \end{pmatrix} \\ \begin{pmatrix} 1 - x_1 & 1 - x_2 \\ 1 & 0 \end{pmatrix} & \text{if } A_S^{(2)}(\mathbf{x}) = S_a S_b = \begin{pmatrix} 1 & 0 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 0 \end{pmatrix} \\ \begin{pmatrix} 1 & 0 \\ 1 - x_1 & -x_2 \end{pmatrix} & \text{if } A_S^{(2)}(\mathbf{x}) = S_b S_a = \begin{pmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 1 & 0 \end{pmatrix} \\ \begin{pmatrix} 1 & 0 \\ 1 - x_1 & 1 - x_2 \end{pmatrix} & \text{if } A_S^{(2)}(\mathbf{x}) = S_b^2 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 1 & 1 \end{pmatrix} \end{cases}$$

$$\|D_S^{(2)}(\mathbf{x})\|_\infty = 1 \text{ for all } \mathbf{x} \in \Delta \Rightarrow \lambda_2(A_S) \leq 0$$

$$\|D_S^{(4)}(\mathbf{x})\|_\infty < 1 \text{ for some } \mathbf{x} \in \Delta \Rightarrow \lambda_2(A_S) < 0$$

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# Selmer $d = 2$ and $d = 3$

## Theorem (Berthé, Steiner, T. '21)

For Selmer with  $d = 2$  we have

$$\lambda_2(\mathbf{A}_S) \leq \frac{1}{50} \int_{\Delta} \log \|D_S^{(50)}(\mathbf{x})\| d\mu_S(\mathbf{x}) < -0.052435991.$$

*This implies exponential convergence.*

Extensive (exact) computations using GPUs yield:

## Theorem (Berthé, Steiner, T. '21)

For Selmer with  $d = 3$  we have

$$\lambda_2(\mathbf{A}_S) \leq \frac{1}{52} \int_{\Delta} \log \|D_S^{(52)}(\mathbf{x})\| d\mu_S(\mathbf{x}) < -0.000436459.$$

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# Selmer matrices

## Theorem (Berthé, Steiner, T. '21)

Let  $d = 2$  and  $M \in \{S_a, S_b\}^n$  for some  $n \geq 1$ . The following assertions are equivalent.

- 1  $M$  is a primitive matrix,
- 2  $M$  is a Pisot matrix,
- 3  $M^2 \notin \{S_a S_b, S_b^2\}^n$ .

Cf. results of [Avila–Delecroix'19](#)) for [Brun](#) matrices ( $d = 2$ ) and [Arnoux–Rauzy](#) matrices ( $d \geq 2$ )

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# Selmer algorithm, computer simulations

Calculating  $D_S^{(n)}(\mathbf{x})$  and  $A_S^{(n)}(\mathbf{x})$  for  $n = 2^{30}$  for randomly chosen points  $\mathbf{x} \in \Delta$  gives the following estimates for  $d \in \{2, 3, 4, 5\}$  (without guaranteed accuracy)

$d$	$\lambda_2(A_S)$	$1 - \frac{\lambda_2(A_S)}{\lambda_1(A_S)}$
2	-0.07072	1.3871
3	-0.02283	1.1444
4	+0.00176	0.9866
5	+0.01594	0.8577

**Lagarias 1993:** Uniform approximation exponent

$$\sup \left\{ \delta : \left\| \mathbf{x} - \frac{\mathbf{p}_i^{(n)}}{q_i^{(n)}} \right\| = O\left(\frac{1}{(q_i^{(n)})^\delta}\right), 0 \leq i \leq d \right\} = 1 - \frac{\lambda_2(A)}{\lambda_1(A)} \quad (\text{a.e. } \mathbf{x} \in \Delta)$$

# Computer simulations

We calculate  $D^{(n)}(\mathbf{x})$  for randomly chosen points  $\mathbf{x} \in \Delta$  by a C program with double precision floating point arithmetic. We have to renormalize the matrices before they get too small or too large. For  $k|n$ , e.g.,  $k = 2^{10}$ ,  $n = 2^{30}$ , we have

$$\begin{aligned} & \frac{1}{D_{1,1}^{(n)}(\mathbf{x})} D^{(n)}(\mathbf{x}) \\ &= \frac{D_{1,1}^{(n-k)}(\mathbf{x})}{D_{1,1}^{(n)}(\mathbf{x})} D^{(k)}(T^{n-k}\mathbf{x}) \cdots \frac{D_{1,1}^{(k)}(\mathbf{x})}{D_{1,1}^{(2k)}(\mathbf{x})} D^{(k)}(T^k\mathbf{x}) \frac{1}{D_{1,1}^{(k)}(\mathbf{x})} D^{(k)}(\mathbf{x}) \\ \log |D_{1,1}^{(n)}(\mathbf{x})| &= \log \frac{|D_{1,1}^{(n)}(\mathbf{x})|}{|D_{1,1}^{(n-k)}(\mathbf{x})|} + \cdots + \log \frac{|D_{1,1}^{(2k)}(\mathbf{x})|}{|D_{1,1}^{(k)}(\mathbf{x})|} + \log |D_{1,1}^{(k)}(\mathbf{x})| \end{aligned}$$

$(D_{1,1}^{(n)}(\mathbf{x}))$  is the top left coefficient of  $D^{(n)}(\mathbf{x})$

# Jacobi–Perron algorithm

A projectivized version of the **Jacobi–Perron algorithm** is given by

$$T_J : [0, 1]^d \rightarrow [0, 1]^d,$$

$$(x_1, x_2, \dots, x_d) \mapsto \left( \frac{x_2}{x_1} - \left\lfloor \frac{x_2}{x_1} \right\rfloor, \dots, \frac{x_d}{x_1} - \left\lfloor \frac{x_d}{x_1} \right\rfloor, \frac{1}{x_1} - \left\lfloor \frac{1}{x_1} \right\rfloor \right).$$

Its matrix version is therefore

$$(x_0, x_1, \dots, x_d) \mapsto \left( x_1, x_2 - \left\lfloor \frac{x_2}{x_1} \right\rfloor x_1, \dots, x_d - \left\lfloor \frac{x_d}{x_1} \right\rfloor x_1, x_0 - \left\lfloor \frac{x_0}{x_1} \right\rfloor x_1 \right)$$

(**not ordered**; thus it is defined on the whole cube).

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## Jacobi–Perron algorithm, computer simulations

$$A_{\text{JP}}(x_1, \dots, x_d) = \begin{pmatrix} \lfloor \frac{1}{x_1} \rfloor & 1 & \lfloor \frac{x_2}{x_1} \rfloor & \dots & \lfloor \frac{x_d}{x_1} \rfloor \\ 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & \vdots & & \ddots & 1 \\ 1 & 0 & \dots & \dots & 0 \end{pmatrix}$$

$$D_{\text{JP}}(x_1, \dots, x_d) = \begin{pmatrix} 0 & 1 & \dots & 0 \\ \vdots & \ddots & \ddots & \vdots \\ 0 & \dots & 0 & 1 \\ -x_1 & \dots & -x_{d-1} & -x_d \end{pmatrix}$$

$d$	$\lambda_2(A_{\text{JP}})$	$1 - \frac{\lambda_2(A_{\text{JP}})}{\lambda_1(A_{\text{JP}})}$	$d$	$\lambda_2(A_{\text{JP}})$	$1 - \frac{\lambda_2(A_{\text{JP}})}{\lambda_1(A_{\text{JP}})}$
2	-0.44841	1.3735	7	-0.02819	1.0243
3	-0.22788	1.1922	8	-0.01470	1.0127
4	-0.13062	1.1114	9	-0.00505	1.0044
5	-0.07880	1.0676	10	+0.00217	0.9981
6	-0.04798	1.0413	11	+0.00776	0.9933

(conjecture in [Hardcastle–Khanin'00](#):  $\lambda_2(A_{\text{JP}}) \approx \frac{c}{d} < 0$ )

## Jacobi–Perron algorithm, computer simulations

$$A_{JP}(x_1, \dots, x_d) = \begin{pmatrix} \lfloor \frac{1}{x_1} \rfloor & 1 & \lfloor \frac{x_2}{x_1} \rfloor & \dots & \lfloor \frac{x_d}{x_1} \rfloor \\ 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & \vdots & & \ddots & 1 \\ 1 & 0 & \dots & \dots & 0 \end{pmatrix}$$

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$d$	$\lambda_2(A_{JP})$	$1 - \frac{\lambda_2(A_{JP})}{\lambda_1(A_{JP})}$	$d$	$\lambda_2(A_{JP})$	$1 - \frac{\lambda_2(A_{JP})}{\lambda_1(A_{JP})}$
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# Nearest integer Jacobi–Perron algorithm

$$\Delta = [-1/2, 1/2]^d$$

$$A_{\text{NIJP}}(x_1, \dots, x_d) = \begin{pmatrix} \lfloor \frac{1}{x_1} + \frac{1}{2} \rfloor & 1 & \lfloor \frac{x_2}{x_1} + \frac{1}{2} \rfloor & \cdots & \lfloor \frac{x_d}{x_1} + \frac{1}{2} \rfloor \\ 0 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & \vdots & & \ddots & 1 \\ 1 & 0 & \cdots & \cdots & 0 \end{pmatrix}$$

$$D_{\text{NIJP}}(x_1, \dots, x_d) = \begin{pmatrix} 0 & 1 & \cdots & 0 \\ \vdots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & 1 \\ -x_1 & \cdots & -x_{d-1} & -x_d \end{pmatrix}$$

$d$	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$	$d$	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$	$d$	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$	$d$	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$
2	1.40145	6	1.06898	10	1.01737	14	0.99924
3	1.22519	7	1.04944	11	1.01125	15	0.99657
4	1.14373	8	1.03551	12	1.00639		
5	1.09786	9	1.02521	13	1.00246		

# Nearest integer Jacobi–Perron algorithm

$$\Delta = [-1/2, 1/2]^d$$

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$$D_{\text{NIJP}}(x_1, \dots, x_d) = \begin{pmatrix} 0 & 1 & \cdots & 0 \\ \vdots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & 1 \\ -x_1 & \cdots & -x_{d-1} & -x_d \end{pmatrix}$$

$d$	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$	$d$	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$	$d$	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$	$d$	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$
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## Brun and modified Jacobi–Perron algorithms (ordered)

$$\Delta = \{(x_1, \dots, x_d) \in \mathbb{R}^d : 1 \geq x_1 \geq \dots \geq x_d \geq 0\}$$

$$T_B(x_1, \dots, x_d) = \kappa(\text{ord}(1 - x_1, x_2, \dots, x_d))$$

$$A_B(\mathbf{x}) = B_k \quad \text{if } x_k > 1 - x_1 > x_{k+1} \quad (0 \leq k \leq d, x_0 = 1, x_{d+1} = 0)$$

$$B_0 = \begin{pmatrix} 1 & 0 & \cdots & \cdots & 0 \\ 1 & 1 & \ddots & & \vdots \\ 0 & 0 & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & 1 & 0 \\ 0 & 0 & \cdots & 0 & 1 \end{pmatrix}, \quad B_k = \begin{pmatrix} 1 & 1 & 0 & \cdots & \cdots & \cdots & 0 \\ k-1 \left\{ \begin{array}{l} 0 & 0 & 1 & \ddots & & & \vdots \\ \vdots & 0 & \ddots & \ddots & \ddots & & \vdots \\ 0 & \vdots & \ddots & \ddots & 1 & \ddots & \vdots \\ 1 & \vdots & & \ddots & 0 & 0 & \ddots \\ 0 & \vdots & & & \ddots & 1 & \ddots \\ 0 & 0 & \cdots & \cdots & \cdots & \cdots & 0 & 1 \end{array} \right. \end{pmatrix}$$

$$A_{\text{MJP}}(\mathbf{x}) = A_B^{(n)}(\mathbf{x}) \quad \text{with } n = \min\{k \geq 1 : A_B(T^{k-1}\mathbf{x}) \neq B_0\}$$

$d$	$\lambda_2(A_B)$	$1 - \frac{\lambda_2(A_B)}{\lambda_1(A_B)}$	$d$	$\lambda_2(A_B)$	$1 - \frac{\lambda_2(A_B)}{\lambda_1(A_B)}$
2	-0.11216	1.3683	7	-0.01210	1.0493
3	-0.07189	1.2203	8	-0.00647	1.0283
4	-0.04651	1.1504	9	-0.00218	1.0102
5	-0.03051	1.1065	10	+0.00115	0.9943
6	-0.01974	1.0746	11	+0.00381	0.9799

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$$A_B(\mathbf{x}) = B_k \quad \text{if } x_k > 1 - x_1 > x_{k+1} \quad (0 \leq k \leq d, x_0 = 1, x_{d+1} = 0)$$

$$B_0 = \begin{pmatrix} 1 & 0 & \cdots & \cdots & 0 \\ 1 & 1 & \ddots & & \vdots \\ 0 & 0 & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & 1 & 0 \\ 0 & 0 & \cdots & 0 & 1 \end{pmatrix}, \quad B_k = \begin{pmatrix} 1 & 1 & 0 & \cdots & \cdots & \cdots & 0 \\ k-1 \left\{ \begin{array}{l} 0 & 0 & 1 & \ddots & & & \vdots \\ \vdots & 0 & \ddots & \ddots & \ddots & & \vdots \\ 0 & \vdots & \ddots & \ddots & 1 & \ddots & \vdots \\ 1 & \vdots & & \ddots & 0 & 0 & \ddots \\ 0 & \vdots & & & \ddots & 1 & \ddots \\ 0 & 0 & \cdots & \cdots & \cdots & \cdots & 0 & 1 \end{array} \right. \end{pmatrix}$$

$$A_{\text{MJP}}(\mathbf{x}) = A_B^{(n)}(\mathbf{x}) \quad \text{with } n = \min\{k \geq 1 : A_B(T^{k-1}\mathbf{x}) \neq B_0\}$$

$d$	$\lambda_2(A_B)$	$1 - \frac{\lambda_2(A_B)}{\lambda_1(A_B)}$	$d$	$\lambda_2(A_B)$	$1 - \frac{\lambda_2(A_B)}{\lambda_1(A_B)}$
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# Garrity's triangle algorithm (ordered)

$$\Delta = \{(x_1, \dots, x_d) \in \mathbb{R}^d : 1 \geq x_1 \geq \dots \geq x_d \geq 0\}$$

$$T_G(x_1, \dots, x_d) = \begin{cases} \kappa(\text{ord}(1 - \sum_{j=1}^k x_j, x_1, \dots, x_d)) & \text{if } 0 < 1 - \sum_{j=1}^k x_j < x_{k+1}, 1 \leq k \leq d-2 \\ \frac{1}{x_1}(x_2, \dots, x_d, 1 - \sum_{j=1}^{d-1} x_j - \ell x_d) & \text{if } 0 < 1 - \sum_{j=1}^{d-1} x_j - \ell x_d < x_d, \ell \geq 0 \end{cases}$$

$d = 2 :$

$$A_G(\mathbf{x}) = \begin{pmatrix} 1 & 1 & 0 \\ \lfloor \frac{1-x_1}{x_2} \rfloor & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix}, \quad D_G(\mathbf{x}) = \begin{pmatrix} -\lfloor \frac{1-x_1}{x_2} \rfloor x_1 & 1 - \lfloor \frac{1-x_1}{x_2} \rfloor x_2 \\ -x_1 & -x_2 \end{pmatrix}$$

Problem:  $\lfloor \frac{1-x_1}{x_2} \rfloor x_1$  can be large

$d$	$\lambda_2(A_G)$	$1 - \frac{\lambda_2(A_G)}{\lambda_1(A_G)}$	$d$	$\lambda_2(A_G)$	$1 - \frac{\lambda_2(A_G)}{\lambda_1(A_G)}$
2	+0.34434	0.6859	7	-0.00644	1.0225
3	+0.37673	0.5798	8	-0.00768	1.0304
4	+0.25232	0.6286	9	-0.00435	1.0189
5	+0.10677	0.7778	10	-0.00074	1.0035
6	+0.01859	0.9468	11	+0.00237	0.9880

# Intermediate algorithm between Arnoux–Rauzy and Brun

$$\Delta = \{(x_1, \dots, x_d) \in \mathbb{R}^d : 1 \geq x_1 \geq \dots \geq x_d \geq 0\}$$

$$T_{\text{BST}}(x_1, \dots, x_d) = \kappa(\text{ord}(1 - \sum_{j=1}^k x_j, x_1, \dots, x_d))$$

if  $0 < 1 - \sum_{j=1}^k x_j < x_{k+1}$  ( $1 \leq k \leq d, x_{d+1} = 1$ )

$d$	$\lambda_2(A_{\text{BST}})$	$1 - \frac{\lambda_2(A_{\text{BST}})}{\lambda_1(A_{\text{BST}})}$	$d$	$\lambda_2(A_{\text{BST}})$	$1 - \frac{\lambda_2(A_{\text{BST}})}{\lambda_1(A_{\text{BST}})}$
2	-0.13648	1.3606	7	-0.02033	1.0729
3	-0.10803	1.2430	8	-0.01175	1.0468
4	-0.07540	1.1817	9	-0.00563	1.0246
5	-0.05035	1.1388	10	-0.00114	1.0054
6	-0.03263	1.1034	11	+0.00224	0.9886

# Comparison between (simulations of) algorithms

uniform approximation coefficients  $1 - \frac{\lambda_2(A)}{\lambda_1(A)} \leq 1 + \frac{1}{d}$

$d$	Selmer	Brun	JP	Intermed.	Garrity	NIJP
2	1.3871	1.3683	1.3735	1.3606	0.6859	<b>1.40145</b>
3	1.1444	1.2203	1.1922	<b>1.2430</b>	0.5798	1.22519
4	0.9866	1.1504	1.1114	<b>1.1817</b>	0.6286	1.14373
5	0.8577	1.1065	1.0676	<b>1.1388</b>	0.7778	1.09786
6	0.7442	1.0746	1.0413	<b>1.1034</b>	0.9468	1.06898
7	0.6437	1.0493	1.0243	<b>1.0729</b>	1.0225	1.04944
8	0.5561	1.0283	1.0127	<b>1.0468</b>	1.0304	1.03551
9	0.4810	1.0102	1.0044	<b>1.0246</b>	1.0189	<b>1.02521</b>
10	0.4173	0.9943	0.9981	<b>1.0054</b>	1.0035	<b>1.01737</b>
11	0.3636	0.9799	<b>0.9933</b>	0.9886	0.9880	<b>1.01125</b>
12						<b>1.00639</b>
13						<b>1.00246</b>
14						<b>0.99924</b>

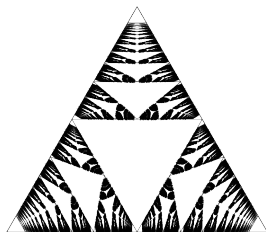
**bold face:** best algorithm or best positive algorithm in dimension  $d$

# The good one: Arnoux–Rauzy algorithm

The **Arnoux–Rauzy algorithm (1991)** is defined by

$$[1 : \alpha : \beta] \mapsto \text{sort}[1 - \alpha - \beta : \alpha : \beta] \quad (1 > \alpha > \beta > 0)$$

on the set  $\{[1 : \alpha : \beta] \in \mathbb{P}^2 : (\alpha, \beta) \in \Delta\}$ .



The Rauzy Gasket  $\Delta$  (taken from **Arnoux and Starosta 2013**)

Avila, Delecroix'19:  $\lambda_2(A_{AR}) < 0$  in all dimensions.

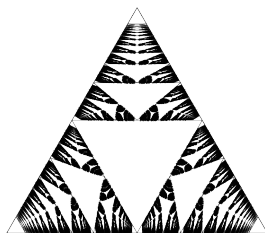


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